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of
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CS/ECE 552: Pipelining and Exceptions

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Lecture notes based in part on slides created by
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Exceptions

- What happens?
 - Instruction fetch page fault
 - Illegal opcode
 - Privileged opcode
 - Arithmetic overflow
 - Data page fault
 - I/O device status change
 - Power-on/reset

Exceptions

- For some, we could test for the condition
 - Arithmetic overflow
 - I/O device ready (polling)
- But most tests uselessly say “no”
- Solution:
 - Surprise “procedure call”
 - Called an exception

Exceptions: Big Picture

- Two types:
 - Interrupt (asynchronous) or
 - Trap (synchronous)
- Hardware handles initial reaction
- Then invokes a software exception handler
 - By convention, at e.g., 0xC00
 - O/S kernel provides code at the handler address

Exceptions: Hardware

- Sets state that identifies cause of exception
 - MIPS: in `exception_code` field of Cause register
- Changes to kernel mode for dangerous work ahead
- Disables interrupts
 - MIPS: recorded in status register
- Saves current PC (MIPS: exception PC)
- Jumps to specific address (MIPS: `0x80000080`)
 - Like a surprise JAL – so can't clobber \$31

Exceptions: Software

- Exception handler:
 - MIPS: .ktext at 0x80000080
- Set flag to detect incorrect entry
 - Nested exception while in handler
- Save some registers
- Find exception type
 - E.g., I/O interrupt or syscall
- Jump to specific exception handler

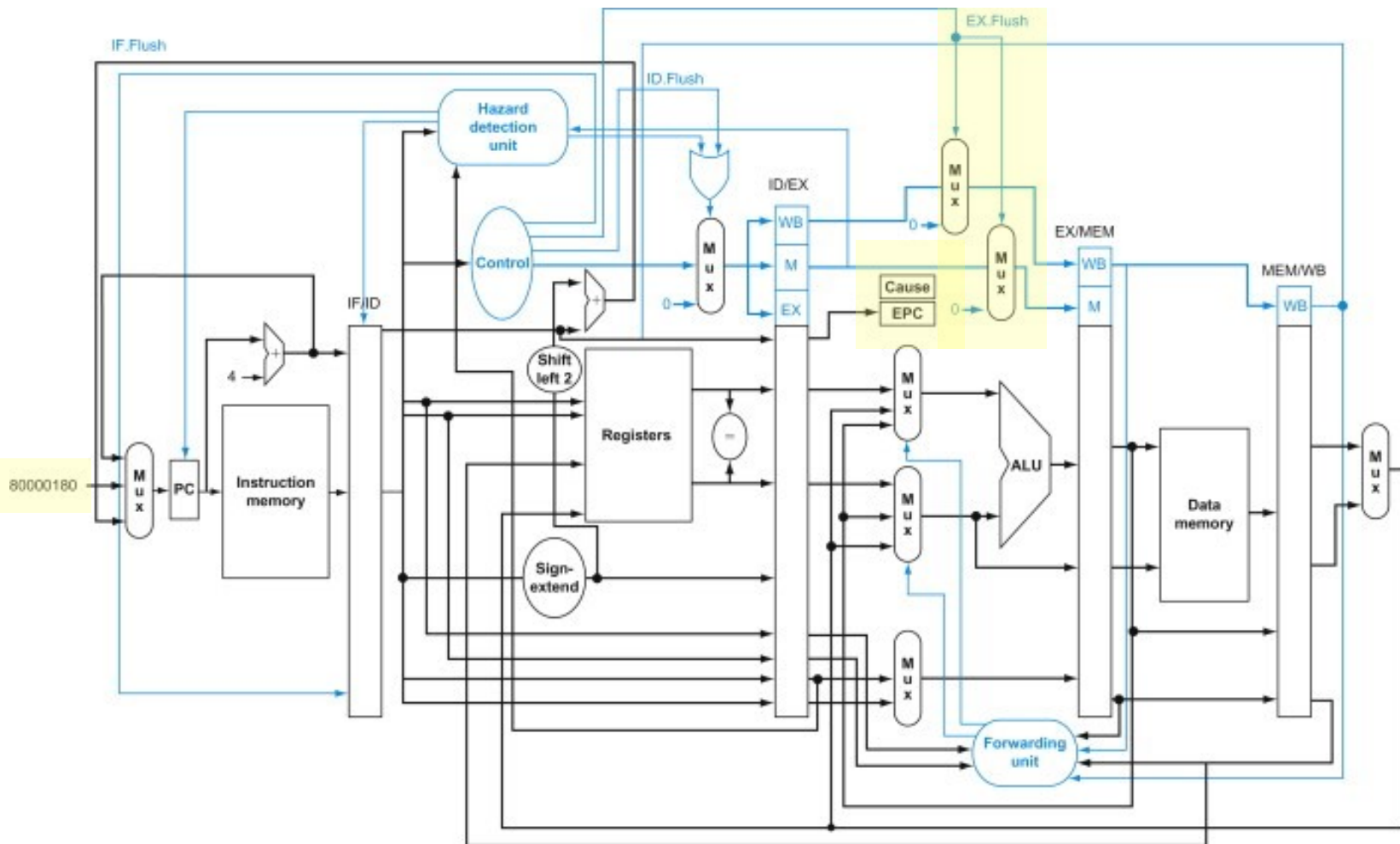
Exceptions: Software, cont'd

- Handle specific exception
- Jump to clean-up to resume user program
- Restore registers
- Reset flag that detects incorrect entry
- Atomically
 - Restore previous mode (user vs. supervisor)
 - Enable interrupts
 - Jump back to program (using EPC)

Implementing Exceptions

- We worry only about hardware, not s/w
- IntCause
 - 0 undefined instruction
 - 1 arithmetic overflow
- Changes to the datapath (e.g., project extra credit)
 - Detect exception
 - Add additional source for next PC (e.g., another 2-1 mux)
 - Storage for exception cause, return address, spare register
- Additional complexity in control logic

Pipeline With Exceptions (Fig 4.66)



Summary and Review

- Exceptions
- Handling exceptions in a pipelined design