

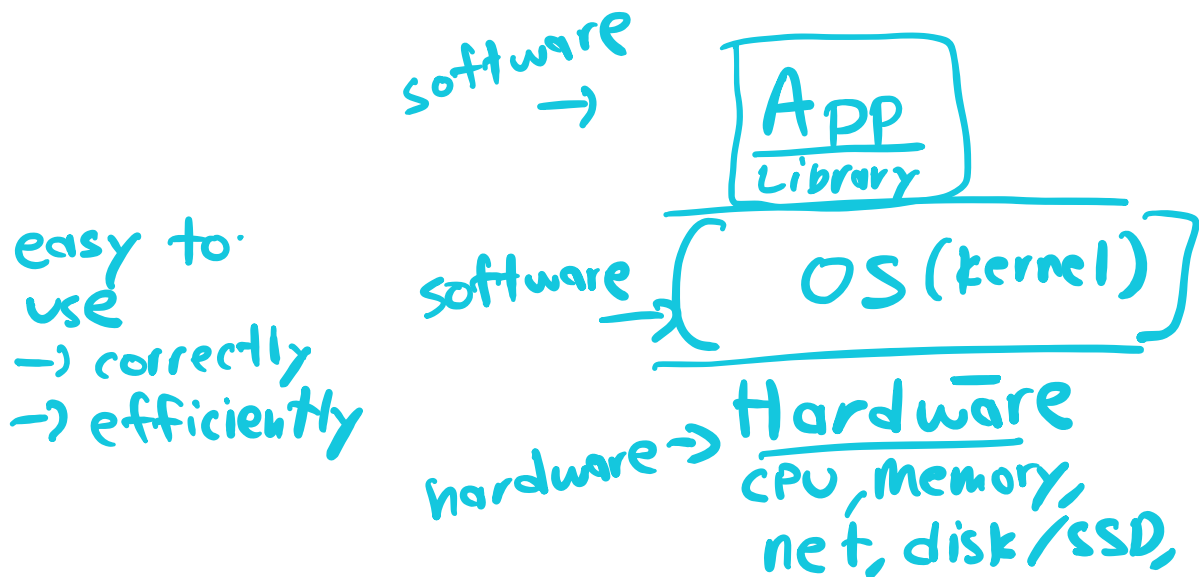
# CS 537 : Day 2

## Operating Systems

Virtualization ←

Concurrency

Persistence



## Virtualization

fact: physical limit in machines

⇒ 1, few CPUs, limited amount of DRAM

build an (illusion)

(=> as many CPUs as needed  
=> for each running program,  
a large memory  
(private)

Virtualization:

CPU

Memory

Goals:

→ Efficiency

→ Security  
isolation

running program or  
"process" ←

P<sub>1</sub> ↔ P<sub>2</sub>

↑ ↓  
OS

Abstraction : Process

What changes when a  
program runs?

1) → registers  
(PC, general purpose)

2)


2) → memory ("address space")

code

heap  
↓

→ I/O (input, output)  
(e.g., open file descriptors)

how to  
CPU: virtualize?



⇒ run  $N$  processes "at once"  
even though we have  $M$  CPUs  
( $N > M$ )

general idea:

1 CPU, 2 processes A, B

↓ ↓  
A B A B ...



"time sharing"

→ Mechanisms: low-level how

→ Policies: which process  
to run?

First attempt: Direct Execution

Boot time: (start up)

⇒ OS is the first prog. to run

set up:

→ free list (to track memory)

→ process list  
(etc.)

want: run one program

OS ⇒ alloc entry on process list

⇒ alloc some memory

⇒ load the program → memory

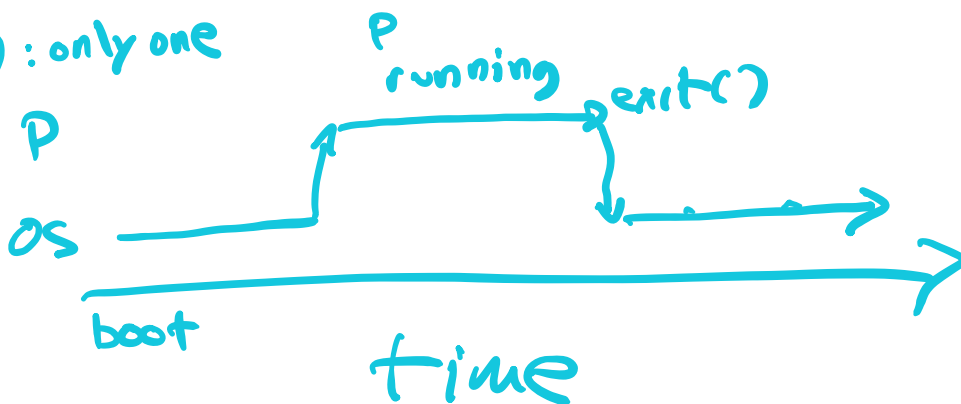
(code, static data)

↘ disk ⇒ memory

⇒ jump to main of program +

start running

CPU: only one



observations:

only one thing running

@ a time

while (1);  
[reboot]

## Problems:

- 1) what if P wants to do something restricted?  
(e.g., issue a read to disk)
- 2) what if OS wants to stop P<sub>A</sub>, run P<sub>B</sub>?
- 3) what if P does something that is "slow"? (disk I/O, net I/O)

## Class

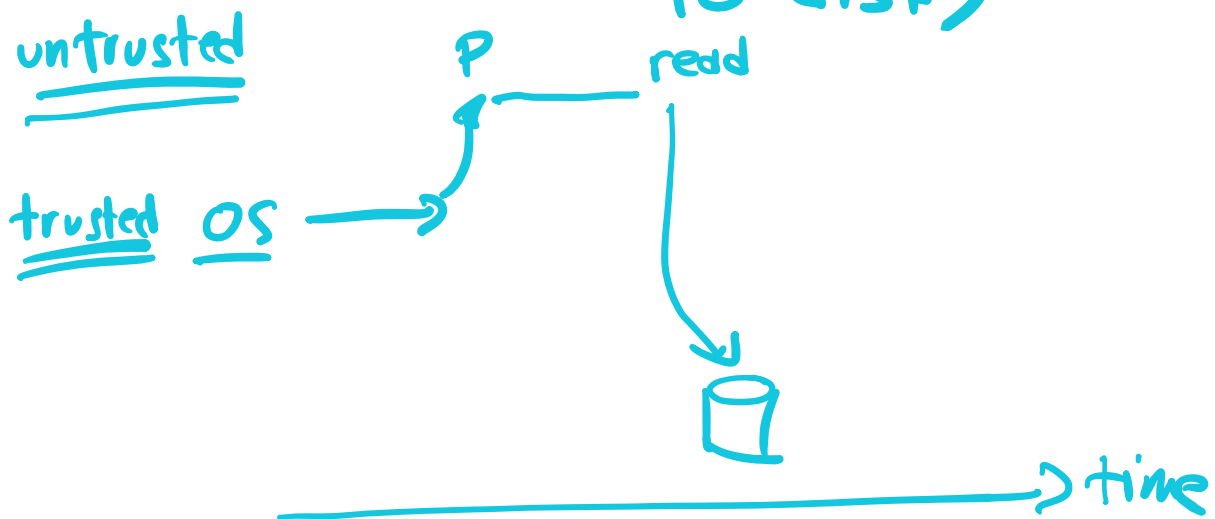
web page ← Canvas

1) ↓  
Piazza → Discord

2) P2d due Monday

3) weekly Canvas quiz

1) What if P wants to do something restricted?  
(e.g., issue a read to disk)



Hardware support:

1) modes of operation (per CPU)

"user mode" (applications)  
restricted, non-privileged

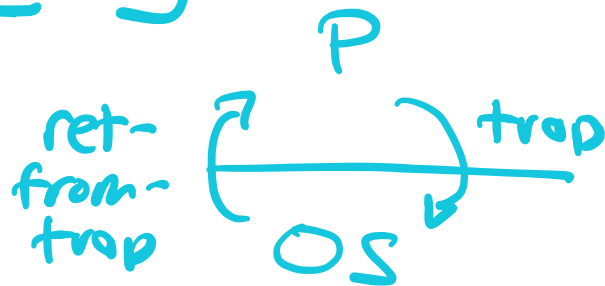
"kernel mode" (OS)

(per CPU) unrestricted, privileged  
bit: 0 or 1 => which mode

2) transition between modes

=> instructions:

user mode → 1) enter kernel mode from user mode: trap  
2) kernel → user: return-from-trap  
kernel mode → 3) set up trap handlers



Time line:

OS @ boot

=> set up  
trap handlers  
(+ other stuff)

Proc @ run time

modern OS:

~ 200 sys calls

# system calls:

open

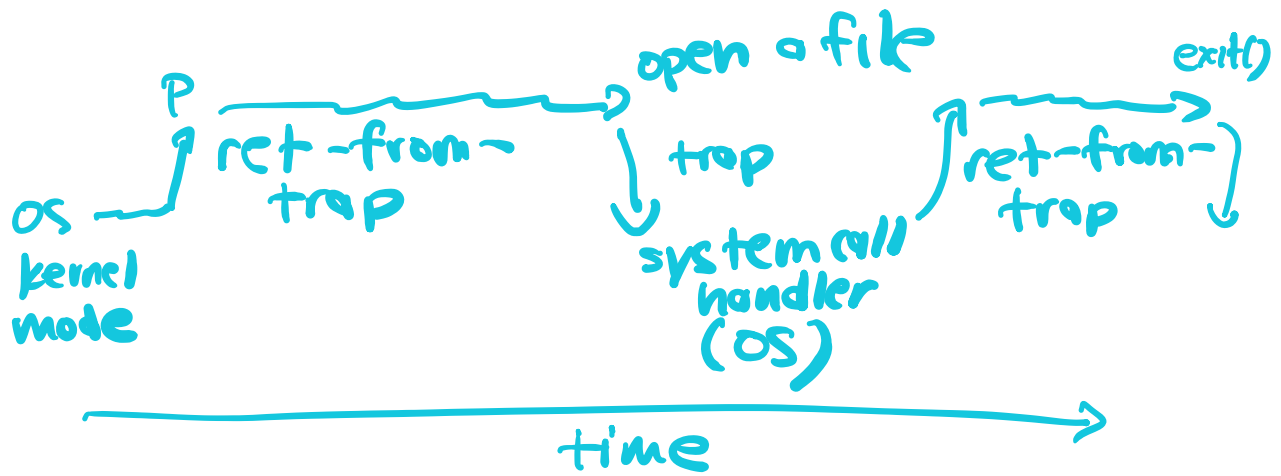
read

close



assembly:

calls trap



=>

Limited Direct Execution