

Today in OS

Last time:

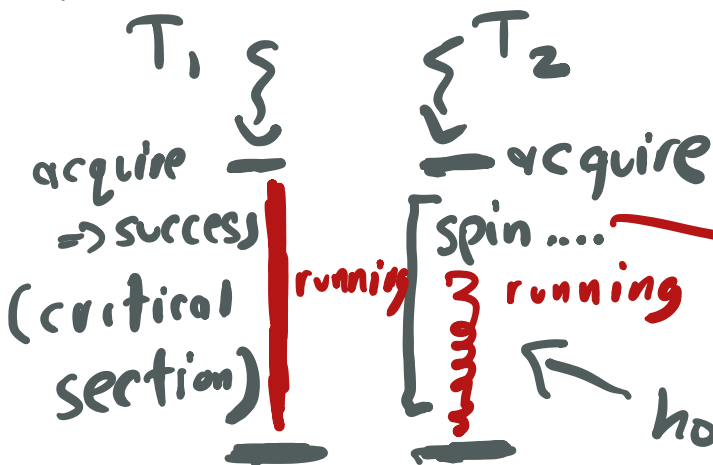
What makes a good lock?

- => correctness
- => overhead
- => fairness

(extreme: starvation)

→ never get lock

Problems: focus on overhead (waste)



spin lock (xchg-based)

```
while (xchg() == 1) yield(); // spin
```

how to reduce waste of spinning

OS support:

⇒ yield() system call

same as READY

when thread calls yield ⇒

moves calling thread from RUNNING

to RUNNABLE

(and push to back

Concurrency (Locks)

⇒ Atomicity

(OS support: new) + hardware support

⇒ Ordering new primitive: condition variable

data

control (when code runs)

CPU: T_1, T_2, \dots, T_{100} of scheduling queue)

acquire

crit. section

timer interrupt

run time slice

T_2 (spins) } 10ms

T_3 spins } 10ms

T_4

...

T_2 yield $\rightarrow T_3$

yield

$\rightarrow T_4 \dots T_{100}$

T_{100}

10ms

worst case analysis:

num of threads x

length of time slice

[waste]

now, with yield:

worst case:

(num of threads x context switch cost)

worst case (yield)

\ll

worst case (spin)

[much faster]

C + Concurrency = SEGV FAULT

Admin

→ midterm scores released

→ solutions coming ASAP
(nice version)

→ P2b : due Monday

(grace period : 2 days)

→ discussion : office hour (not recorded)

Atomicity ↷

Locks : ↷

h/w support : e.g. exchange inst.

OS support : e.g. yield,
park/unpark

Problem #2 :

Ordering ↷

+ control races

(condition variables)