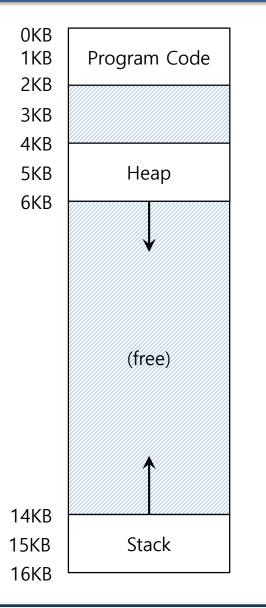
16. Segmentation

Operating System: Three Easy Pieces

Inefficiency of the Base and Bound Approach



- **Big chunk of "free**" space
- "free" space **takes up** physical memory.
- Hard to run when an address space does not fit into physical memory

Segmentation

- Segment is just a contiguous portion of the address space of a particular length.
 - Logically-different segment: code, stack, heap
- **•** Each segment can be **placed** in **different part of physical memory**.
 - Base and bounds exist per each segment.

Placing Segment In Physical Memory

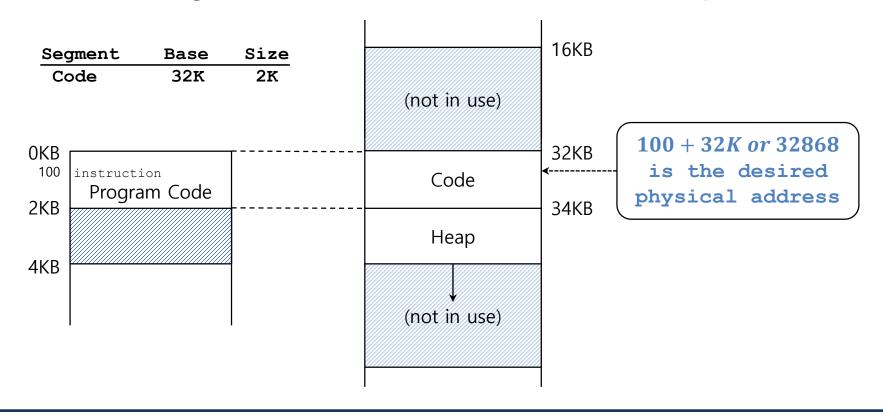
0KB **Operating System** 16KB (not in use) Stack (not in use) 32KB Code Heap 48KB (not in use) 64KB **Physical Memory**

Segment	Base	Size
Code	32K	2К
Heap	34K	2К
Stack	28K	2К

Address Translation on Segmentation

physical address = offset + base

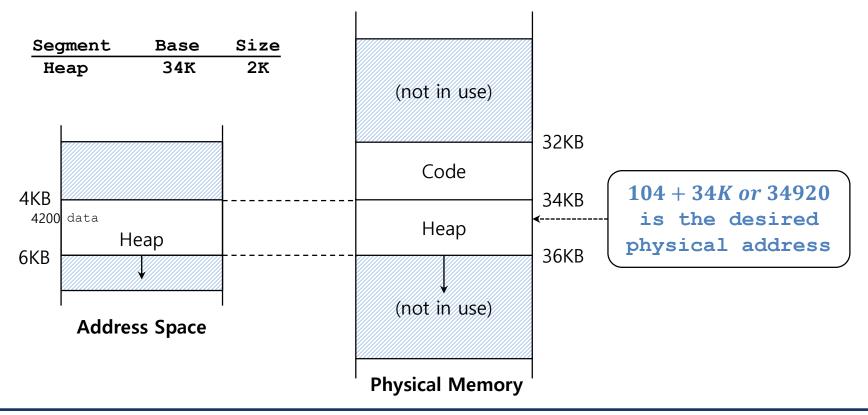
- The offset of virtual address 100 is 100.
 - The code segment **starts at virtual address 0** in address space.



Virtual address + base is not the correct physical address.

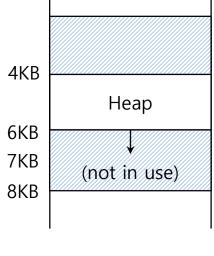
The offset of virtual address 4200 is 104.

• The heap segment starts at virtual address 4096 in address space.



Segmentation Fault or Violation

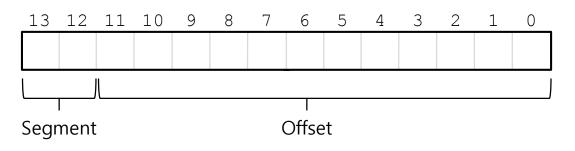
- If an illegal address such as 7KB which is beyond the end of heap is referenced, the OS occurs segmentation fault.
 - The hardware detects that address is **out of bounds**.



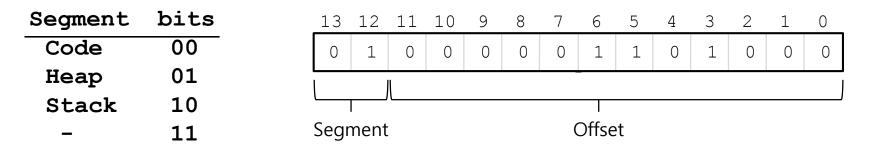
Referring to Segment

Explicit approach

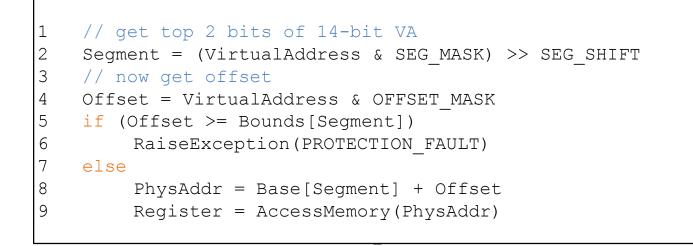
 Chop up the address space into segments based on the top few bits of virtual address.



Example: virtual address 4200 (01000001101000)

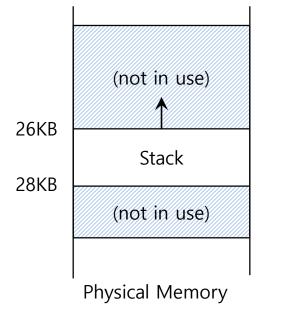


Referring to Segment(Cont.)



- SEG_MASK = 0x3000(110000000000)
- SEG SHIFT = 12
- OFFSET MASK = 0xFFF (001111111111)

- **D** Stack grows **backward**.
- **Extra hardware support** is need.
 - The hardware checks which way the segment grows.
 - 1: positive direction, 0: negative direction



Segment Register(with Negative-Growth Support)

Segment	Base	Size	Grows Positive?
Code	32K	2K	1
Heap	34K	2K	1
Stack	28K	2K	0

Support for Sharing

- **D** Segment can be **shared between address** space.
 - Code sharing is still in use in systems today.
 - by extra hardware support.
- **D** Extra hardware support is need for form of **Protection bits**.
 - A few more bits per segment to indicate permissions of read, write and execute.

Segment	Base	Size	Grows Positive?	Protection
Code	32K	2K	1	Read-Execute
Heap	34K	2K	1	Read-Write
Stack	28K	2K	0	Read-Write

Segment Register Values(with Protection)

Fine-Grained and Coarse-Grained

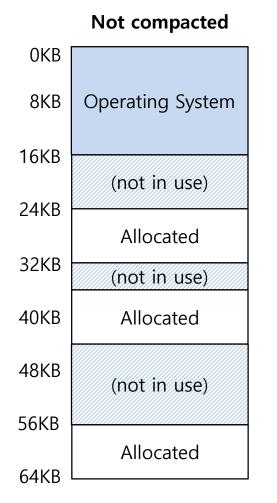
- **Coarse-Grained** means segmentation in a small number.
 - e.g., code, heap, stack.
- Fine-Grained segmentation allows more flexibility for address space in some early system.
 - To support many segments, Hardware support with a segment table is required.

OS support: Fragmentation

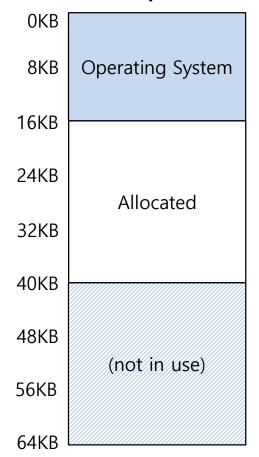
- External Fragmentation: little holes of free space in physical memory that make difficulty to allocate new segments.
 - There is **24KB free**, but **not in one contiguous** segment.
 - The OS cannot satisfy the **20KB request**.

- **Compaction**: **rearranging** the exiting segments in physical memory.
 - Compaction is **costly**.
 - Stop running process.
 - Copy data to somewhere.
 - Change segment register value.

Memory Compaction



Compacted



 Disclaimer: This lecture slide set was initially developed for Operating System course in Computer Science Dept. at Hanyang University. This lecture slide set is for OSTEP book written by Remzi and Andrea at University of Wisconsin.