CS/ECE 752:
Advanced Computer Architecture I

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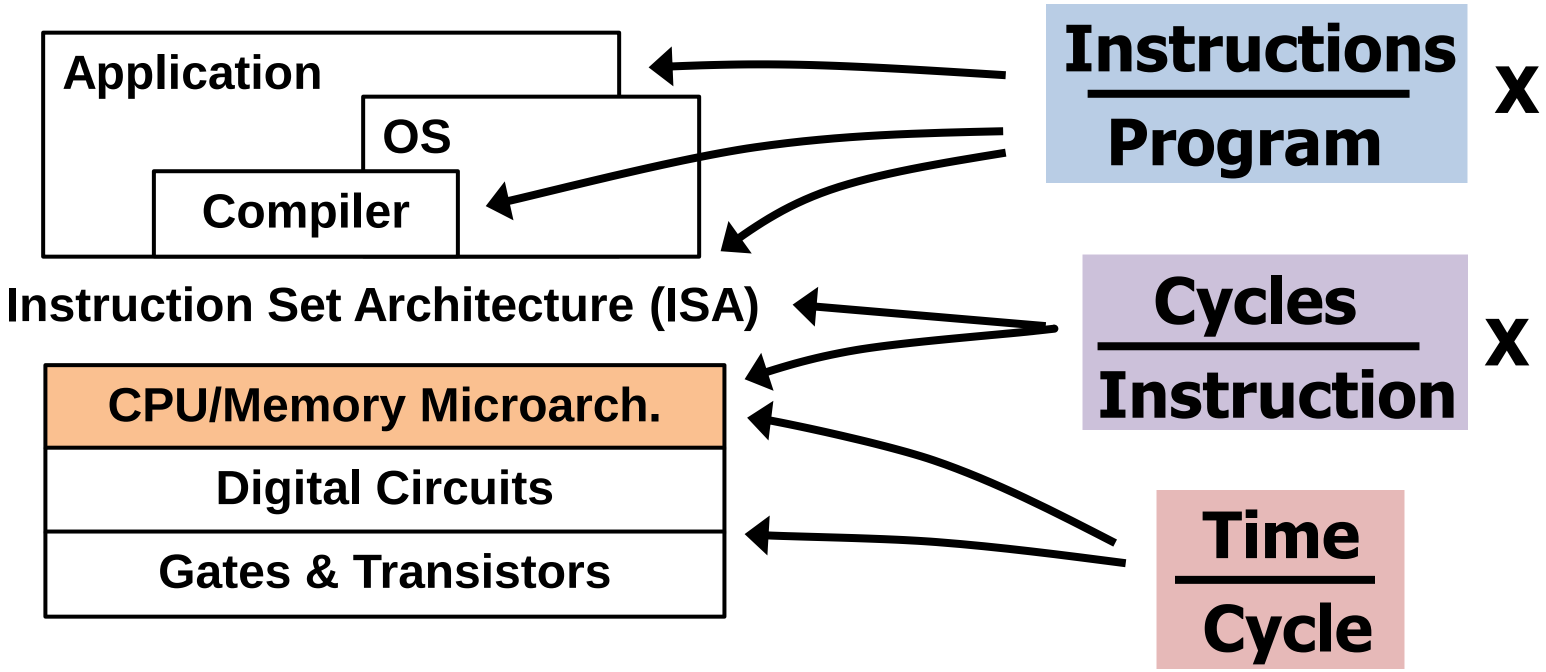
Pipelining

Slide History/Attribution Diagram:

Announcements

- Advanced Topics Lectures Selected
  - Will post Readings and Review shortly
- HW2 Due Friday
  - **See Piazza for issue with gem5 stats – need to cherry-pick and recompile**

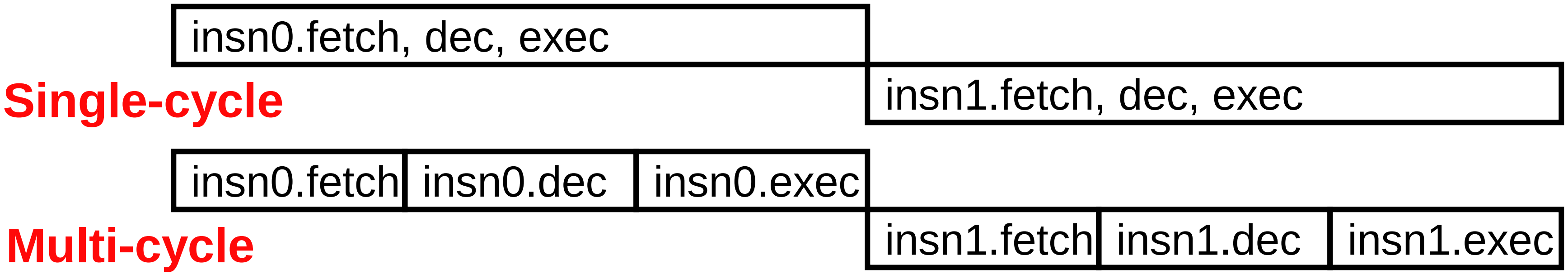
Computer System Layers



This Unit: Pipelining

- Basic Pipelining
  - Single, in-order issue
  - Clock rate vs. IPC
- Data Hazards
  - Hardware: stalling and bypassing
  - Software: pipeline scheduling
- Control Hazards
  - Branch prediction
- Precise state

Single-cycle vs. Multi-cycle



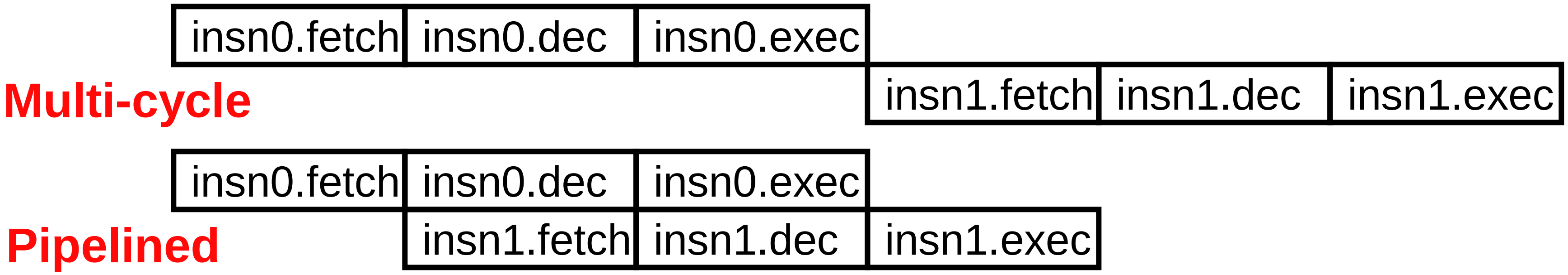
- **Single-cycle datapath:**
  - Fetch, decode, execute one complete instruction every cycle
  - \+ Low CPI: 1 by definition
  - – Long clock period: to accommodate slowest instruction
- **Multi-cycle datapath**: attacks slow clock
  - Fetch, decode, execute one complete insn over multiple cycles
  - \+ Short clock period
  - – High CPI
- Can we have both low CPI and short clock period?
  - Not if datapath executes only one instruction at a time
  - No good way to make a single instruction go faster

Single-cycle vs. Multi-cycle Performance

- Single-cycle
  - Clock period = 50ns, CPI = 1
  - Performance = **50ns/insn**

- Multi-cycle has opposite performance split of single-cycle
  + Shorter clock period
  – Higher CPI
- Multi-cycle
  - Branch: 20% (**3** cycles), load: 20% (**5** cycles), ALU: 60% (**4** cycles)
  - Clock period = **11ns**
    - Why is clock period 11ns and not 10ns?
  - CPI = (20%*3)+(20%*5)+(60%*4) = 4
  - Performance = **44ns/insn**

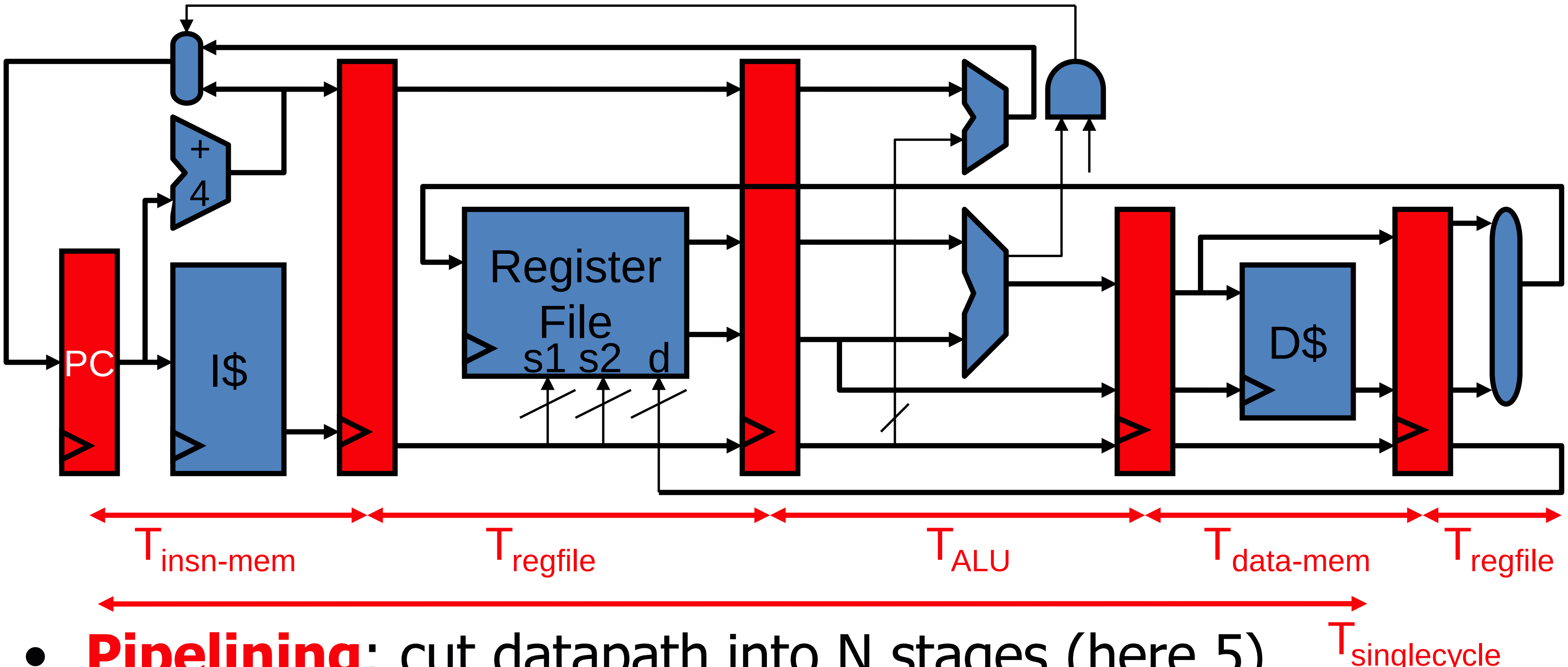
Pipelining



- Important performance technique
  - **Improves instruction throughput rather instruction latency**
- Begin with multi-cycle design
  - When instruction advances from stage 1 to 2
  - Allow next instruction to enter stage 1
  - Form of parallelism: “insn-stage parallelism”
  - Individual instruction takes the same number of stages

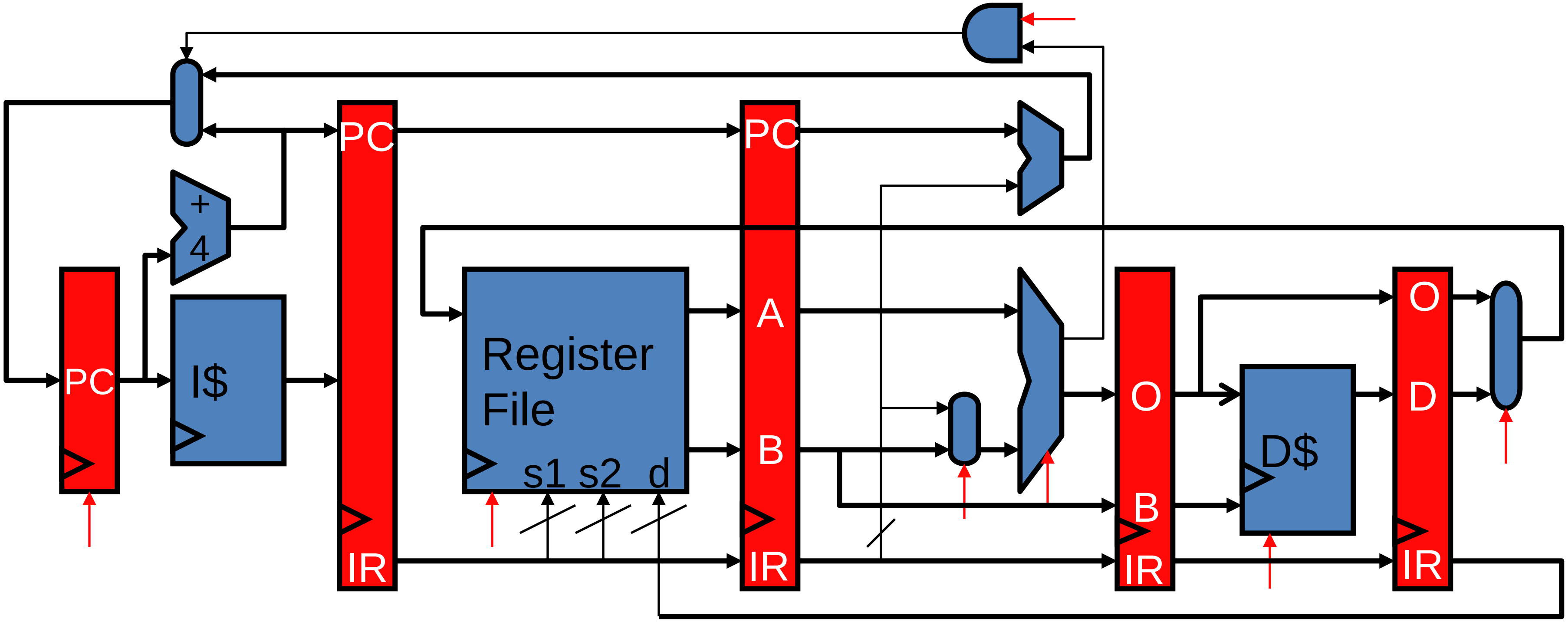
  \+ **But instructions enter and leave at a much faster rate**

Five Stage Pipeline Performance



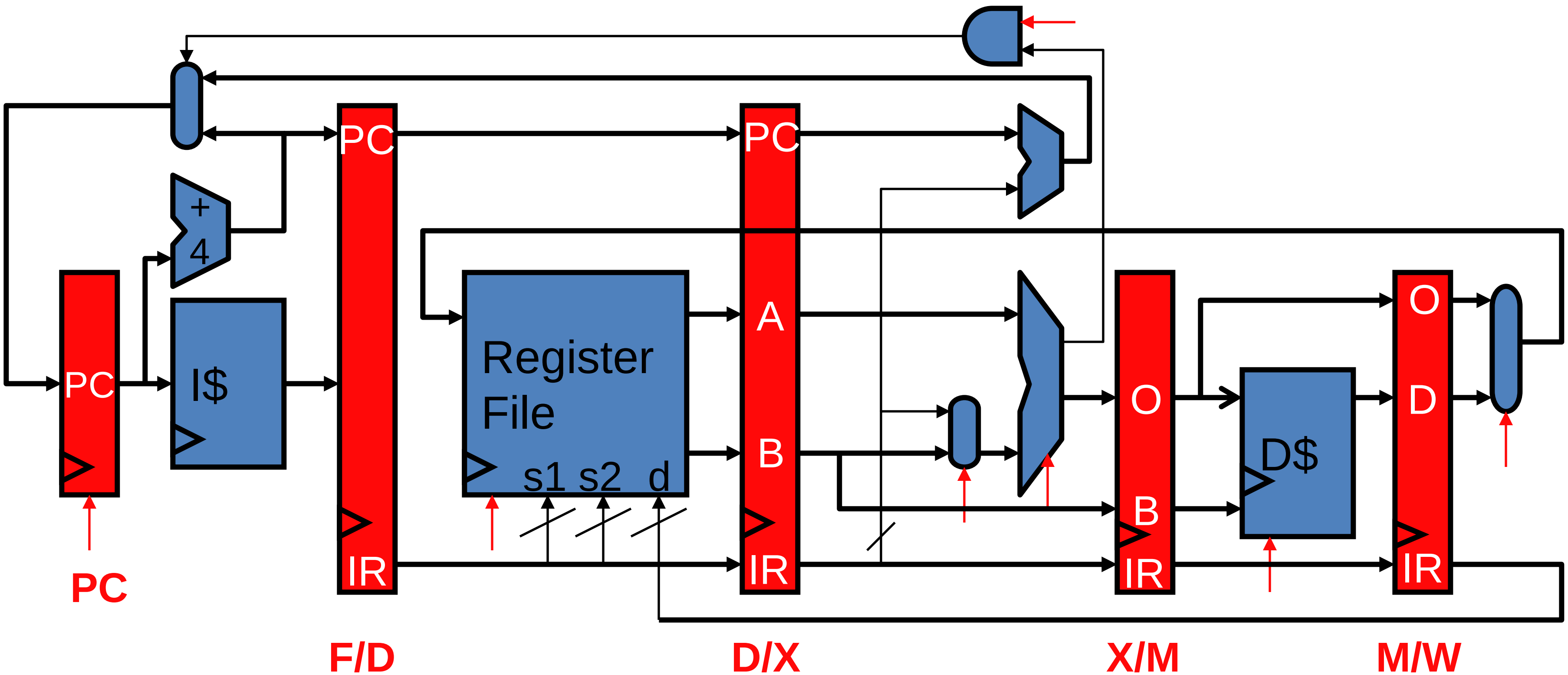
- **Pipelining**: cut datapath into N stages (here 5)
  - One insn in each stage in each cycle
  - \+ Clock period = MAX($T_{insn-mem}$, $T_{regfile}$, T_{ALU}, $T_{data-mem}$)
  - \+ Base CPI = 1: insn enters and leaves every cycle
  - – Actual CPI > 1: pipeline must often stall
  - Individual insn latency increases (pipeline overhead), not the point

5 Stage Pipelined Datapath



- Temporary values (PC,IR,A,B,O,D) re-latched every stage
  - Why? 5 insns may be in pipeline at once, they share a single PC?
  - Notice, PC not latched after ALU stage (why not?)

Pipeline Terminology



- Five stage: **F**etch, **D**ecode, e**X**ecute, **M**emory, **W**riteback
  - Nothing magical about the number 5 (Pentium 4 has 22 stages)
- Latches (pipeline registers) named by stages they separate
  - **PC**, **F/D**, **D/X**, **X/M**, **M/W**

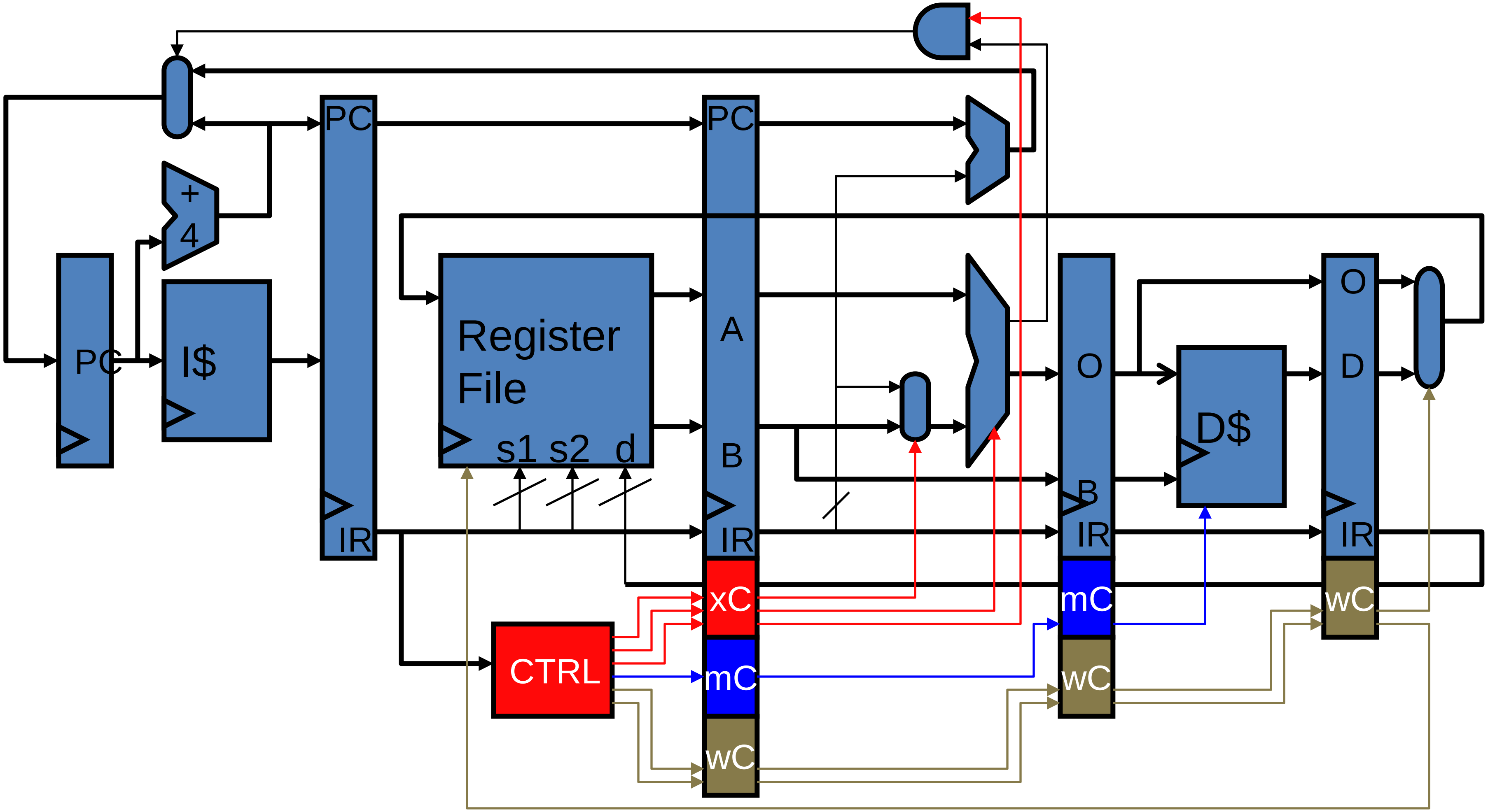
More Terminology & Foreshadowing

- **Scalar pipeline**: one insn per stage per cycle
  - Alternative: “superscalar” (later)

- **In-order pipeline**: insns enter execute stage in order
  - Alternative: “out-of-order” (later)

- **Pipeline depth**: number of pipeline stages
  - Nothing magical about five
  - Contemporary high-performance cores have ~15 stage pipelines
    - (even Intel atom, an in-order core, uses 16 stages)

Pipeline Control



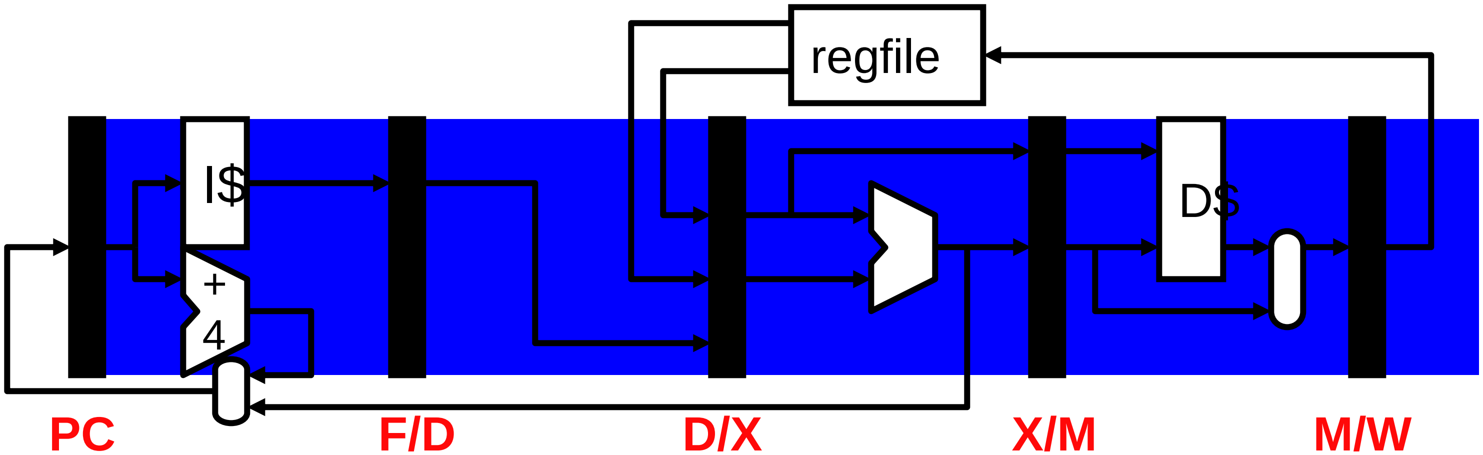
- One single-cycle controller, but pipeline the control signals

Pipeline Diagram

	1	2	3	4	5	6	7	8	9
`add r3,r2,r1`	F	D	X	M	W				
`ld r4,0(r5)`		F	D	**X**	M	W			
`st r6,4(r7)`			F	D	X	M	W		

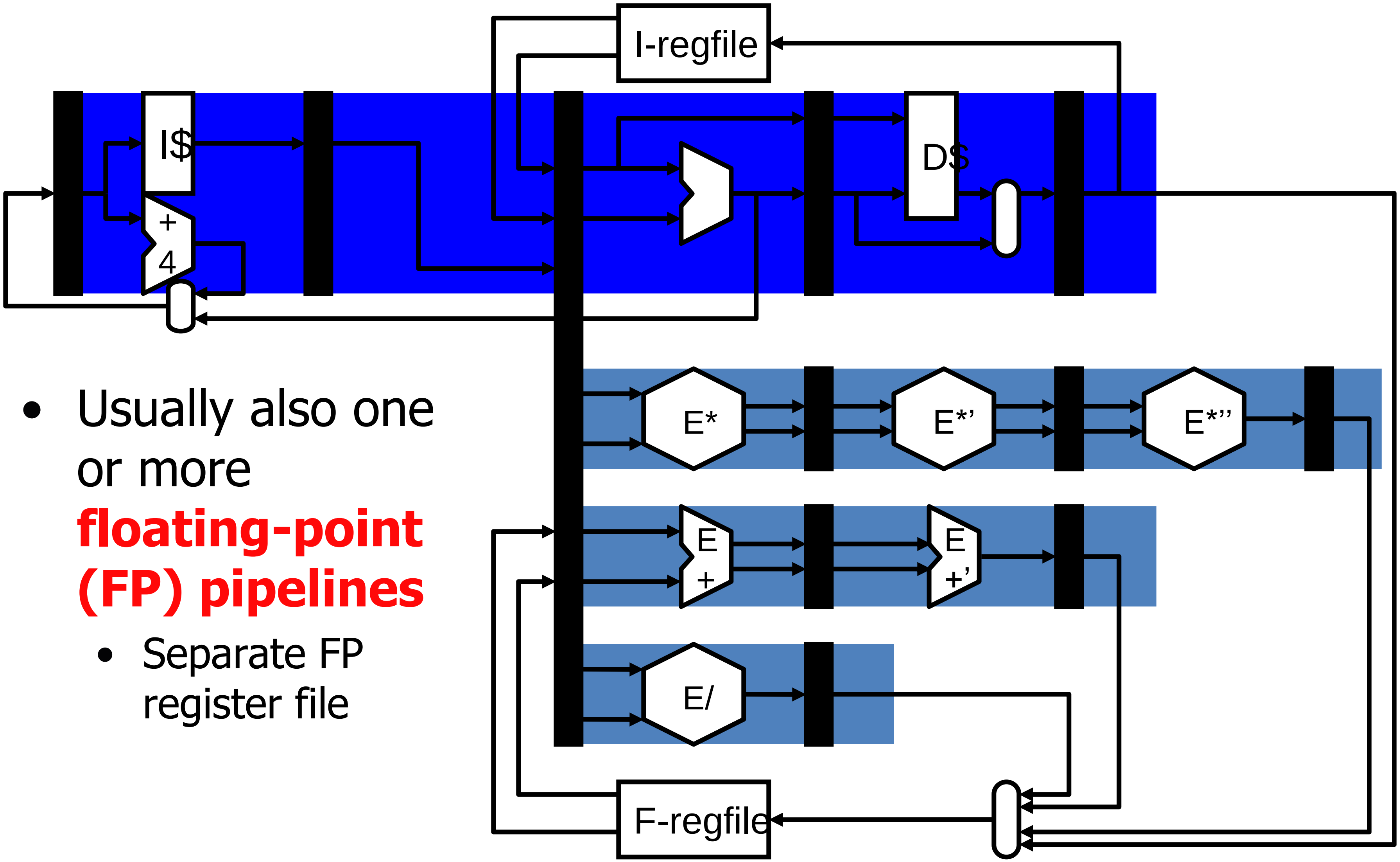
- **Pipeline diagram**
  - Cycles across, insns down
  - Convention: **X** means **`ld r4,0(r5)`** finishes execute stage and writes into X/M latch at end of cycle 4

Abstract Pipeline



- This is an **integer pipeline**
  - Execution stages are X,M,W

Floating Point Pipelines



- Usually also one or more **floating-point (FP) pipelines**
  - Separate FP register file

Pipeline Performance Calculation

- Single-cycle
  - Clock period = 50ns, CPI = 1
  - Performance = 50ns/insn
- Pipelined
  - Clock period = **12ns** (50ns / 5 stages) + overheads
  - Optimistic Model:
    - CPI = **1** (each insn takes 5 cycles, but 1 completes each cycle)
    - Performance = **12ns/insn**
  - Realistic Model: (adds pipeline penalty)
    - CPI = **1.5** (on average insn completes every 1.5 cycles)
    - Performance = **18ns/insn**
    - Much higher performance than single-cycle or multi-cycle

Data Dependences, Pipeline Hazards, and Bypassing

Dependences and Hazards

- **Dependence**: relationship between two insns
  - **Data**: two insns use same storage location
  - **Control**: one insn affects whether another executes at all
  - Programs differ depending on data/control dependences
  - Enforced by making older insn go before younger one
    - Happens naturally in single-/multi-cycle designs
    - But not in a pipeline

- **Hazard**: dependence & possibility of wrong insn order
  - Effects of wrong insn order cannot be externally visible
  - Hazards are a bad thing: stalls reduce performance

Managing a Pipeline

- Proper flow requires two pipeline operations
  - Mess with latch write-enable and clear signals to achieve
- Operation I: **stall**
  - Effect: stops some insns in their current stages
  - Use: make younger insns wait for older ones to complete
  - Implementation: de-assert write-enable
- Operation II: **flush**
  - Effect: removes insns from current stages
  - Use: see later
  - Implementation: assert clear signals
- Both stall and flush must be propagated to younger insns

Structural Hazards

	1	2	3	4	5	6	7	8	9
`ld r2,0(r1)`	F	D	X	**M**	W				
`add r1,r3,r4`		F	D	X	M	W			
`sub r1,r3,r5`			F	D	X	M	W		
`st r6,0(r1)`				**F**	D	X	M	W	

- **Structural hazard**: resource needed twice in one cycle
  - Example: shared I/D$

Fixing Structural Hazards

	1	2	3	4	5	6	7	8	9
`ld r2,0(r1)`	F	D	X	M	W				
`add r1,r3,r4`		F	D	X	M	W			
`sub r1,r3,r5`			F	D	X	M	W		
`and r6,r1,r2`				*	F	D	X	M	W

- Can fix structural hazards by stalling
  - * = structural stall
  - Q: which one to stall: **`ld`** or **`and`**?
    - Always safe to stall younger instruction (here **`and`**)
      - Fetch stall logic: (X/M.op == **`ld`** || X/M.op == **`st`**)
    - But not always the best thing to do performance wise (?)
  + Low cost, simple
  – Decreases IPC
  - Upshot: better to avoid by design than to fix by stalling

Avoiding Structural Hazards

- Pipeline the contended resource
  - + No IPC degradation, low area, power overheads
  - For multi-cycle resources (e.g., multiplier)
  - – Doesn't help for single-cycle resources…
- Replicate the contended resource
  - + No IPC degradation
  - – Increased area, power, latency (interconnect delay?)
  - For cheap, divisible, or highly contended resources (e.g., I$/D$)
- Schedule pipeline to reduce structural hazards (RISC)
  - Design ISA so insn uses a resource at most once
    - Eliminate same insn hazards
  - Always in same pipe stage (hazards between two of same insn)
    - Reason why integer operations forced to go through M stage
  - And always for one cycle

Data Hazards

- Real insn sequences pass values via registers/memory
  - Three kinds of **data dependences** (where's the fourth?)

`add r2,r3➔r1`
`sub r1,r4➔r2`
`or r6,r3➔r1`
	`add r2,r3➔r1`
	`sub r5,r4➔r2`
	`or r6,r3➔r1`
		`add r2,r3➔r1`
		`sub r1,r4➔r2`
		`or r6,r3➔r1`

Read-after-write (RAW)	Write-after-read (WAR)	Write-after-write (WAW)
True-dependence	Anti-dependence	Output-dependence

  - Only one dependence matters between any two insns (RAW has priority)
  - Dependence is property of the program and ISA
- **Data hazards**: function of data dependences and pipeline
  - Potential for executing dependent insns in wrong order
  - Require both insns to be in pipeline ("in flight") simultaneously

RAW

- **Read-after-write (RAW)**

```
add r2,r3→r1
sub r1,r4→r2
or r6,r3→r1
```

- Problem: swap would mean **`sub`** uses wrong value for **`r1`**
- **True**: value flows through this dependence
  - Using different output register for **`add`** doesn't help

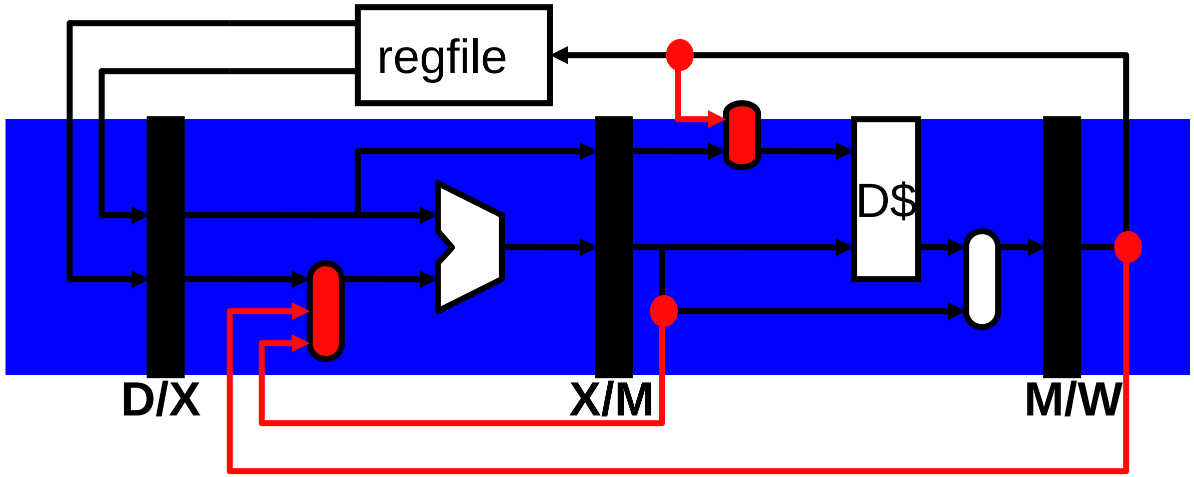
Stall Timing

- Stall Types:
  - data stall,
  - propagated stall
- D and W stages share regfile

	1	2	3	4	5	6	7	8	9	10
add r2,r3➔r1	F	D	X	M	W					
sub r1,r4➔r2		F	**D***	**D***	D	X	M	W		
add r5,r6➔r7			**F***	**F***	F	D	X	M	W	

(assumes RF bypassing: 1st half W writes, 2nd half D reads 2 cycle stall. Also, see backup slides for more on this.)

Reducing RAW Stalls with Bypassing



- Why wait until W stage? Data available after X or M stage
  - **Bypass** (aka **forward**) data directly to input of X or M
    - **X → X**: from beginning of M (X output) to input of X
    - **M → X**: from beginning of W (M output) to input of X
    - **M → M**: from beginning of W (M output) to data input of M
    - **"full bypassing"**:
      - Two each of X → X, M → X (figure shows 1) + M → M =

\+ Reduces stalls in a big way

– Additional wires and muxes may increase clock cycle

Multi Cycle/Pipelined Functional Units

Multiplier Write Port Structural Hazard

- What about…
  - Two instructions trying to write register file in same cycle?
  - Structural hazard!
- Must prevent:

	1	2	3	4	5	6	7	8	9
`mul r3,r5➔r4`	F	D	P0	P1	P2	P3	W		
`addi r1,1➔r6`		F	D	X	M	W			
`add r6,r10➔r7`			F	D	X	M	**W**		

- Solution: stall the offending instruction:

	1	2	3	4	5	6	7	8	9
`mul r3,r5➔r4`	F	D	P0	P1	P2	P3	W		
`addi r1,1➔r6`		F	D	X	M	W			
`add r6,r10➔r7`			F	D	**D***	X	M	**W**	

WAW Hazards

- **Write-after-write (WAW)**

```
add r2,r3→r1
sub r1,r4→r2
or r6,r3→r1
```

- **Artificial**: no value flows through dependence
  - Eliminate using different output register name for `or`
- Compiler effects
  - Scheduling problem: reordering would leave wrong value in `r1`
    - Later instruction reading `r1` would get wrong value
- Pipeline effects
  - Doesn't affect in-order pipeline with single-cycle operations
    - One reason for making ALU operations go through M stage
  - Can happen with multi-cycle operations (e.g., FP or cache misses)

WAW and Precise Interrupts

Optimizing WAW Hazards

	1	2	3	4	5	6	7	8	9	10
`divf f0,f1➔f2`	F	D	E/	E/	E/	E/	E/	**W**		
`stf f2➔(r1)`		F	D	d*	d*	d*	X	M	W	
`addf f0,f1➔f2`			F	D	E+	E+	**W**			

- What to do?
  - Option I: stall younger instruction (**`addf`**) at writeback
    - \+ Intuitive, simple
    - – Lower performance, cascading W structural hazards
  - Option II: cancel older instruction (**`divf`**) writeback
    - \+ No performance loss
    - – What if **`divf`** or **`stf`** cause an exception (e.g., /0, page fault)?

Handling Interrupts/Exceptions

- How are interrupts/exceptions handled in a pipeline?
  - **Interrupt**: external, e.g., timer, I/O device requests
  - **Exception**: internal, e.g., /0, page fault, illegal instruction
  - We care about **restartable** interrupts (e.g. `stf` page fault)

	1	2	3	4	5	6	7	8	9	10
`divf f0,f1➔f2`	F	D	E/	E/	E/	E/	E/	W		
`stf f2➔(r1)`		F	D	D*	D*	D*	X	**M**	W	
`addf f0,f1➔f2`			F	D	E+	E+	W			

- Von Neumann says
  - "Insn execution should appear sequential and atomic"
    - Insn X should complete before instruction X+1 should begin
    - \+ Doesn't physically have to be this way (e.g., pipeline)
    - But be ready to restore to this state at a moments notice
  - Called **precise state** or **precise interrupts**

Handling Interrupts

	1	2	3	4	5	6	7	8	9	10
`divf f0,f1➔f2`	F	D	E/	E/	E/	E/	E/	W		
`stf f2➔(r1)`		F	D	D*	D*	D*	X	**M**	W	
`addf f0,f1➔f2`			F	D	E+	E+	W			

- In this situation
  - Make it appear as if **`divf`** finished and **`stf`**, **`addf`** haven't started
    - Allow **`divf`** to writeback
    - **Flush** **`stf`** and **`addf`** (so that's what a flush is for)
      - But **`addf`** has already written back
      - – Keep an "undo" register file? Complicated
      - – Force in-order writebacks? Slow
    - Invoke exception handler
    - Restart **`stf`**

More Interrupt Nastiness

	1	2	3	4	5	6	7	8	9	10
`divf f0,f1→f2`	F	D	E/	E/	E/	E/	E/	W		
`stf f2→(r1)`		F	D	D*	D*	D*	X	**M**	W	
`divf f0,f4→f2`			F	D	**E/**	E/	E/	E/	E/	W

- What about two simultaneous in-flight interrupts
  - Example: **`stf`** page fault, **`divf`** /0
  - Interrupts must be handled in program order (**`stf`** first)
    - Handler for **`stf`** must see program as if **`divf`** hasn't started
  - Must defer interrupts until writeback **and** force in-order writeback
- In general: interrupts are really nasty
  - Some processors (Alpha) only implement precise integer interrupts
  - Easier because fewer WAW scenarios
  - Most floating-point interrupts are non-restartable anyway
    - **`divf`** /0 → rescale computation to prevent underflow
    - Typically doesn't restart computation at excepting instruction

WAR Hazards

- **Write-after-read (WAR)**

```
add r2,r3→r1
sub r5,r4→r2
or r6,r3→r1
```

- Compiler effects
  - Scheduling problem: reordering would mean `add` uses wrong value for `r2`
  - **Artificial**: solve using different output register name for `sub`
- Pipeline effects
  - Can't happen in simple in-order pipeline
  - Can happen with out-of-order execution

Memory Data Hazards

- So far, have seen/dealt with register dependences
  - Dependences also exist through memory

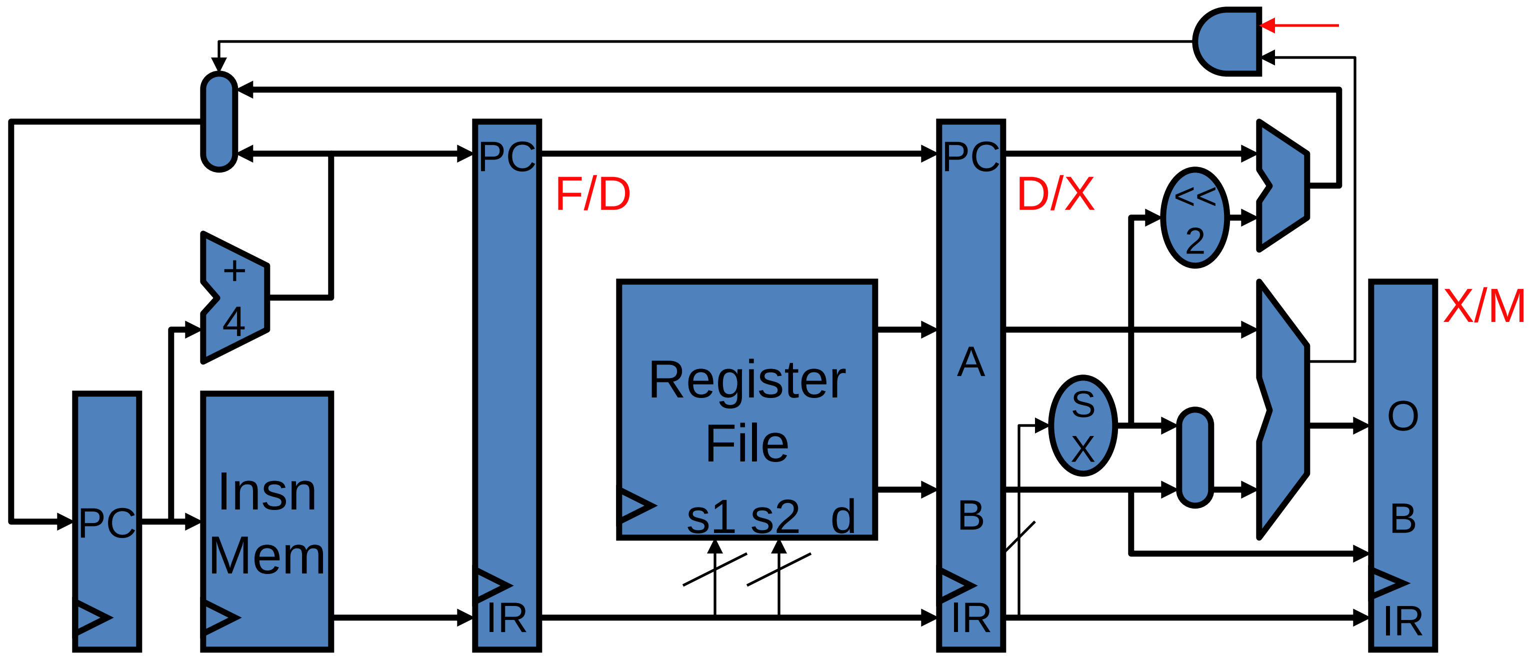
`st r2→(r1)`
`ld (r1)→r4`
`st r5→(r1)`
Read-after-write (RAW)
	`st r2→(r1)`
	`ld (r1)→r4`
	`st r5→(r1)`
	Write-after-read (WAR)
		`st r2→(r1)`
		`ld (r1)→r4`
		`st r5→(r1)`
		Write-after-write (WAW)

  - But in an in-order pipeline like ours, they do not become hazards
  - Memory read and write happen at the same stage
    - Register read happens three stages earlier than register write
  - In general: memory dependences more difficult than register

	1	2	3	4	5	6	7	8	9	10
`st r2→(r1)`	F	D	X	**M**	W					
`ld (r1)→r4`		F	D	X	**M**	W				

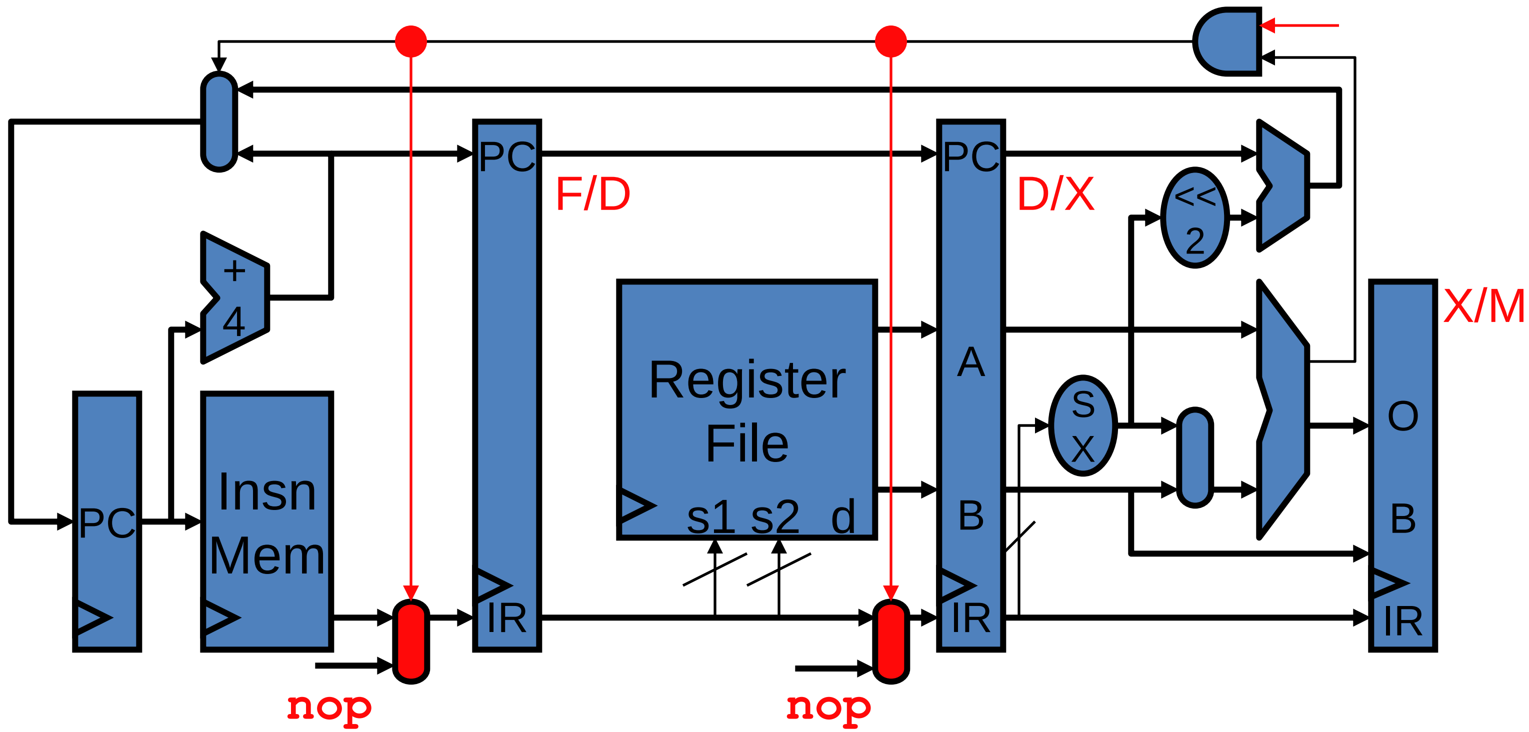
Control Dependences and Branch Prediction

What About Branches?



- **Control hazards options**
  1. Could just stall to wait for branch outcome (two-cycle penalty)
  2. **Fetch past branch insns before branch outcome is known**
     - Default: assume "**not-taken**" (at fetch, can't tell it's a branch)

Branch Recovery



- **Branch recovery**: what to do when branch is actually taken
  - Insns that will be written into F/D and D/X are wrong
  - **Flush them**, i.e., replace them with `nops`
  + They haven't had written permanent state yet (regfile, DMem)
  – Two cycle penalty for taken branches

Control Hazards

- **Control hazards**
  - Control hazards indicated with **F*** (or not at all)
  - Taken branch penalty is 2 cycles

	1	2	3	4	5	6	7	8	9
`addi r1,1➔r3`	F	D	X	M	W				
`bnez r3,targ`		F	D	X	M	W			
`st r6➔[r7+4]`			**F***	**F***	F	D	X	M	W

- Back of the envelope calculation
  - **Branch: 20%**, other: 80%,
  - Say, **75% of branches are taken**
  - $CPI_{BASE} = 1$
  - $CPI_{BASE+BRANCH} = 1 + 0.20*0.75*2 = 1.3$
  - **Branches cause 30% slowdown**
    - Worse with deeper pipelines (higher misprediction penalty)

ISA Branch Techniques

- **Fast branch**: resolves at D, not X
  - Test must be comparison to zero or equality, no time for ALU
  - \+ New taken branch penalty is 1
  - – Must bypass into decode now, too e.g., `cmplt`, `slt`
  - – Complex tests still 2-cycle delay? Or just split into compare + branch?
- **Delayed branch**: branch that takes effect one insn later
  - Insert insns that are independent of branch into "branch delay slot"
  - Preferably from before branch (always helps then)
  - But from after branch OK too
    - As long as no undoable effects (e.g., a store)
  - Upshot: short-sighted feature (e.g., MIPS regrets it)
    - – Not a big win in today's pipelines
    - – Complicates interrupt handling

Big Idea: Speculation

- **Speculation**
  - "Engagement in risky transactions on the chance of profit"

- **Speculative execution**
  - Execute before all parameters known with certainty
- **Correct speculation**
  + Avoid stall, improve performance
- **Incorrect speculation (mis-speculation)**
  – Must abort/flush/squash incorrect instructions
  – Must undo incorrect changes (recover pre-speculation state)

The "game": **[$\%_{correct}$ * gain] > [($1 - \%_{correct}$) * penalty]**

Control Hazards: Control Speculation

- Deal with control hazards with **control speculation**
  - Unknown parameter: are these the correct insns to execute next?
- Mechanics
  - Guess branch target, start fetching at guessed position
  - Execute branch to verify (check) guess
    - Correct speculation? keep going
    - Mis-speculation? Flush mis-speculated insns
  - Don’t write registers or memory until prediction verified
- Speculation game for in-order 5 stage pipeline
  - Gain = 2 cycles
  - Penalty = 0 cycles
    - No penalty → mis-speculation no worse than stalling
  - $\%_{correct}$ = **branch prediction**
    - Static (compiler) ~85%, **dynamic** (hardware) >95%
    - Not much better? Static has 3X mispredicts!

Control Speculation and Recovery

Correct:

	1	2	3	4	5	6	7	8	9
`addi r1,1→r3`	F	D	X	M	W				
`bnez r3,targ`		F	D	X	M	W			
`st r6→(r7+4)`			**F**	**D**	X	M	W		
`targ:add r4,r5→r4`				**F**	D	X	M	W	

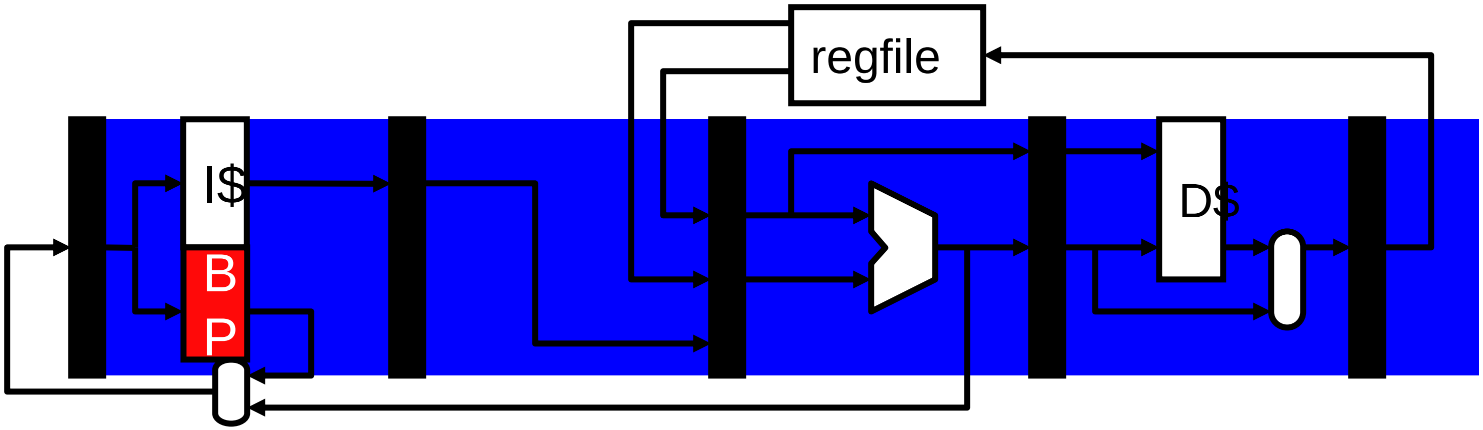
speculative

- **Mis-speculation recovery**: what to do on wrong guess
  - Not too painful in an in-order pipeline
  - Branch resolves in X
  - \+ Younger insns (in F, D) haven't changed permanent state
  - **Flush** insns currently in F/D and D/X (i.e., replace with `nops`)

Recovery:

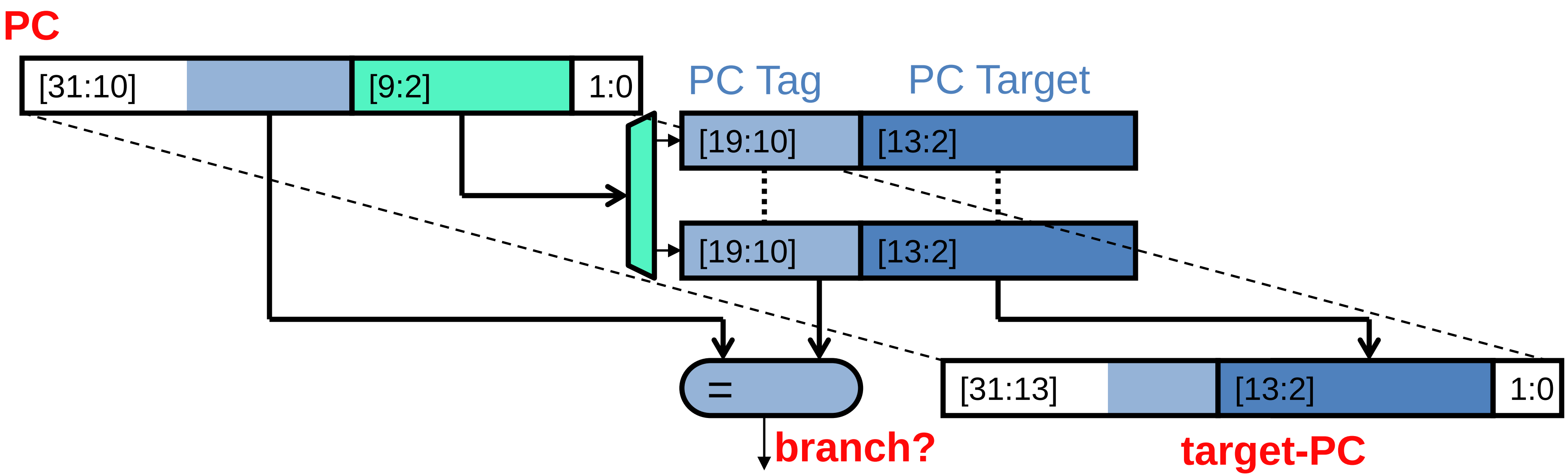
	1	2	3	4	5	6	7	8	9
`addi r1,1→r3`	F	D	X	M	W				
`bnez r3,targ`		F	D	**X**	M	W			
~~`st r6→(r7+4)`~~			**F**	**D**	--	--	--		
~~`targ:add r4,r5→r4`~~				**F**	--	--	--	--	
`targ:add r4,r5→r4`					**F**	D	X	M	W

Dynamic Branch Prediction



- BP part I: **target predictor** **(if taken)**
  - Applies to all control transfers
  - Supplies target PC, tells if insn is a branch prior to decode
  - \+ Easy
- BP part II: **direction predictor**
  - Applies to conditional branches only
  - Predicts taken/not-taken
  - – Harder (or at least more options)

Branch Target Buffer (BTB)



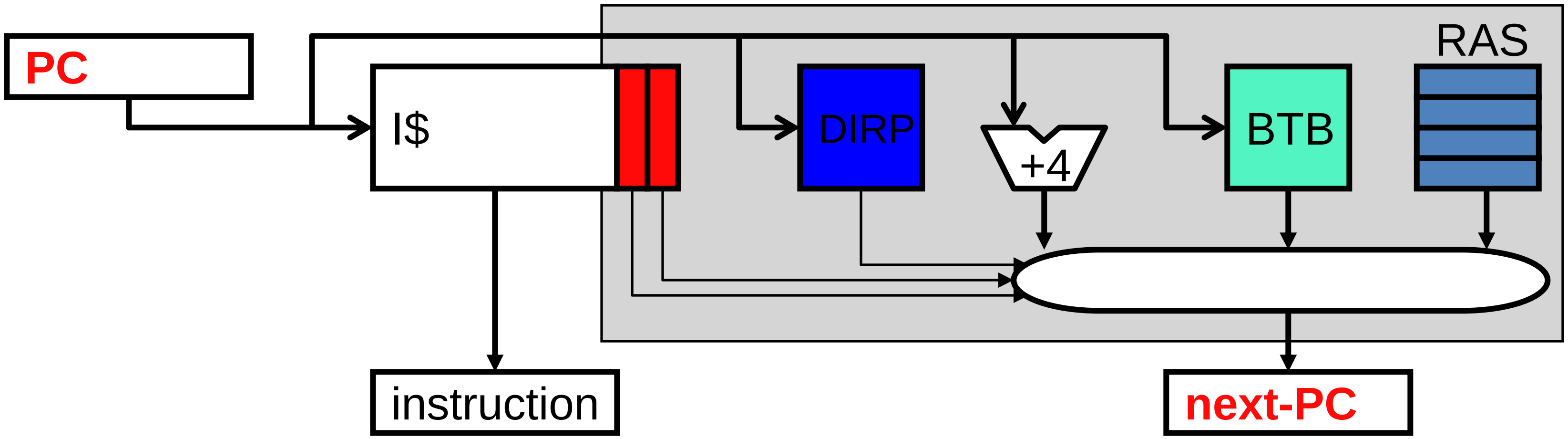
- A small cache: address = PC, data = target-PC
  - Hit? This is a control insn and it's going to target-PC (if "taken")
  - Miss? Not a control insn, or one I have never seen before
- Partial data/tags: full tag not necessary, target-PC is just a guess
  - **Aliasing**: tag match, but not actual match (OK for BTB)
- Insert into BTB when (taken) branch is resolved
- Pentium4 BTB: 2K entries, 4-way set-associative

Why Does a BTB Work?

- Because control insn targets are stable
  - **Direct** means constant target, **indirect** means register target
  - \+ Direct conditional branches? ✓
  - \+ Direct calls? ✓
  - \+ Direct unconditional jumps? ✓

  - \+ Indirect conditional branches? Not that useful→not widely supported
  - Indirect calls? Two idioms:
    - \+ Dynamically linked functions (DLLs)? ✓
    - \+ Dynamically dispatched (virtual) functions? ✓–
  - Indirect unconditional jumps? Two idioms
    - – Switches? ✗ but these are rare
    - – Returns? ✓– but… we should know based on the program where we are returning!

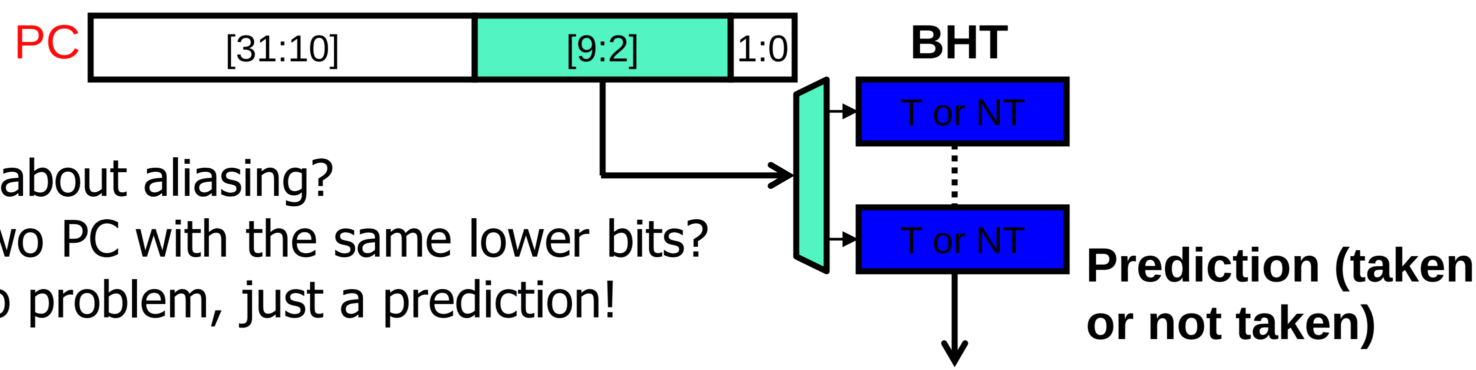
Return Address Stack (RAS)



- Return addresses are easy to predict without a BTB
  - Hardware **return address stack (RAS)** tracks call sequence
  - Calls push PC+4 onto RAS
  - Prediction for returns is RAS[TOS]
  - Q: how can you tell if an insn is a return before decoding it?
  - A1: Add tags to make RAS a cache (have to check it…)
  - A2: (Better) attach **pre-decode bits** to I$
    - Written after first time insn executes
    - Two useful bits: return?, conditional-branch?

Branch Direction Prediction

- **Direction predictor (DIRP)**
  - Map conditional-branch PC to taken/not-taken (T/N) decision
  - Can be based on additional information
- **Branch history table (BHT)**: simplest predictor
  - PC indexes table of bits (0 = N, 1 = T), no tags
  - Essentially: branch will go same way it went last time



  - What about aliasing?
    - Two PC with the same lower bits?
    - No problem, just a prediction!
- Why: Individual conditional branches often biased or weakly biased
  - 90%+ one way or the other considered **"biased"**
  - Why? Loop back edges, checking for uncommon conditions

Branch History Table (BHT)

- Problem: **inner loop branch** below

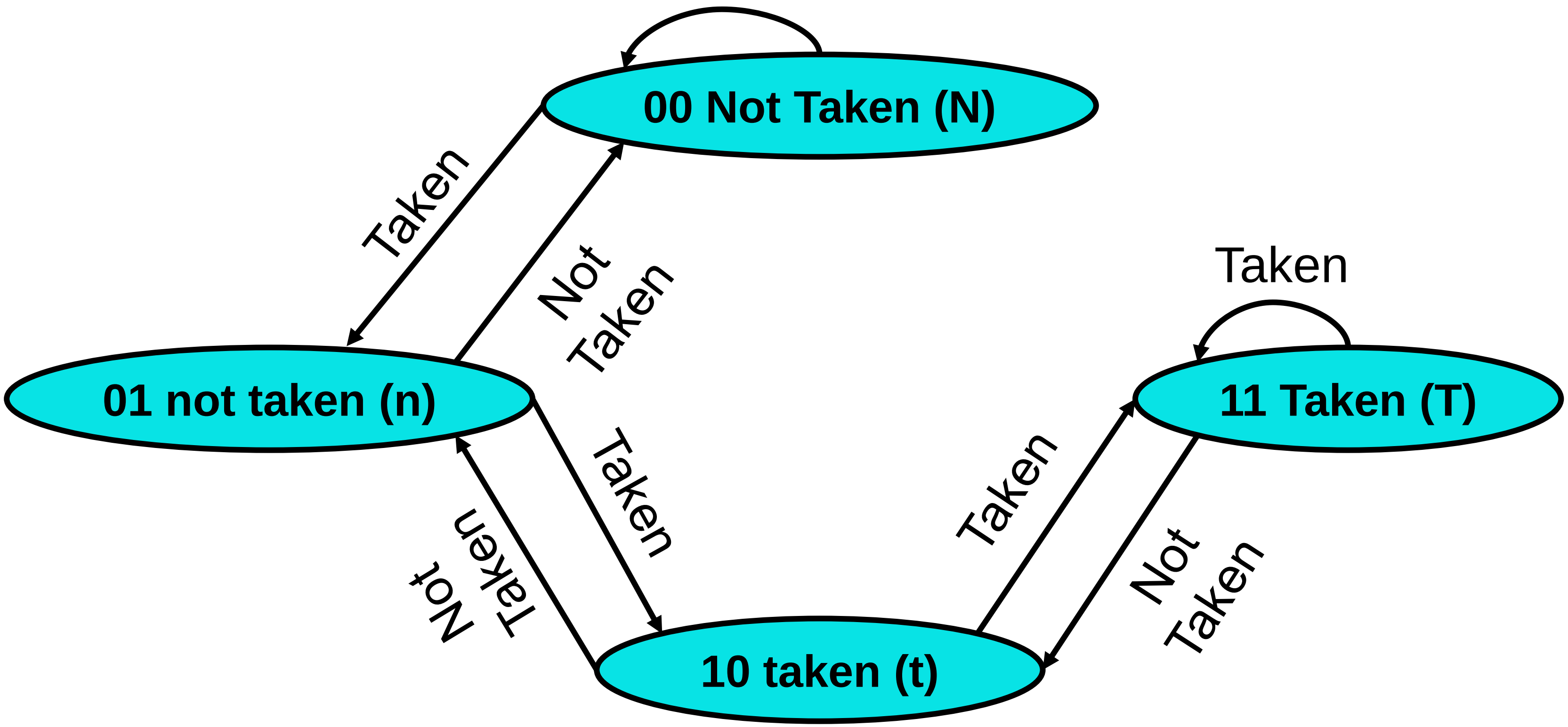
```
for (i=0; i<100; ++i)
    for (j=0; j<3; ++j)
        // whatever
```

  - Two “built-in” mis-predictions per inner loop iteration
  - Branch predictor “changes its mind too quickly”

Time	Outcome
1	**T**
2	**T**
3	**T**
4	**N**
5	**T**
6	**T**
7	**T**
8	**N**
9	**T**
10	**T**
11	**T**
12	**N**

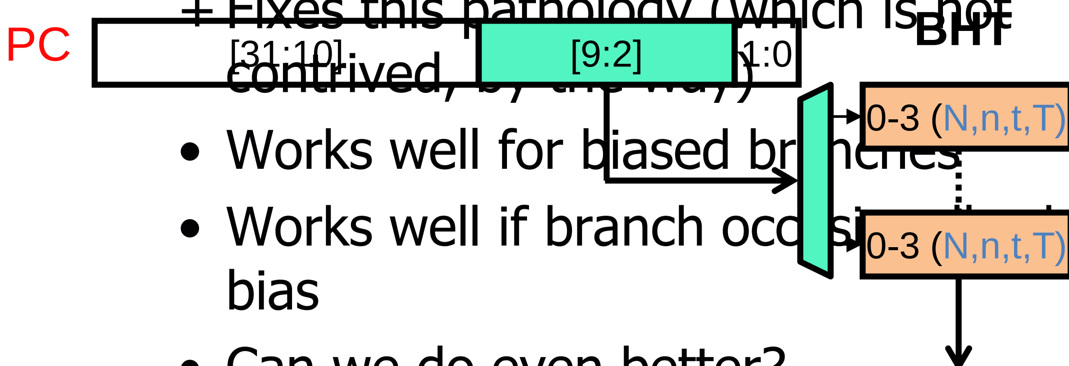
Two-Bit Saturating Counters (2bc)

- **Two-bit saturating counters (2bc)** [Smith 1981]
  - Replace each single-bit prediction
    - (0,1,2,3) = (N,n,t,T)
    - **Strong not-taken**, weak not-taken, weak taken, **strong taken**



Two-Bit Saturating Counters (2bc)

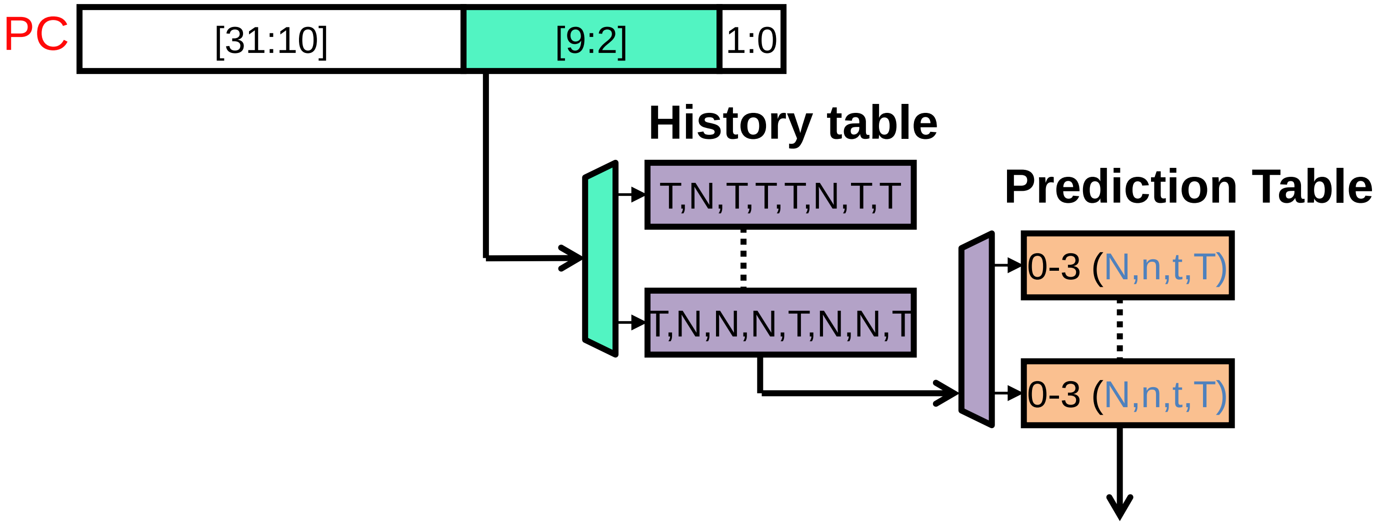
- **Two-bit saturating counters (2bc)** [Smith 1981]
  - Replace each single-bit prediction
    - (0,1,2,3) = (N,n,t,T)
  - Adds "hysteresis"
    - Force predictor to mis-predict twice before "changing its mind"
  - One mispredict each loop execution (rather than two)
    - + Fixes this pathology (which is not contrived, by the way)
    - Works well for biased branches
    - Works well if branch occasionally changes bias
    - Can we do even better?



Time	State	Prediction	Outcome	Result?
1	N	N	T	Wrong
2	n	N	T	Wrong
3	t	T	T	Correct
4	T	T	N	Wrong
5	t	T	T	Correct
6	T	T	T	Correct
7	T	T	T	Correct
8	T	T	N	Wrong
9	t	T	T	Correct
10	T	T	T	Correct
11	T	T	T	Correct
12	T	T	N	Wrong

Two-level Predictor

- **Correlated (two-level) predictor** [Patt 1991]
  - Exploits observation that branch outcomes are correlated
  - Branch history table stores past branches



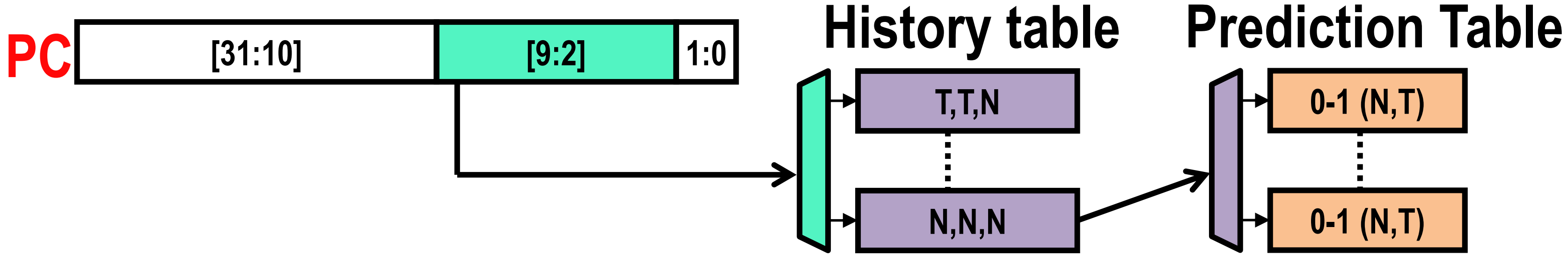
Correlated Predictor – 3 Bit History

```
for (i=0;i<100;i++)
   for (j=0;j<3;j++)
      // whatever
```

- **Actual Pattern:**

  T,T,T,N,T,T,T,N,T

  …
- **Want:**
  - **T,T,T ->**
  - **N,T,T ->**
  - **T,N,T ->**
  - **T,T,N ->**



Correlated Predictor Design

- Design choice I: one **global** BHR or one per PC (**local**)?
  - Each one captures different kinds of patterns
  - Global captures local patterns for tight loop branches

- Design choice II: how many history bits (BHR size)?
  - Tricky one

  \+ Given unlimited resources, longer BHRs are better, but…

  – BHT utilization decreases

    – Many history patterns are never seen

    – Many branches are history independent (don't care)

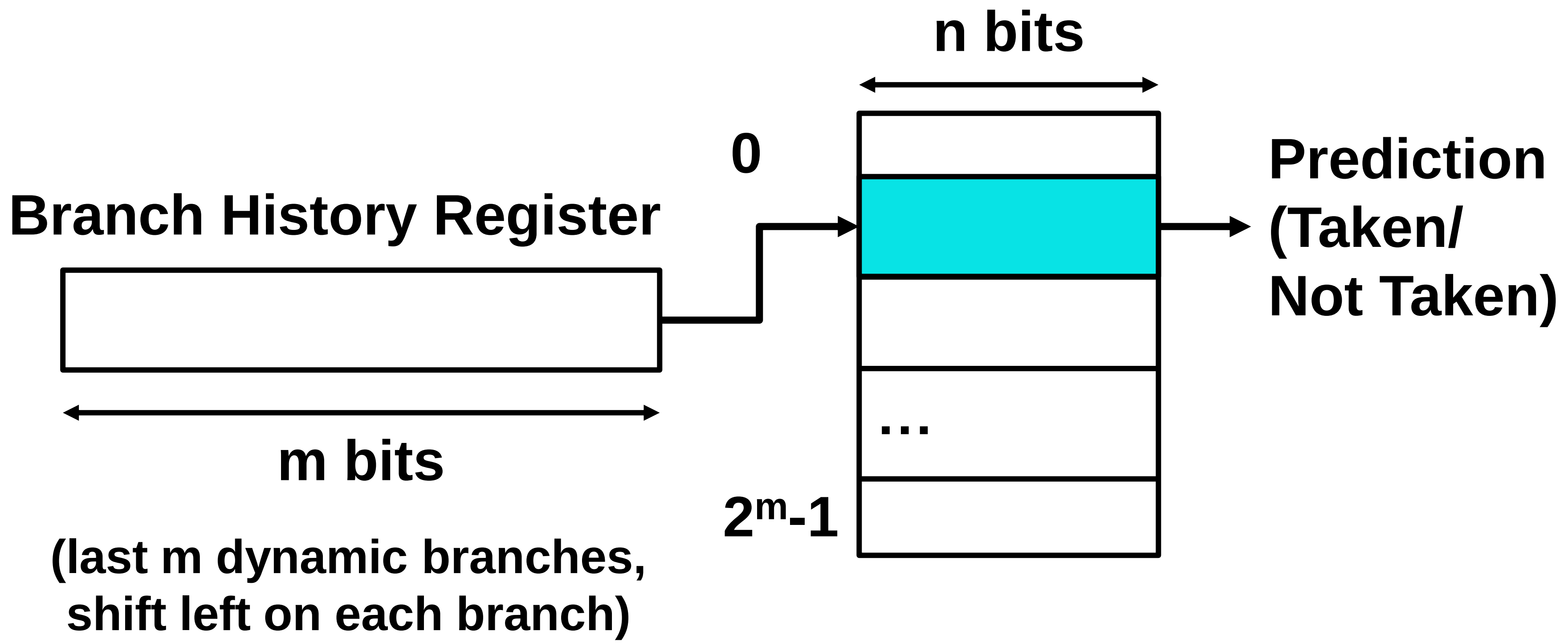
    - PC xor BHR allows multiple PCs to dynamically share BHT
    - BHR length < $\log_2$(BHT size)

  – Predictor takes longer to train

  - Typical length: 8–12

(m,n) Correlated Predictor

- Generalizing, an (m,n) predictor is:
  - N = n-bit saturating counter
    - 2^n counters that can be indexed
  - M = m-bit global history register
    - 2^m locations per PC (e.g., in BHT)



Branch Prediction Performance

- Same parameters
  - **Branch: 20%**, load: 20%, store: 10%, other: 50%
  - 75% of branches are taken
- Dynamic branch prediction
  - Branches predicted with 95% accuracy
  - CPI = 1 + 0.20*0.05*2 = **1.02**

- **So are we done with branch prediction?**
  - **No, not yet … penalties for out-of-order core are VERY HIGH even with 5% mispredictions**

Pipeline Performance Summary

- Base CPI is 1, but hazards increase it
- Nothing magical about a 5 stage pipeline
  - Pentium4 has 22 stage pipeline
- Increasing **pipeline depth**
  - \+ Increases clock frequency (that’s why companies used to do it)
  - – But decreases IPC
    - Branch mis-prediction penalty becomes longer
      - More stages between fetch and whenever branch computes
    - Non-bypassed data hazard stalls become longer
      - More stages between register read and write
  - Ultimate metric is *IPC * frequency*
  - At some point, CPI losses offset clock gains

Dynamic Pipeline Power

- Remember control-speculation game
  - [2 cycles * $\%_{correct}$] – [**0 cycles** * (1–$\%_{correct}$)]
  - No penalty → mis-speculation no worse than stalling
  - This is a performance-only view
  - From a power standpoint, mis-speculation is worse than stalling

- **Power control-speculation game**
  - [0 nJ * $\%_{correct}$] – [**X nJ** * (1–$\%_{correct}$)]
  - No benefit → correct speculation no better than stalling
    - Not exactly, increased execution time increases static power
  - How to balance the two?

Trends…

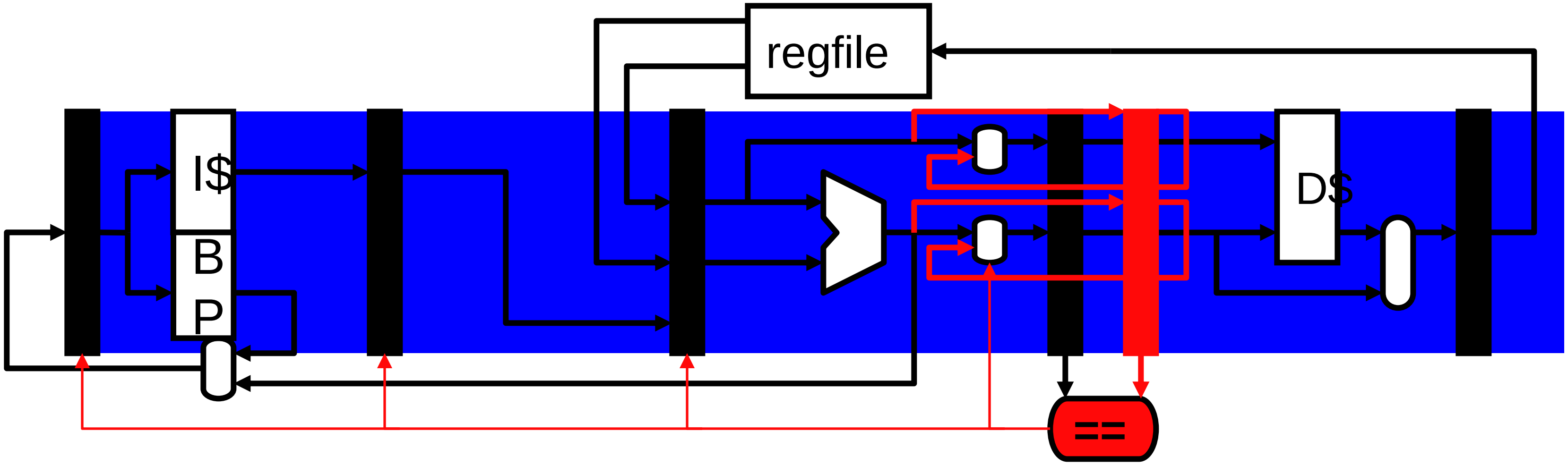
- Trend has been for deeper pipelines
  - Intel example:
    - 486: 5 stages (50+ gate delays / clock)
    - Pentium: 7 stages
    - Pentium II/III: 12 stages
    - Pentium 4: 22 stages (10 gate delays / clock)
    - 800 MHz Pentium III was faster than 1 GHz Pentium4
    - Intel Core2: 14 stages, less than Pentium 4
    - Nehalem (2008): 20-24 Stages
    - Haswell (2013): 14-19 Stages
    - Skylake (2017): 14-19 Stages
    - Cooper Lake (2019): 14-19 Stages

Summary

- Principles of pipelining
  - Effects of overhead and hazards
  - Pipeline diagrams
- Data hazards
  - Stalling and bypassing
- Control hazards
  - Branch prediction
- Power techniques
  - Dynamic power: speculation gating
  - Static and dynamic power: razor latches

Hidden Bonus Slides

Research: Razor



- **Razor** [Uht, Ernst+]
  - Identify pipeline stages with narrow signal margins (e.g., **X**)
  - Add **"Razor"** X/M latch: relatches X/M input signals after safe delay
  - Compare X/M latch with "safe" razor X/M latch, different?
    - Flush F,D,X & M
    - Restart M using X/M razor latch, restart F using D/X latch

  \+ Pipeline will not "break" → reduce V_{DD} until flush rate too high

  \+ Alternatively: "over-clock" until flush rate too high

When to Perform Branch Prediction?

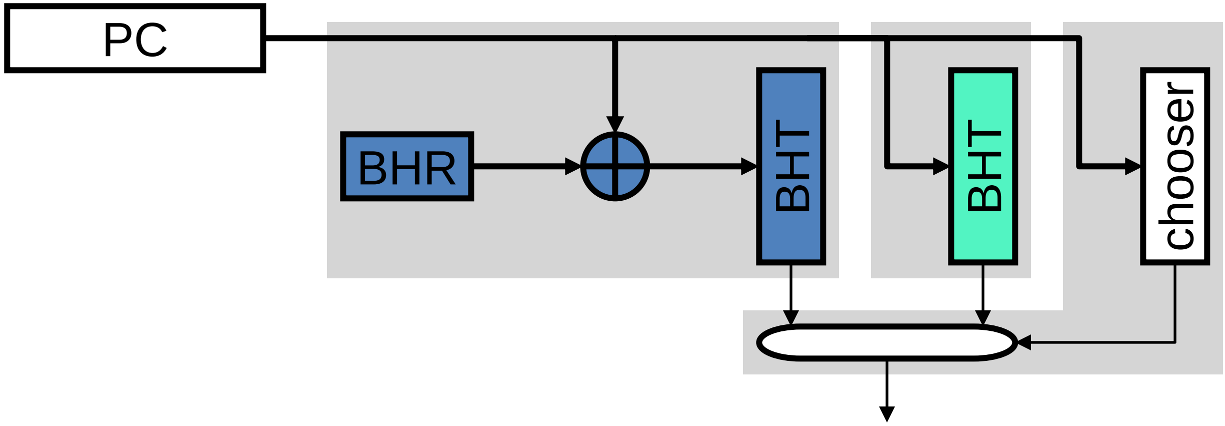
- Option #1: During Decode
  - Look at instruction opcode to determine branch instructions
  - Can calculate next PC from instruction (for PC-relative branches)
  – One cycle "mis-fetch" penalty **even if branch predictor is correct**

	1	2	3	4	5	6	7	8	9
`bnez r3,targ`	F	**D**	X	M	W				
`targ:add r4,r5,r4`			**F**	D	X	M	W		

- Option #2: During Fetch?
  - How do we do that?

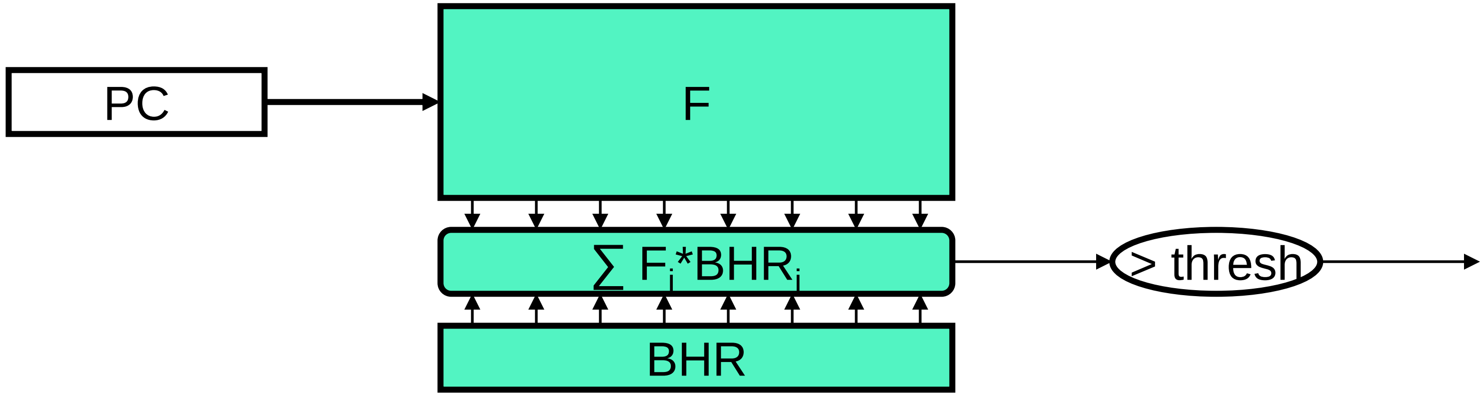
Hybrid Predictor

- **Hybrid (tournament) predictor** [McFarling 1993]
  - Attacks correlated predictor BHT capacity problem
  - Idea: combine two predictors
    - **Simple BHT** predicts history independent branches
    - **Correlated predictor** predicts only branches that need history
    - **Chooser** assigns branches to one predictor or the other
    - Branches start in simple BHT, move mis-prediction threshold
  - \+ Correlated predictor can be made **smaller**, handles fewer branches
  - \+ 90–95% accuracy



Research: Perceptron Predictor

- **Perceptron predictor** [Jimenez]
  - Attacks BHR size problem using machine learning approach
  - BHT replaced by table of function coefficients F_i (signed)
  - Predict taken if $\sum(BHR_i*F_i)>$ threshold
  + Table size #PC*|BHR|*|F| (can use long BHR: ~60 bits)
    – Equivalent correlated predictor would be $\#PC*2^{|BHR|}$
  - How does it learn? Update F_i when branch is taken
    - $BHR_i == 1\ ?\ F_i{+}{+} : F_i{-}{-};$
    - "don't care" F_i bits stay near 0, important F_i bits saturate
  + Hybrid BHT/perceptron accuracy: 95–98%

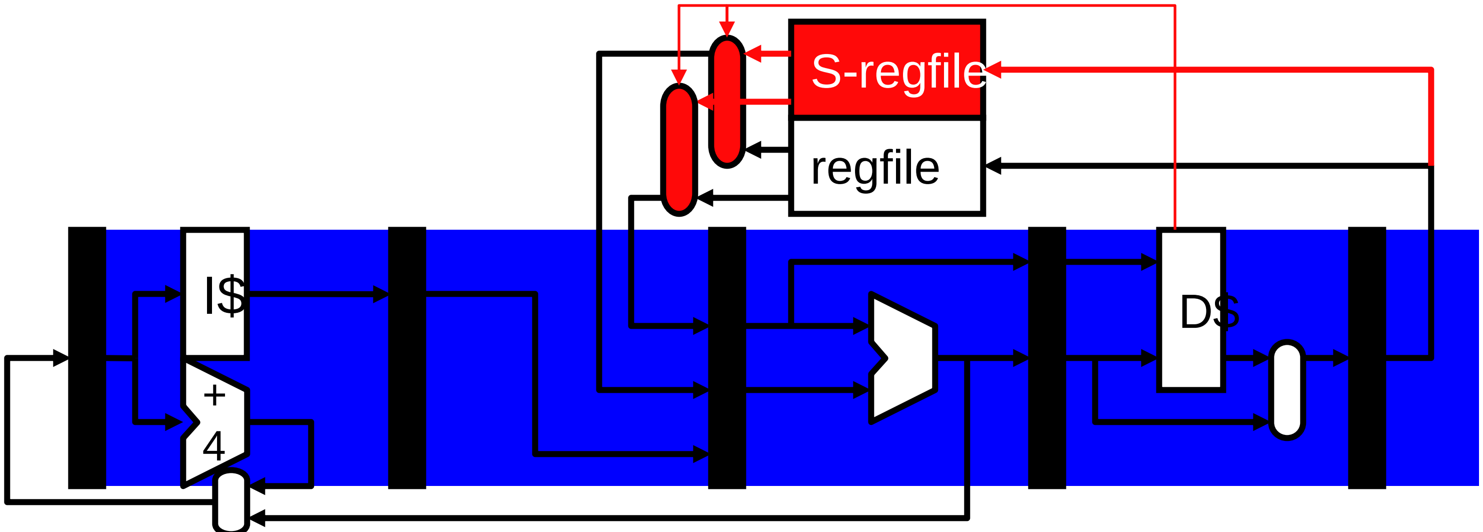


Research: Speculation Gating

- **Speculation gating** [Manne+]
  - Extend branch predictor to give prediction + **confidence**
  - Speculate on high-confidence (mis-prediction unlikely) branches
  - Stall (save energy) on low-confidence branches

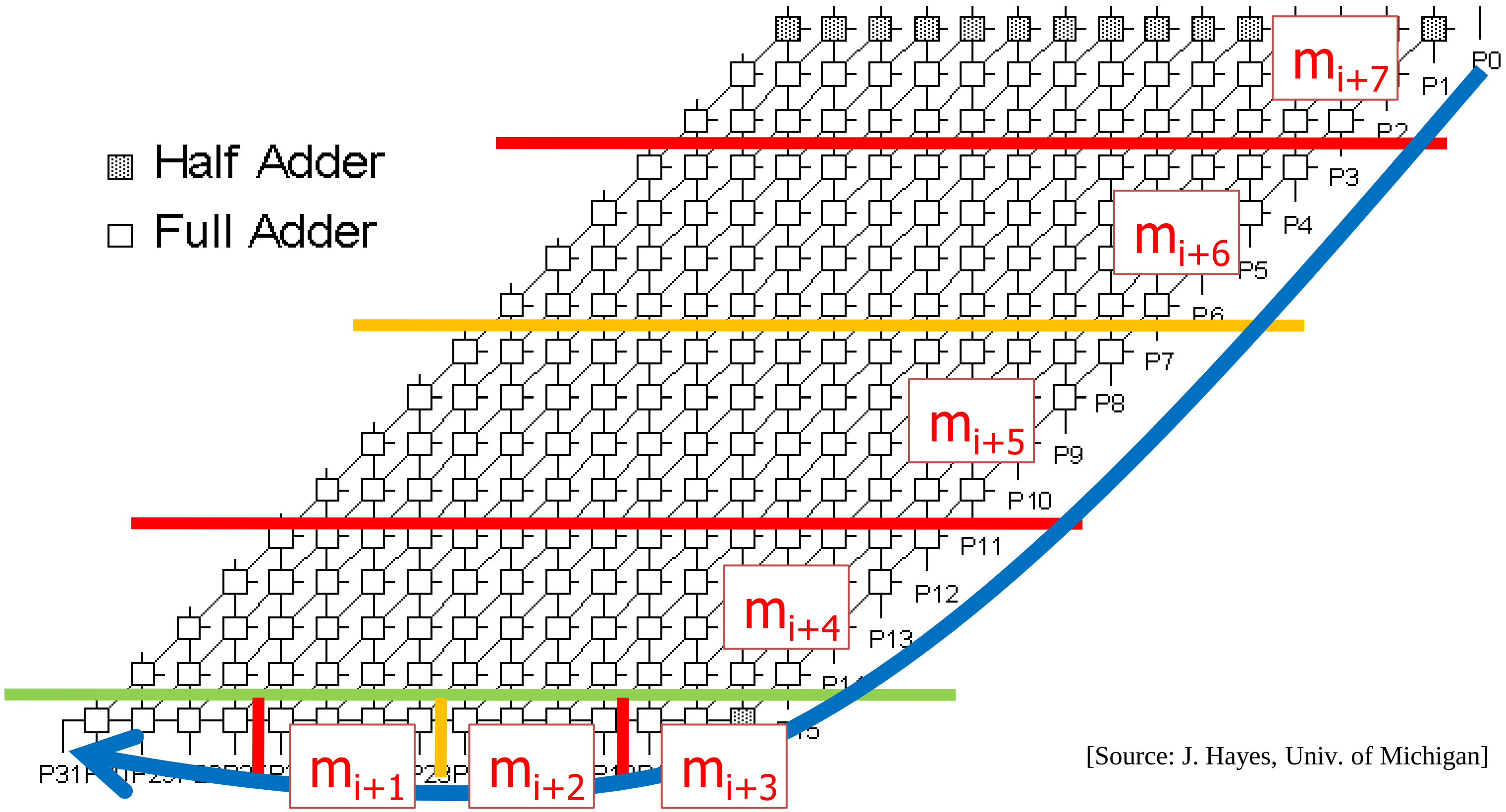
- **Confidence estimation**
  - What kind of hardware circuit estimates confidence?
  - Hard in absolute sense, but easy relative to given threshold
  - Counter-scheme similar to $\%_{miss}$ threshold for cache resizing
  - Example: assume 90% accuracy is high confidence
    - PC-indexed table of confidence-estimation counters
    - Correct prediction? table[PC]+=1 : table[PC]–=9;
    - Prediction for PC is confident if table[PC] > 0;

Research: Runahead Execution



- In-order writebacks essentially imply stalls on D$ misses
  - Can save power … or use idle time for performance
- **Runahead execution** [Dundas+ 97]
  - Shadow regfile kept in sync with main regfile (write to both)
  - D$ miss: continue executing using shadow regfile (disable stores)
  - D$ miss returns: flush pipe and restart with stalled PC
  - \+ Acts like a smart prefetch engine
  - \+ Performs better as cache t_{miss} grows (relative to clock period)

Example: Integer Multiplier



16x16 combinational multiplier

[Source: J. Hayes, Univ. of Michigan]

Dependences and Loops

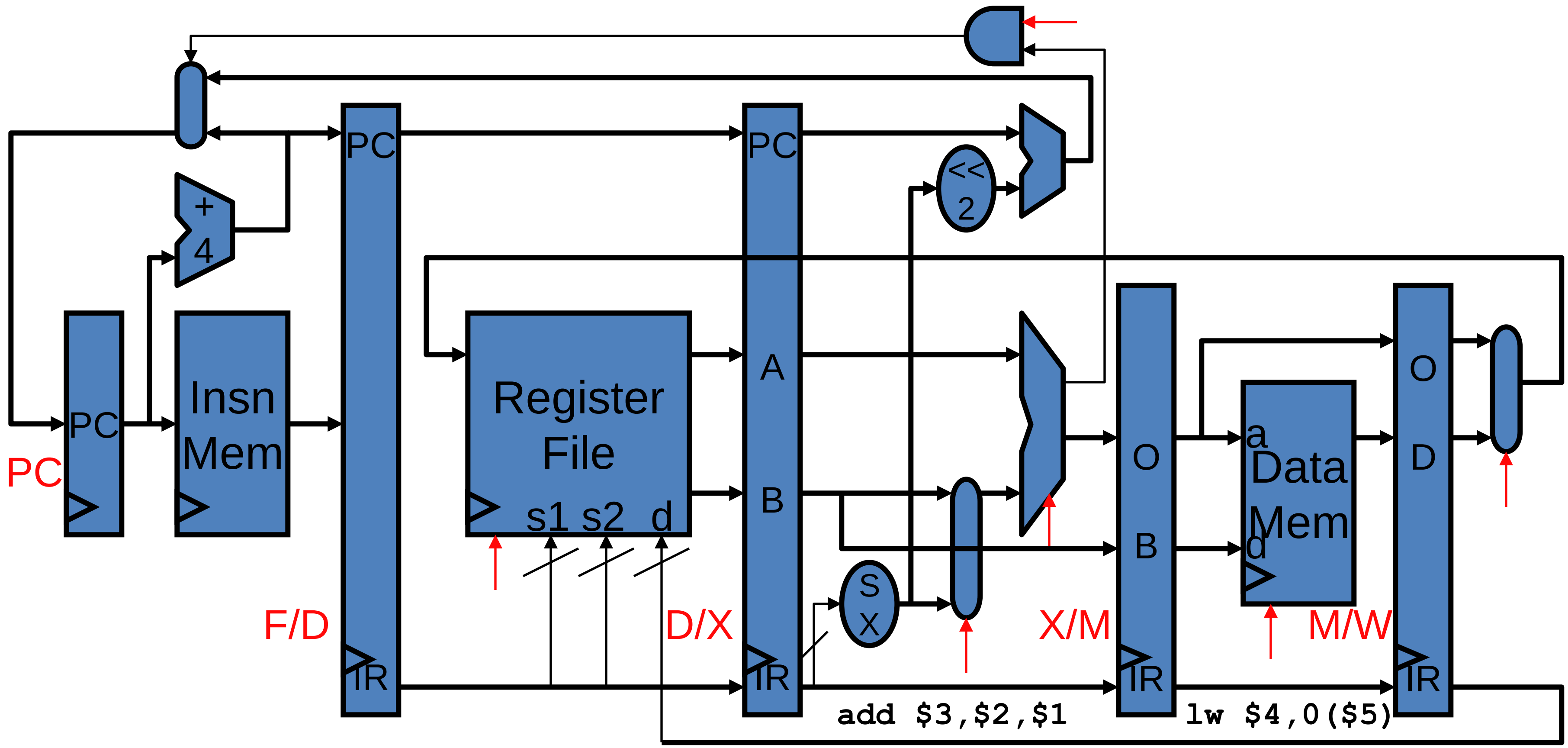
- Data dependences in loops
  - **Intra-loop**: within same iteration
  - **Inter-loop**: across iterations
  - Example: DAXPY (**D**ouble precision **A X P**lus **Y**)

```
for (i=0;i<100;i++)
   Z[i]=A*X[i]+Y[i];

0: ldf f2,X(r1)
1: mulf f2,f0,f4
2: ldf f6,Y(r1)
3: addf f4,f6,f8
4: stf f8,Z(r1)
5: addi r1,8,r1
6: cmplti r1,800,r2
7: beq r2,Loop
```

- RAW intra: 0→1(`f2`), 1→3(`f4`), 2→3(`f6`), 3→4(`f8`), 5→6(`r1`), 6→7(`r2`)
- RAW inter: 5→0(`r1`), 5→2(`r1`), 5→4(`r1`), 5→5(`r1`)
- WAR intra: 0→5(`r1`), 2→5(`r1`), 4→5(`r1`)
- WAR inter: 1→0(`f2`), 3→1(`f4`), 3→2(`f6`), 4→3(`f8`), 6→5(`r1`), 7→6(`r2`)
- WAW intra: none
- WAW inter: 0→0(`f2`), 1→1(`f4`), 2→2(`f6`), 3→3(`f8`), 6→6(`r2`)

Why Does Every Insn Take 5 Cycles?



- Could/should we allow **`add`** to skip M and go to W? No
  - It wouldn't help: peak fetch still only 1 insn per cycle
  - **Structural hazards**: imagine **`add`** follows **`lw`**

Simple Analytical Pipeline Model

- Let: insn execution require **N** stages, each takes t_n time
- Single-cycle execution
  - $\mathbf{L_1}$ (1-insn latency) = Σt_n
  - **T** (throughput) = $1/L_1$
  - $\mathbf{L_M}$ (M-insn latency, where M>>1) = $M*L_1$

- Now: N-stage pipeline
  - $L_{1+P} = L_1$
  - $T_{+P} = 1/\max(t_n) \leq$ **N**$/L_1$
    - If t_n are equal (i.e., $\max(t_n) = L_1/N$), throughput = N/L_1
  - $L_{M+P} = M*\max(t_n) \geq M*L_1/$**N**
  - S_{+P} (speedup) = $[M*L_1 / (\geq M*L_1/N)]$ = **≤ N**

- Q: for arbitrarily high speedup, use arbitrarily high N?

N-stages != ∞ due to Pipeline Overhead

- Let: **O** be extra delay per pipeline stage
  - Latch overhead: pipeline latches take time
  - Clock/data skew

- Now: N-stage pipeline with overhead
  - Assume $\max(t_n) = L_1/N$
  - $L_{1+P+O} = L_1 + N^*O$
  - $T_{+P+O} = 1/(L_1/N + O) = 1/(1/T + O) \leq T$, **$\leq T/O$**
  - $L_{M+P+O} = M^*L_1/N + M^*O = L_{M+P} + M^*O$
  - $S_{+P+O} = [M^*L_1 / (M^*L_1/N + M^*O)] = \leq N = S_{+P}$, **$\leq L_1/O$**

- **O** limits throughput and speedup → useful N

N-stages != due to Hazards

- **Dependence**: relationship that serializes two insns
  - **Data**: two insns use the same value or storage location
  - **Control**: one instruction affects whether another executes at all
  - **Maybe**: two insns *may* have a dependence
- **Hazard**: dependence causes potential incorrect execution
  - Possibility of using or corrupting data or execution flow
  - **Structural**: two insns want to use same structure, one must wait
  - Often fixed with **stalls**: insn stays in same stage for multiple cycles
- Let: **H** be average number of hazard stall cycles per instruction
  - $L_{1+P+H} = L_{1+P}$ (no hazards for one instruction)
  - $T_{+P+H} = [N/(N+H)]*N/L_1 = [N/(N+H)] * T_{+P}$
  - $L_{M+P+H} = M* L_1/N * [(N+H)/N] = [(N+H)/N] * L_{M+P}$
  - $S_{+P+H} = M*L_1 / M*L_1/N*[(N+H)/N] = [N/(N+H)]*S_{+P}$
- **H** also limit throughput, speedup → useful N
  - N↑→ H↑ (more insns "in flight" → more dependences become hazards)
  - Exact H depends on program, requires detailed simulation/model

Compiler Scheduling

- Compiler can schedule (move) insns to reduce stalls
  - **Basic pipeline scheduling**: eliminate back-to-back load-use pairs
  - Example code sequence: `a = b + c; d = f – e;`
  - MIPS Notation:
    - “ld r2,4(sp)” is “ld [sp+4]➔r2” “st r1, 0(sp)” is “st r1➔[sp+0]”

Before

```
ld r2,4(sp)
ld r3,8(sp)
add r3,r2,r1   //stall
st r1,0(sp)
ld r5,16(sp)
ld r6,20(sp)
sub r5,r6,r4   //stall
st r4,12(sp)
```

After

```
ld r2,4(sp)
ld r3,8(sp)
ld r5,16(sp)
add r3,r2,r1   //no stall
ld r6,20(sp)
st r1,0(sp)
sub r5,r6,r4   //no stall
st r4,12(sp)
```

Compiler Scheduling Requires

- **Large scheduling scope**
  - Independent instruction to put between load-use pairs
  - + Original example: large scope, two independent computations
  - – This example: small scope, one computation

Before

```
ld r2,4(sp)
ld r3,8(sp)
add r3,r2,r1   //stall
st r1,0(sp)
```

After

```
ld r2,4(sp)
ld r3,8(sp)
add r3,r2,r1   //stall
st r1,0(sp)
```

Compiler Scheduling Requires

- **Enough registers**
  - To hold additional "live" values
  - Example code contains 7 different values (including **sp**)
  - Before: max 3 values live at any time → 3 registers enough
  - After: max 4 values live → 3 registers not enough → WAR violations

Original

```
ld r2,4(sp)
ld r1,8(sp)
add r1,r2,r1   //stall
st r1,0(sp)
ld r2,16(sp)
ld r1,20(sp)
sub r2,r1,r1   //stall
st r1,12(sp)
```

Wrong!

```
ld r2,4(sp)
ld r1,8(sp)
ld r2,16(sp)
add r1,r2,r1   //WAR
ld r1,20(sp)
st r1,0(sp)    //WAR
sub r2,r1,r1
st r1,12(sp)
```

Compiler Scheduling Requires

- **Alias analysis**
  - Ability to tell whether load/store reference same memory locations
    - Effectively, whether load/store can be rearranged
  - Example code: easy, all loads/stores use same base register (**`sp`**)
  - New example: can compiler tell that **`r8`** = **`sp`**?

Before

```
ld r2,4(sp)
ld r3,8(sp)
add r3,r2,r1   //stall
st r1,0(sp)
ld r5,0(r8)
ld r6,4(r8)
sub r5,r6,r4   //stall
st r4,8(r8)
```

Wrong(?)

```
ld r2,4(sp)
ld r3,8(sp)
ld r5,0(r8)
add r3,r2,r1
ld r6,4(r8)
st r1,0(sp)
sub r5,r6,r4
st r4,8(r8)
```

- Reverse stream analogy
  - "Downstream": earlier stages, younger insns
  - "Upstream": later stages, older insns
  - Reverse? instruction stream fixed, pipeline flows over it
    - Architects see instruction stream as fixed by program/compiler

Two Stall Timings (without bypassing)

- Depend on how D and W stages share regfile
  - Each gets regfile for half a cycle
  – 1st half D reads, 2nd half W writes 3 cycle stall
  - **d*** = data stall, **p*** = propagated stall

	1	2	3	4	5	6	7	8	9	10
`add r2,r3→r1`	F	D	X	M	W					
`sub r1,r4→r2`		F	**d***	**d***	**d***	D	X	M	W	
`add r5,r6→r7`			**p***	**p***	**p***	F	D	X	M	W

\+ 1st half W writes, 2nd half D reads 2 cycle stall

- How does the stall logic change here?

	1	2	3	4	5	6	7	8	9	10
`add r2,r3→r1`	F	D	X	M	W					
`sub r1,r4→r2`		F	**d***	**d***	D	X	M	W		
`add r5,r6→r7`			**p***	**p***	F	D	X	M	W	