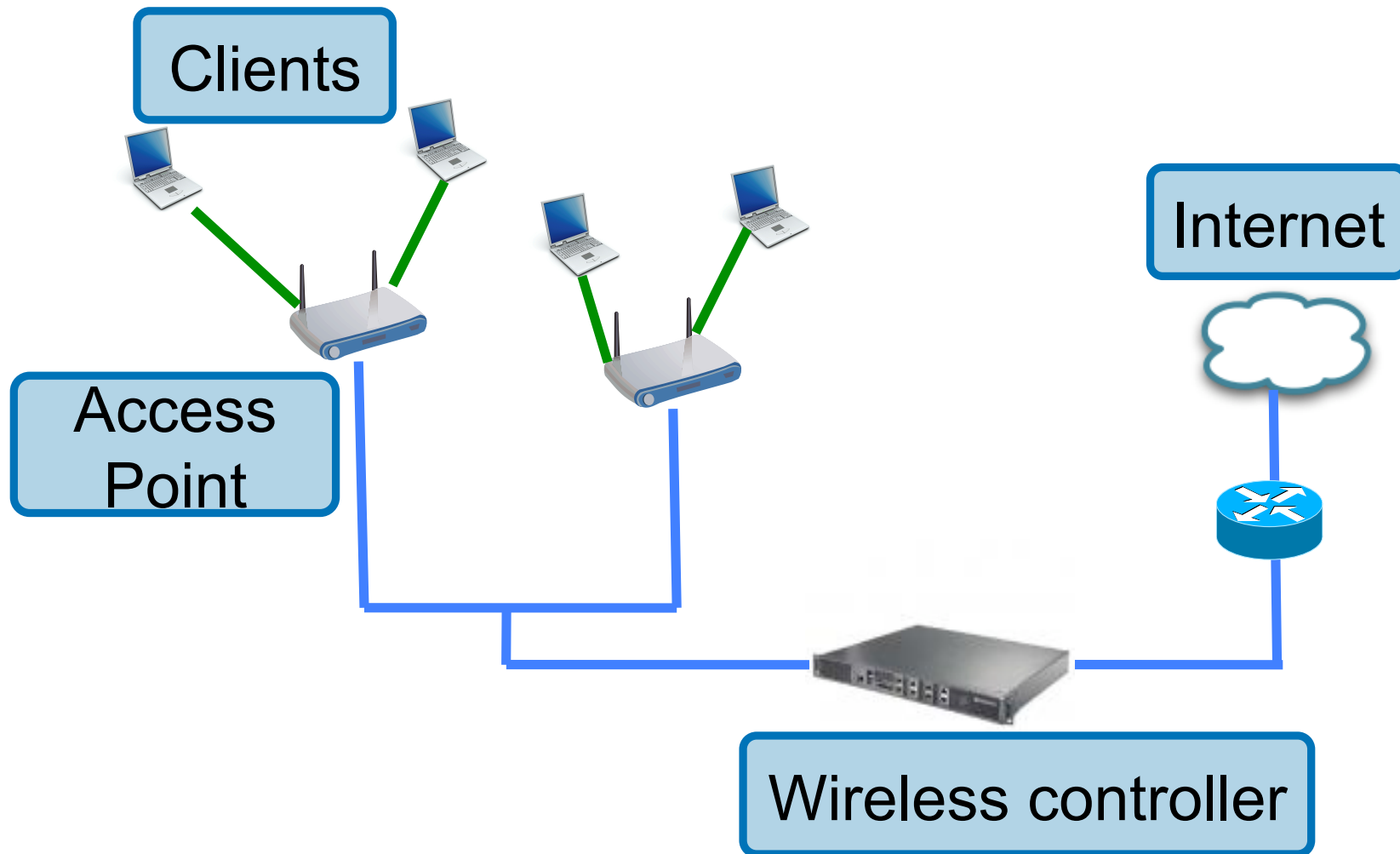


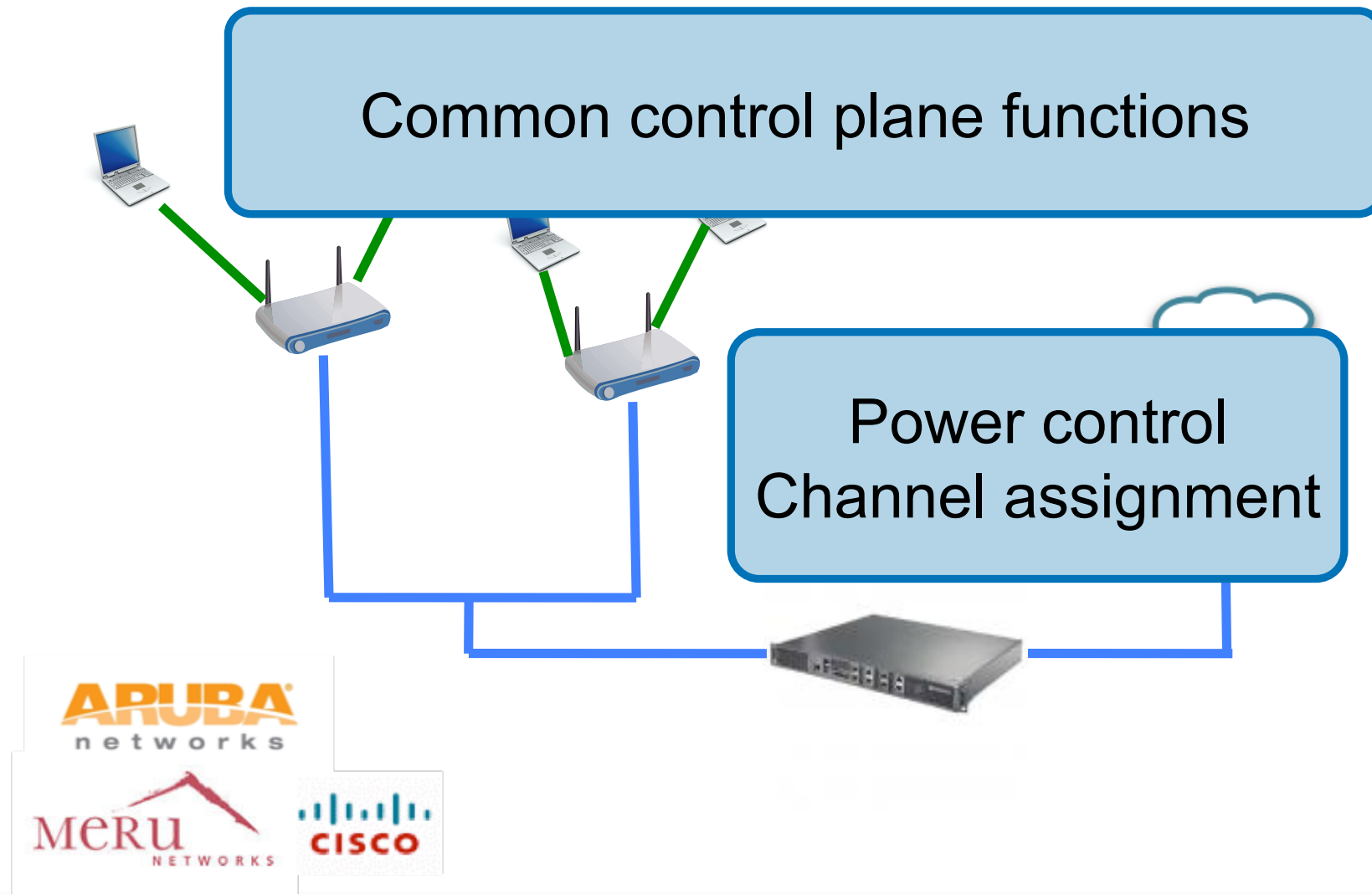
CENTAUR: Realizing the Full Potential of Centralized WLANs Through a Hybrid Data Path

Suman Banerjee
(talk originally prepared by Vivek Shrivastava)

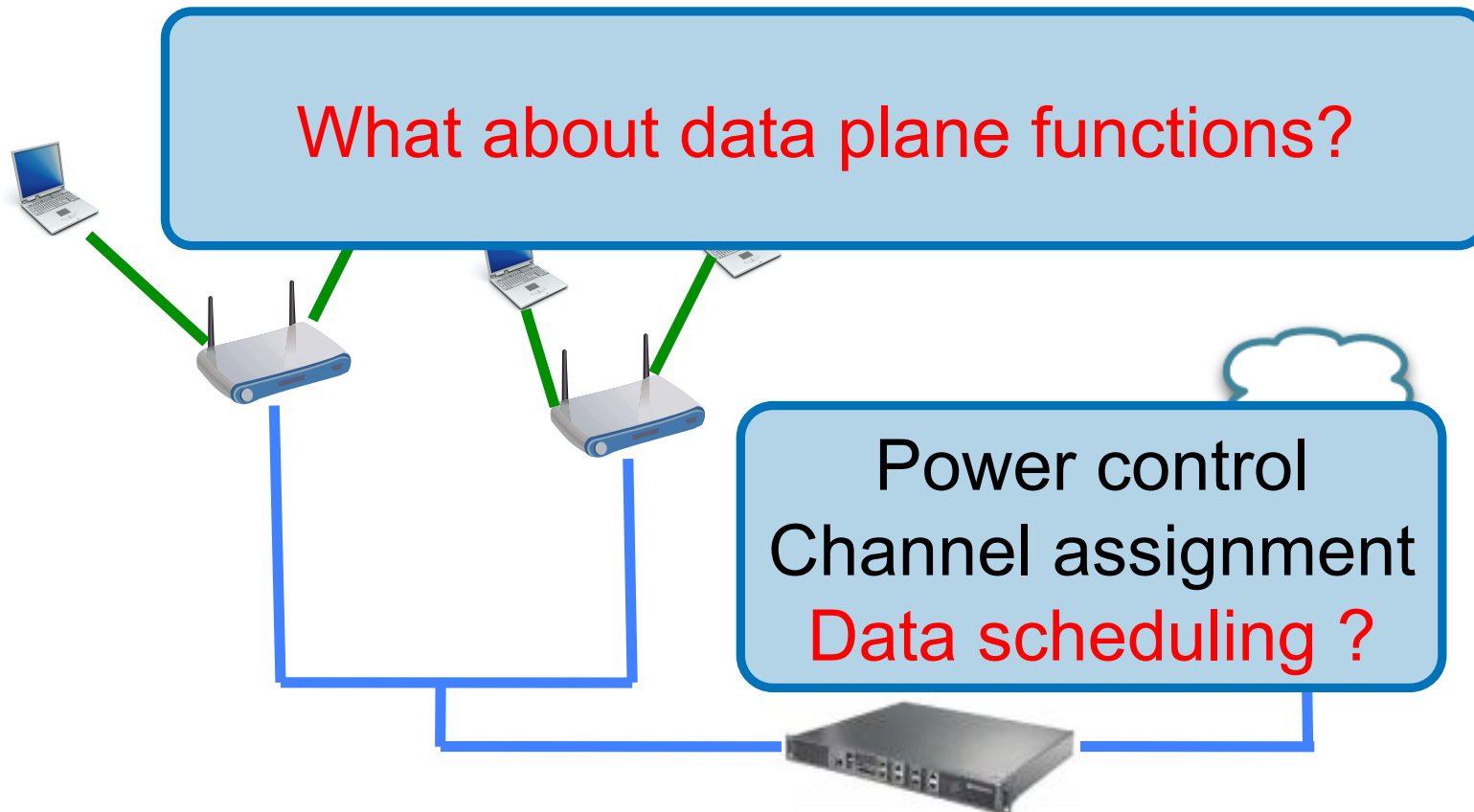
Centralization of Enterprise WLANs



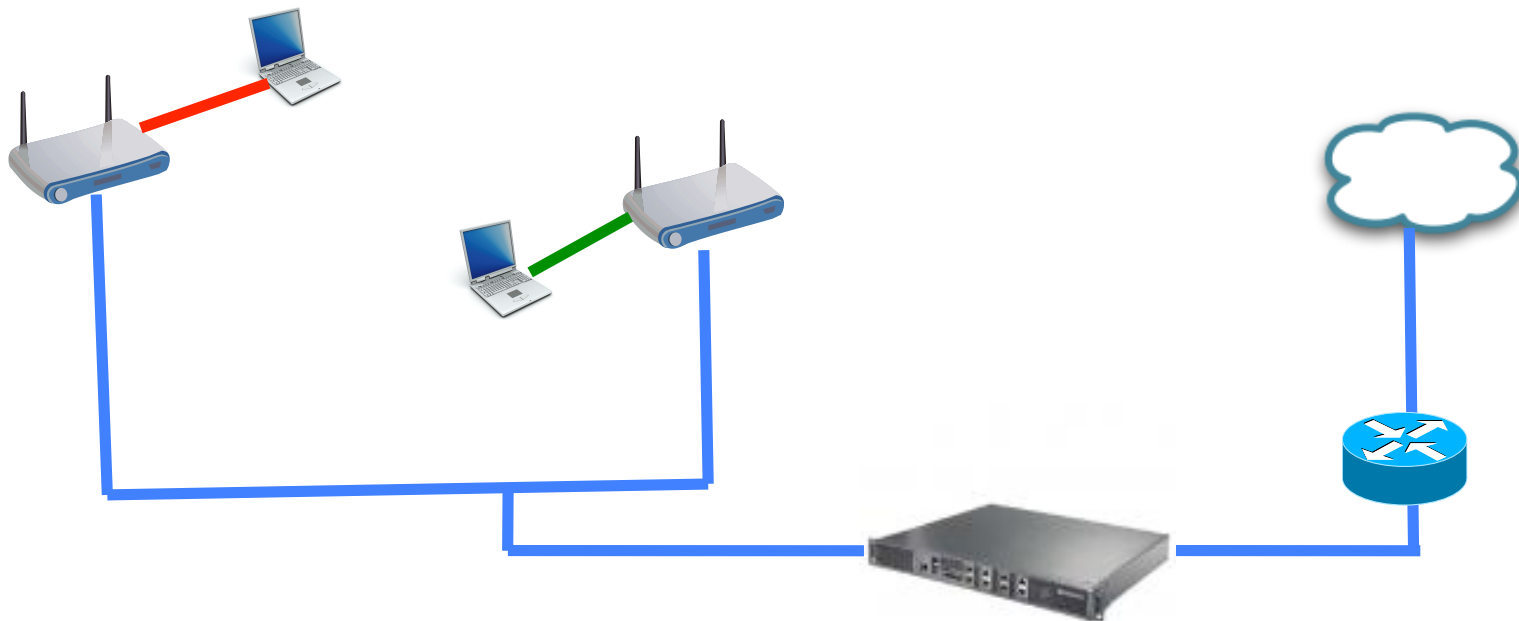
Centralization of Enterprise WLANs



Centralization of Enterprise WLANs

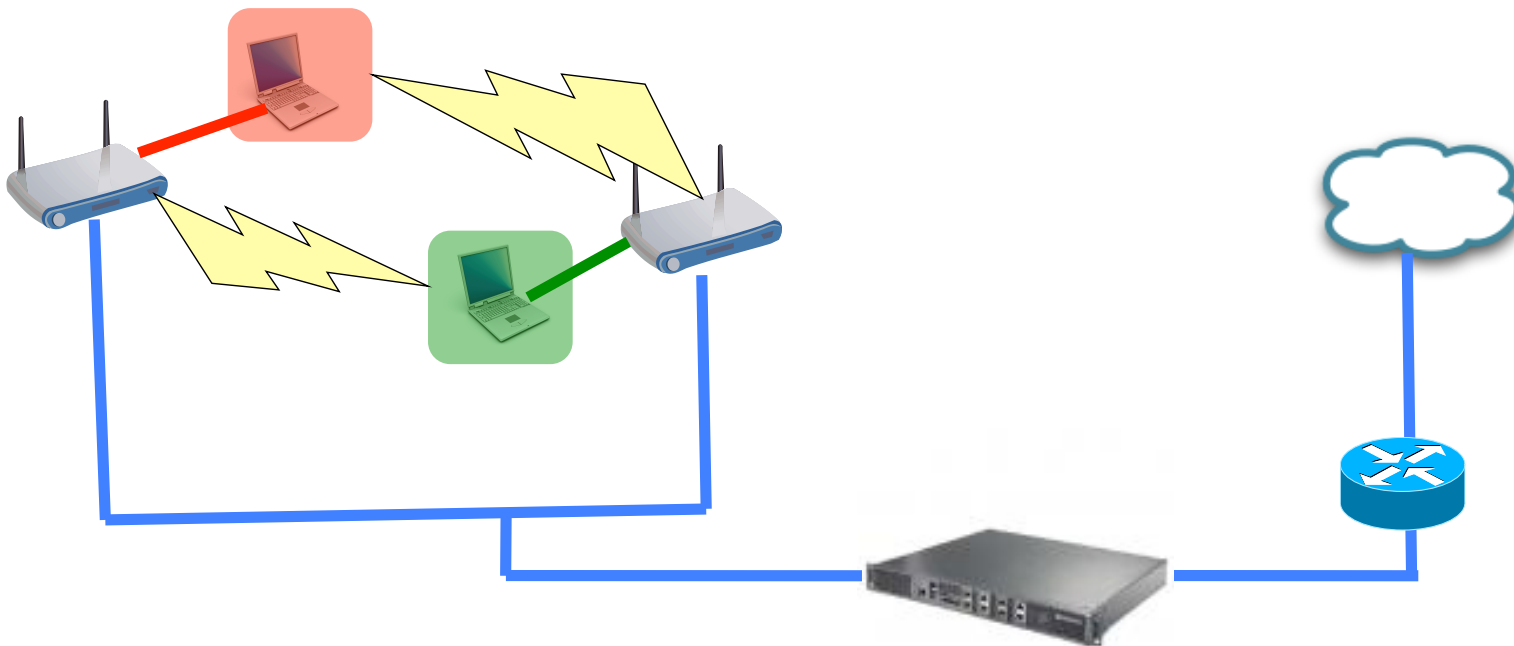


Can centralized scheduling help?



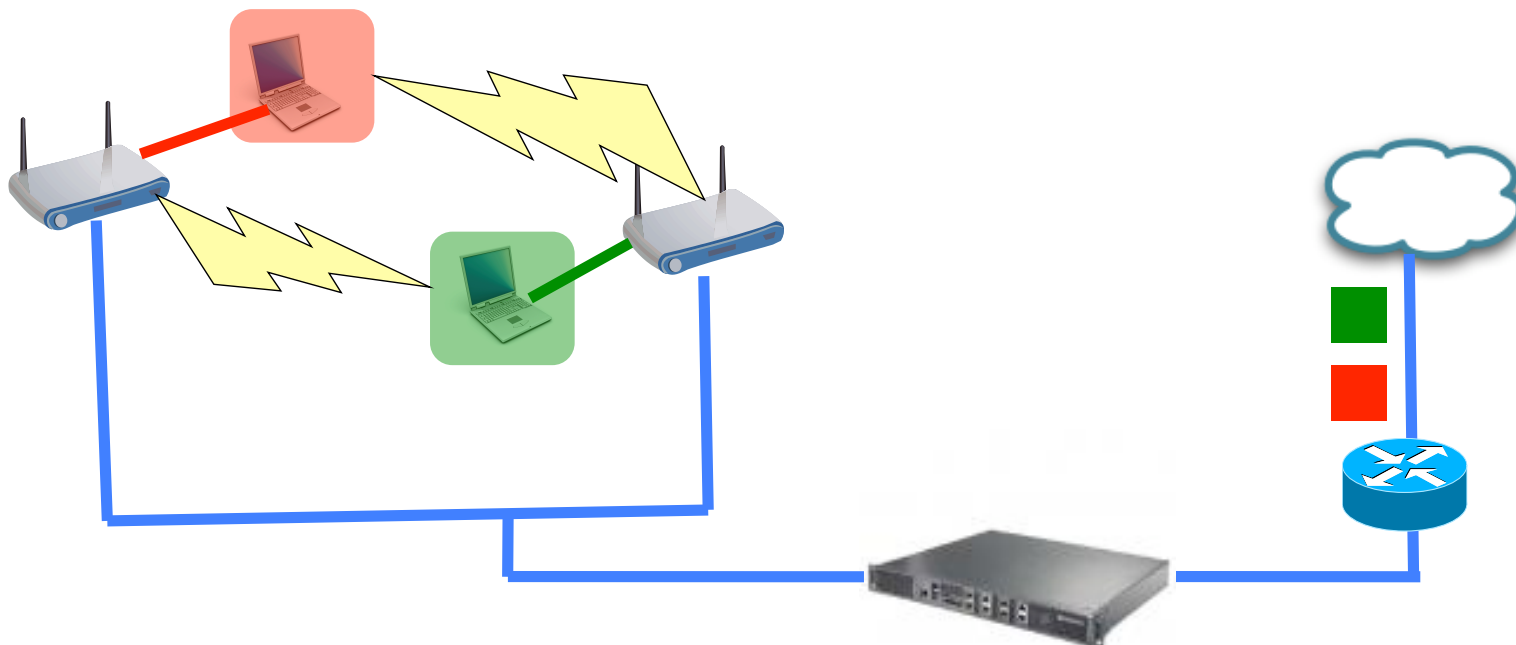
Can centralized scheduling help?

Hidden terminals



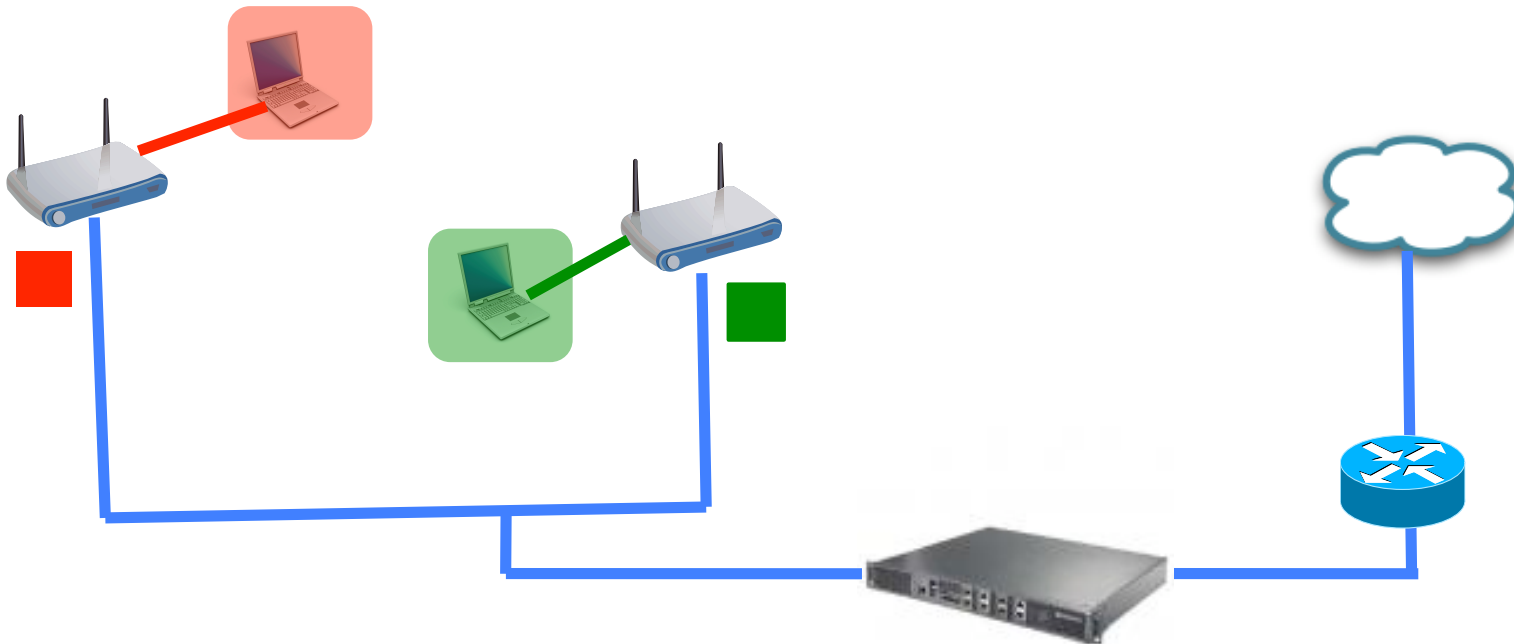
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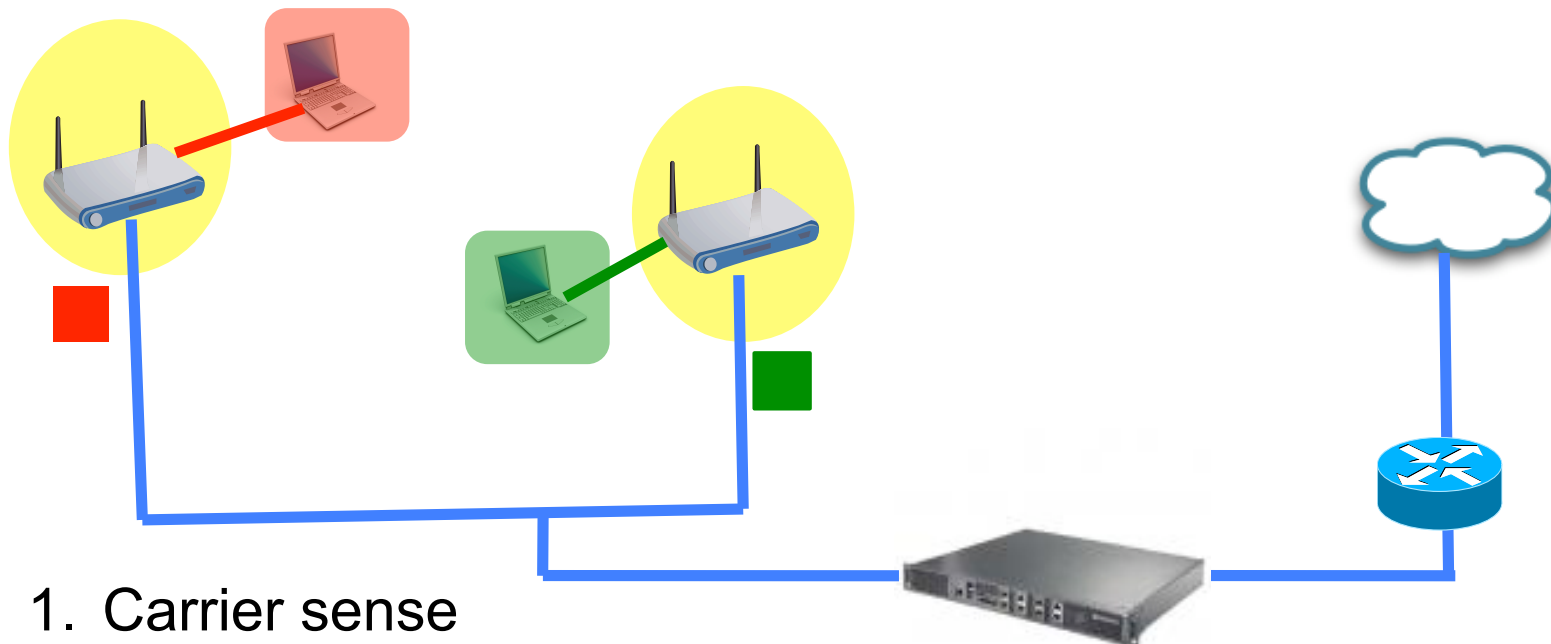
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Hidden terminals



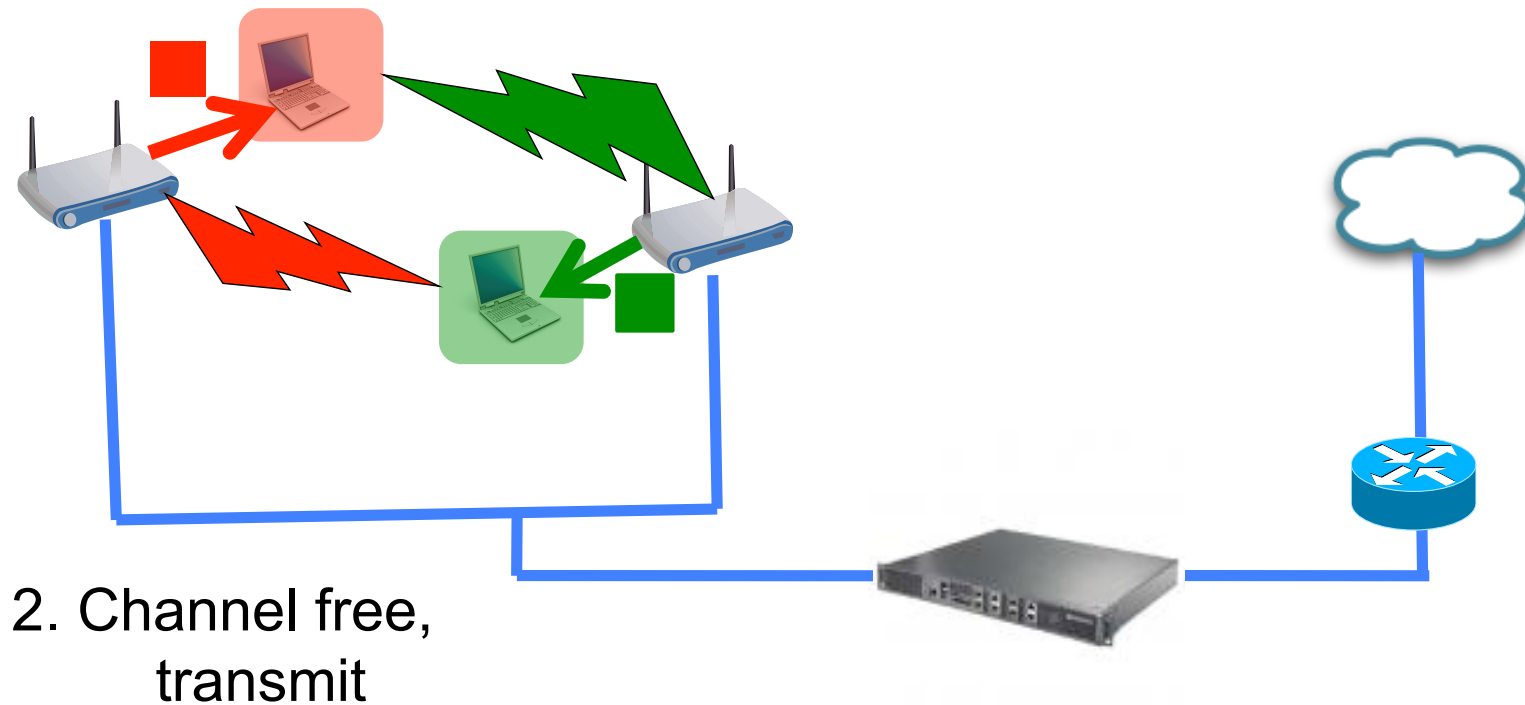
Can centralized scheduling help?

Hidden terminals

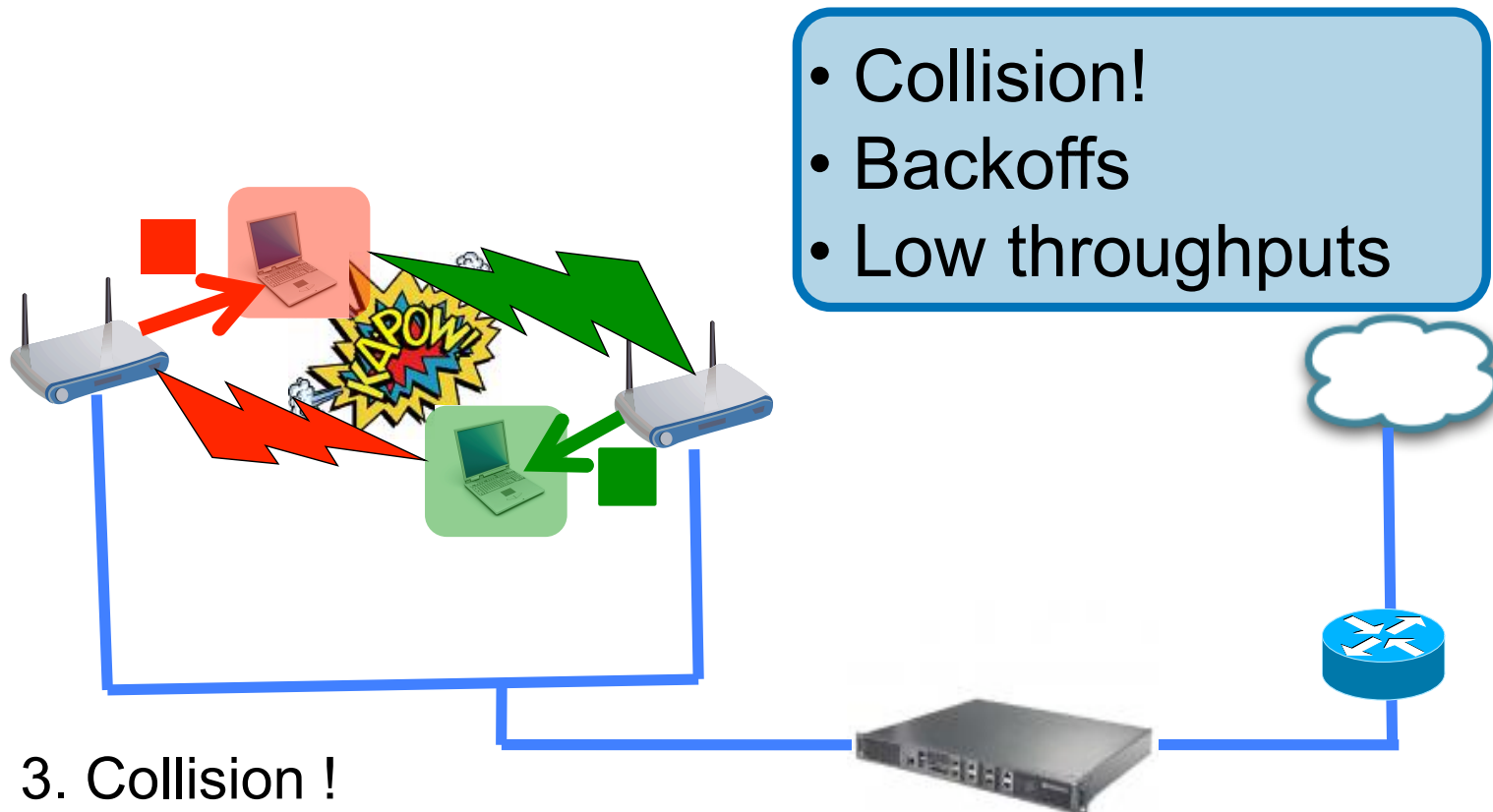


Can centralized scheduling help?

Hidden terminals



Can centralized scheduling help?

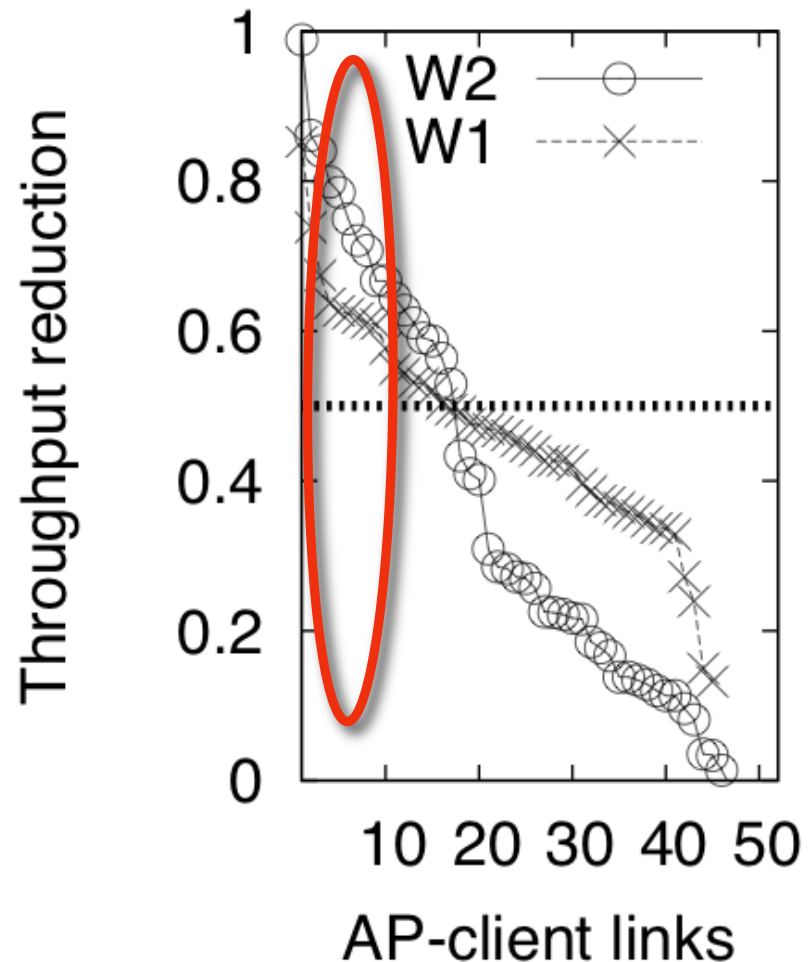


How bad is it ?

Experiments on production building-wide WLANs

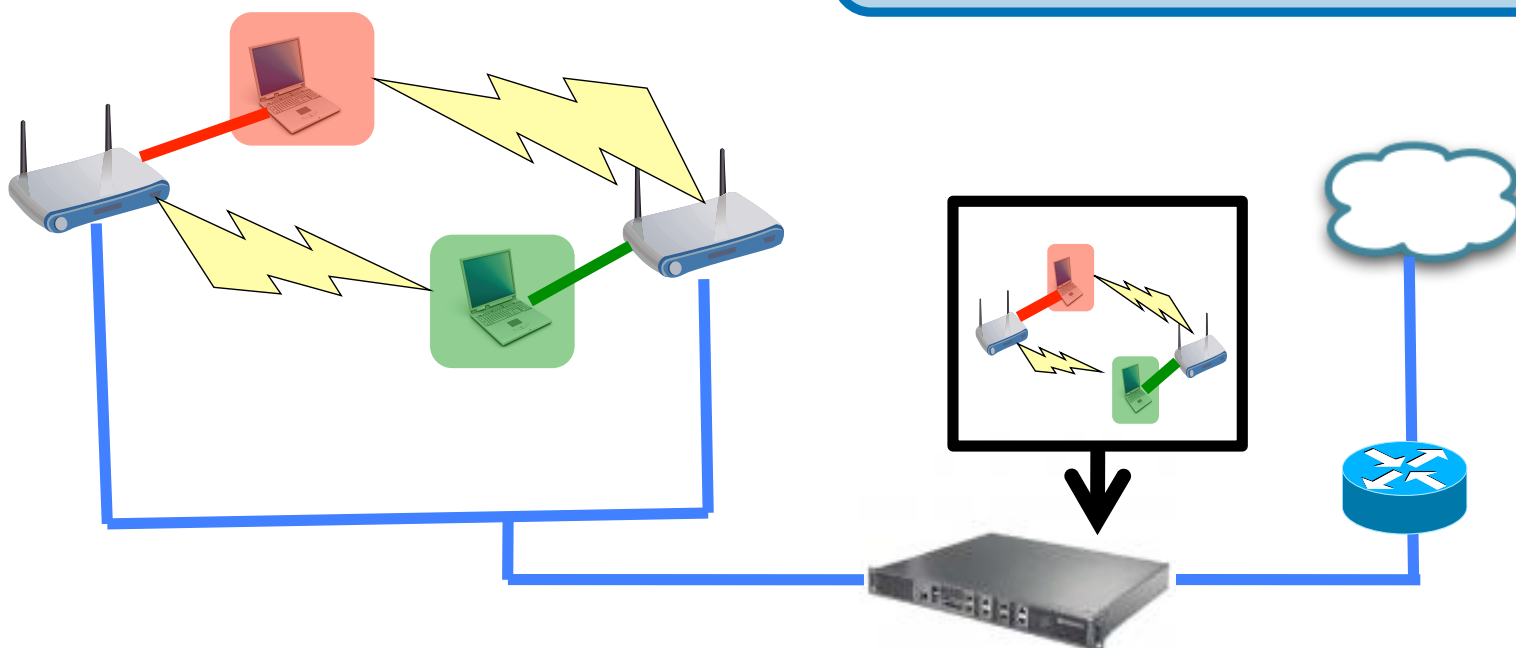
- W1: 5 floors
9 APs, 45 clients
- W2: 1 floor
21 APs, 51 clients

10% links suffer more than 70% throughput reduction



A lost opportunity ?

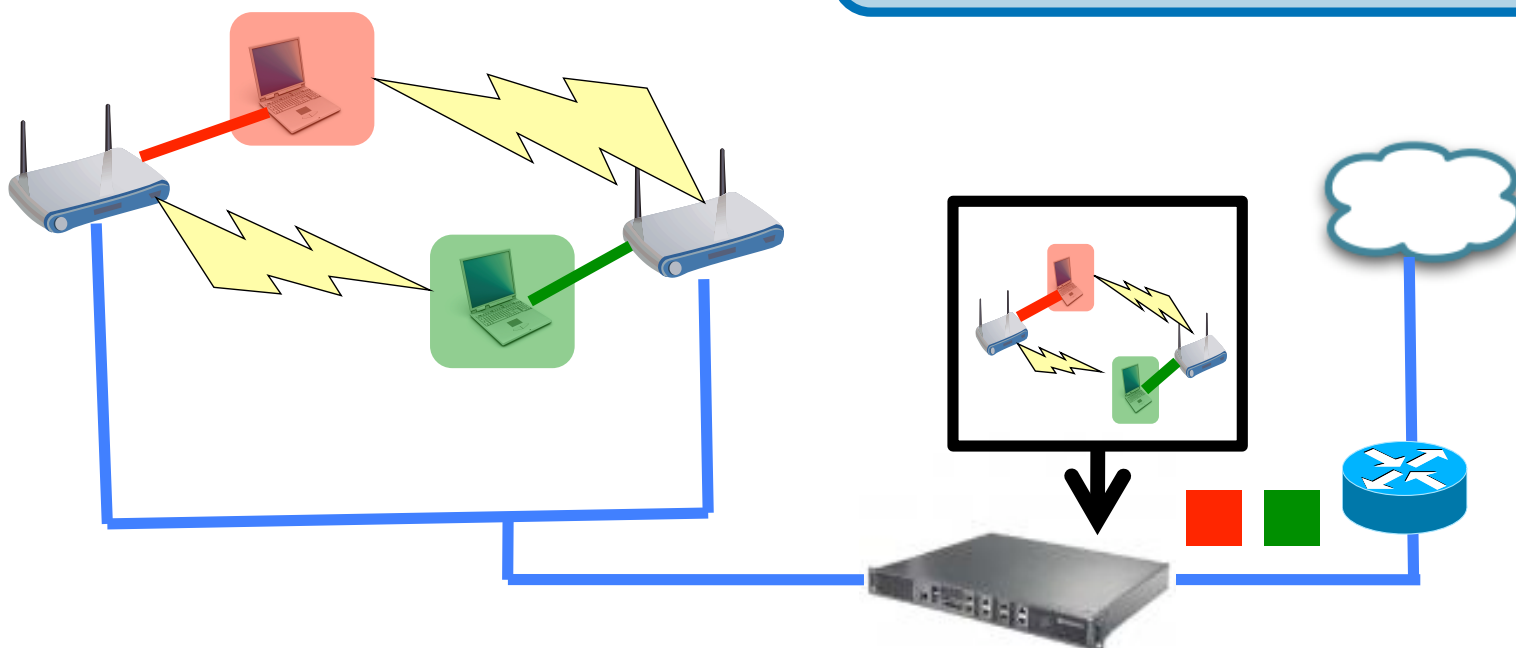
Suppose infrastructure
can gather conflict data



A lost opportunity ?

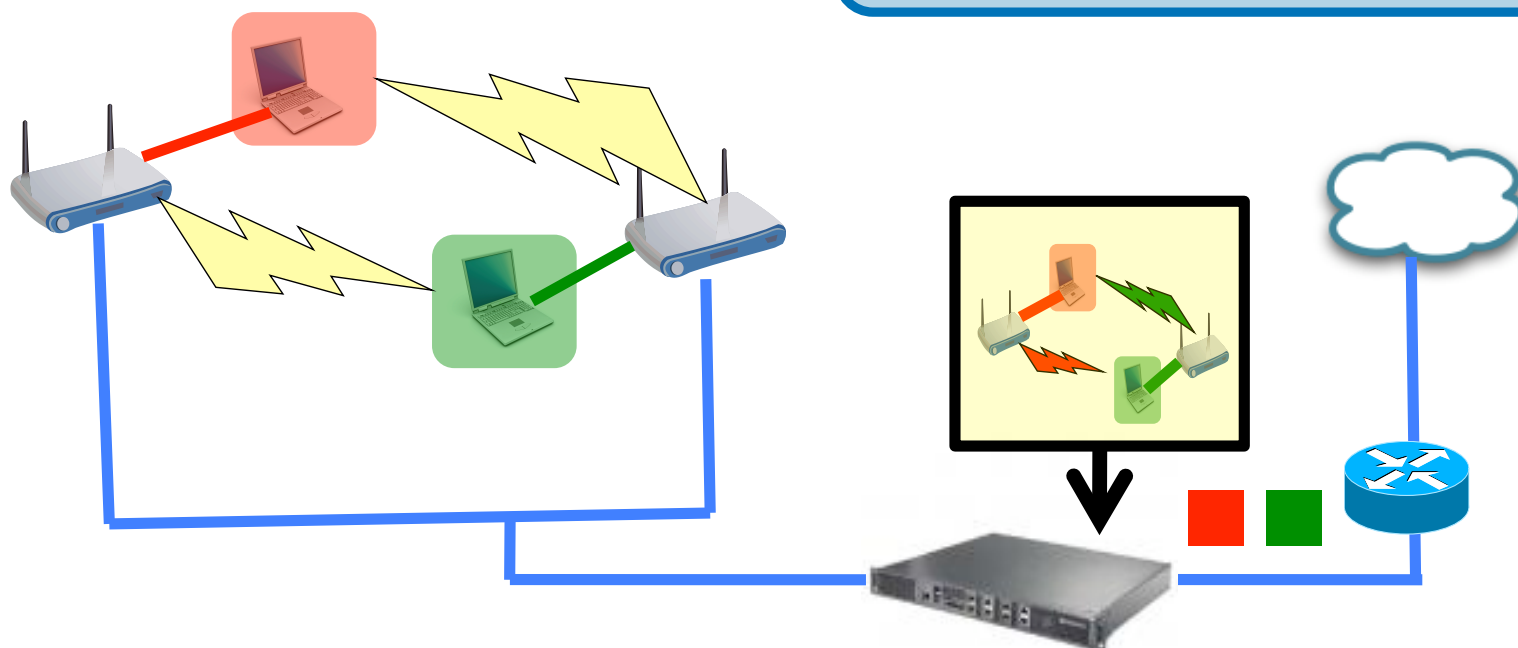
And when packets arrive

...



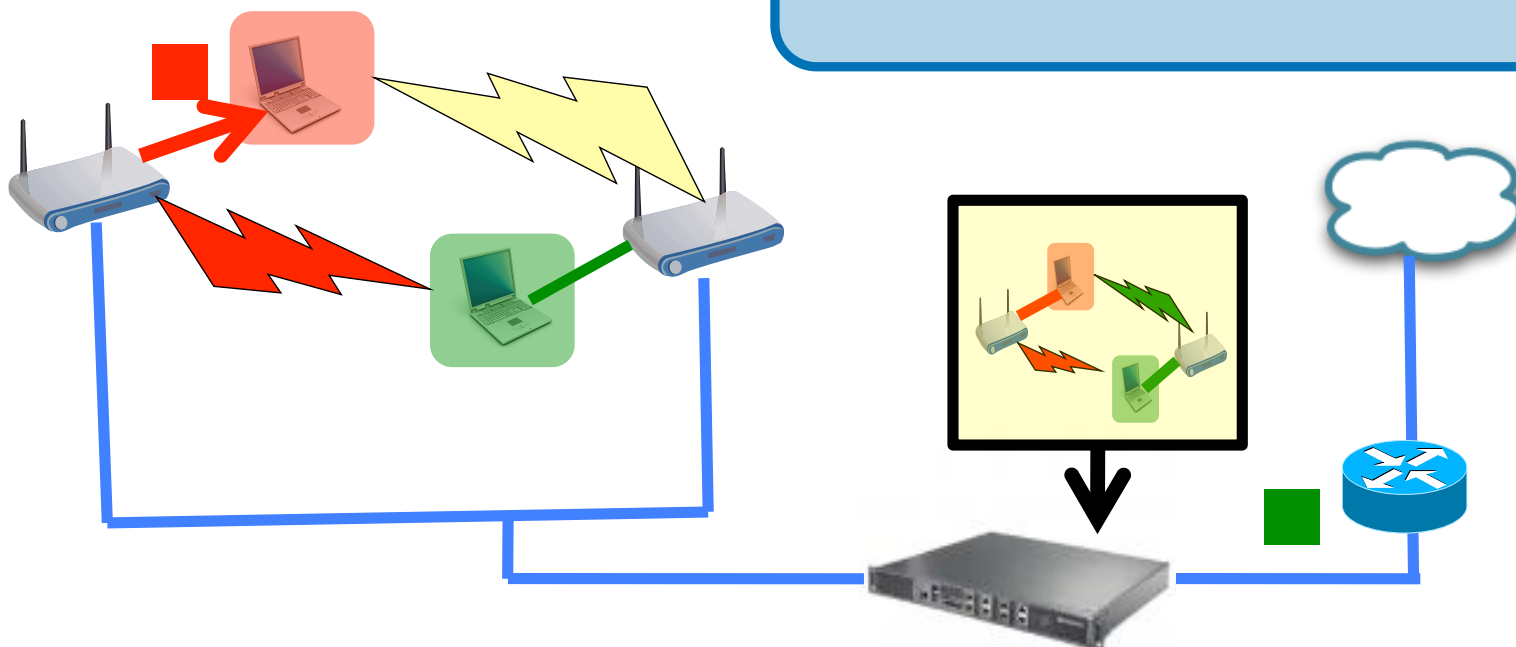
A lost opportunity ?

... realize interference
will happen ...



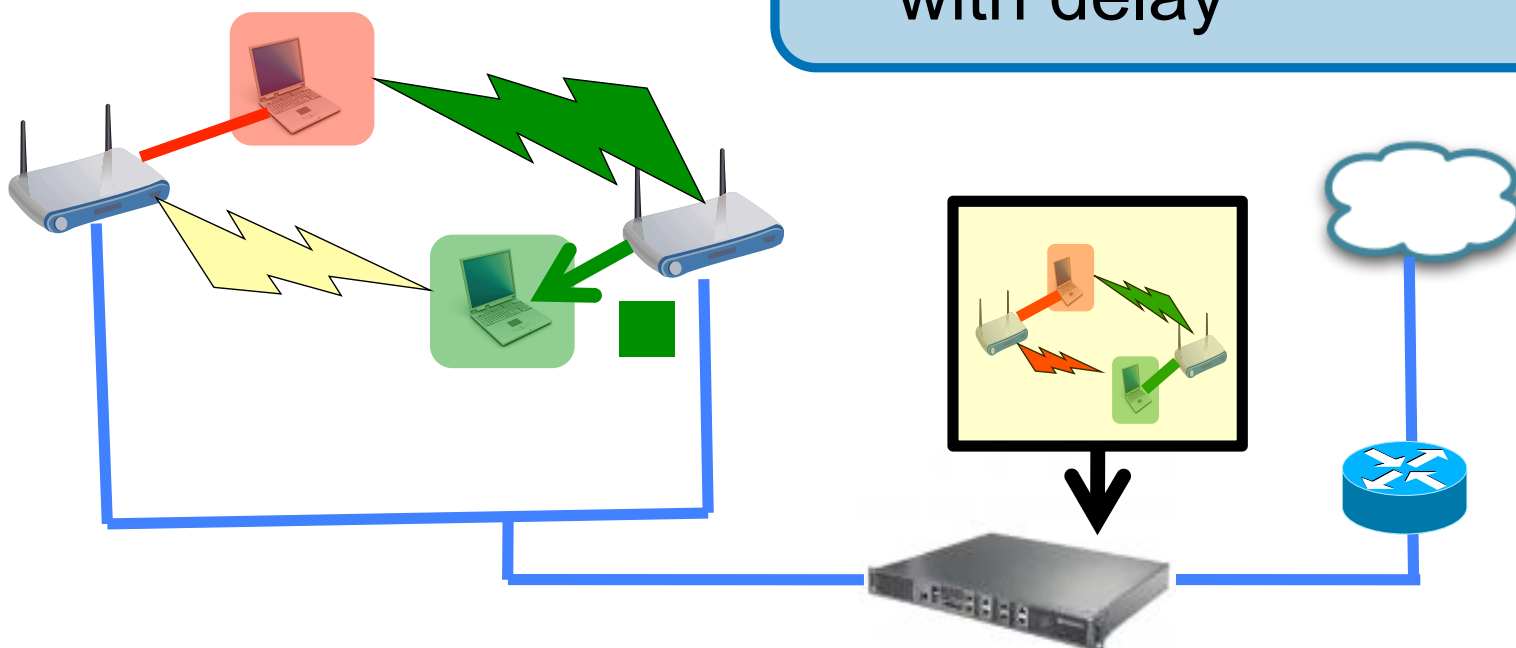
A lost opportunity ?

1. Transmit first packet



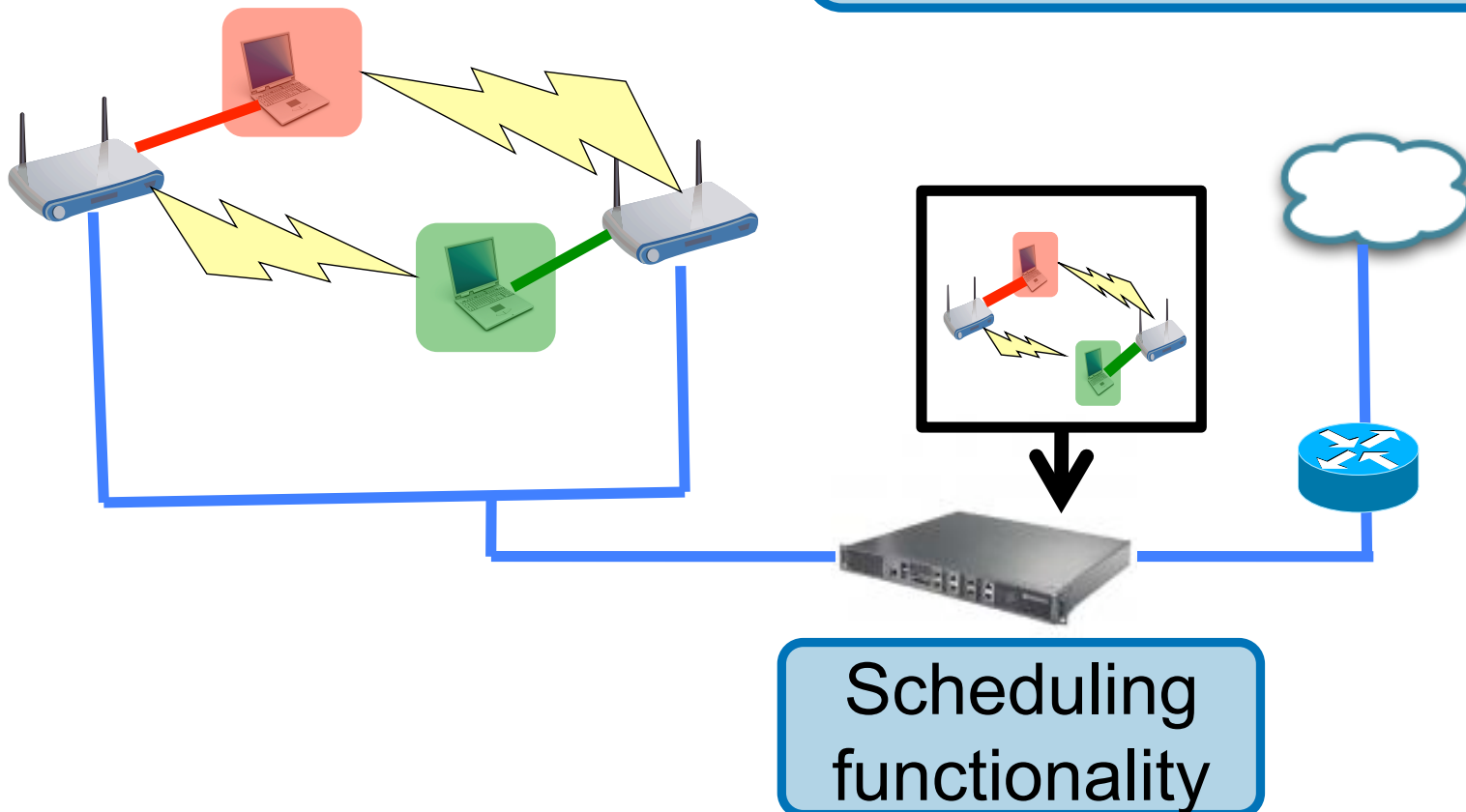
A lost opportunity ?

1. Transmit first packet
2. Transmit second packet with delay



Use an in-band scheduler

Simple FIFO schedule
with interference
avoidance



What about exposed terminals ?

- In our experiments, about 41% of link pairs suffer from exposed terminal interference
- Disabling carrier sense to solve the problem can be dangerous for uplink, non-enterprise traffic
- We will show how centralization can help even here

Challenge

- Centralization has obvious overheads
- How to make this feasible and useful under
 - Real applications
 - Common large-scale wireless environments
 - Presence of uplink and non-enterprise traffic
 - Requirements of no client modifications

Contributions

1.Design CENTAUR, a hybrid (partly centralized, partly distributed) scheduling approach

- Resolves hidden and exposed terminals
- Requires no client modifications
- Requires no carrier sense disabling

2.Evaluate CENTAUR on two WLAN testbeds with real-world traffic traces

- 1.48x greater throughput for bulk data traffic
- 1.38x reduction in web-transaction times

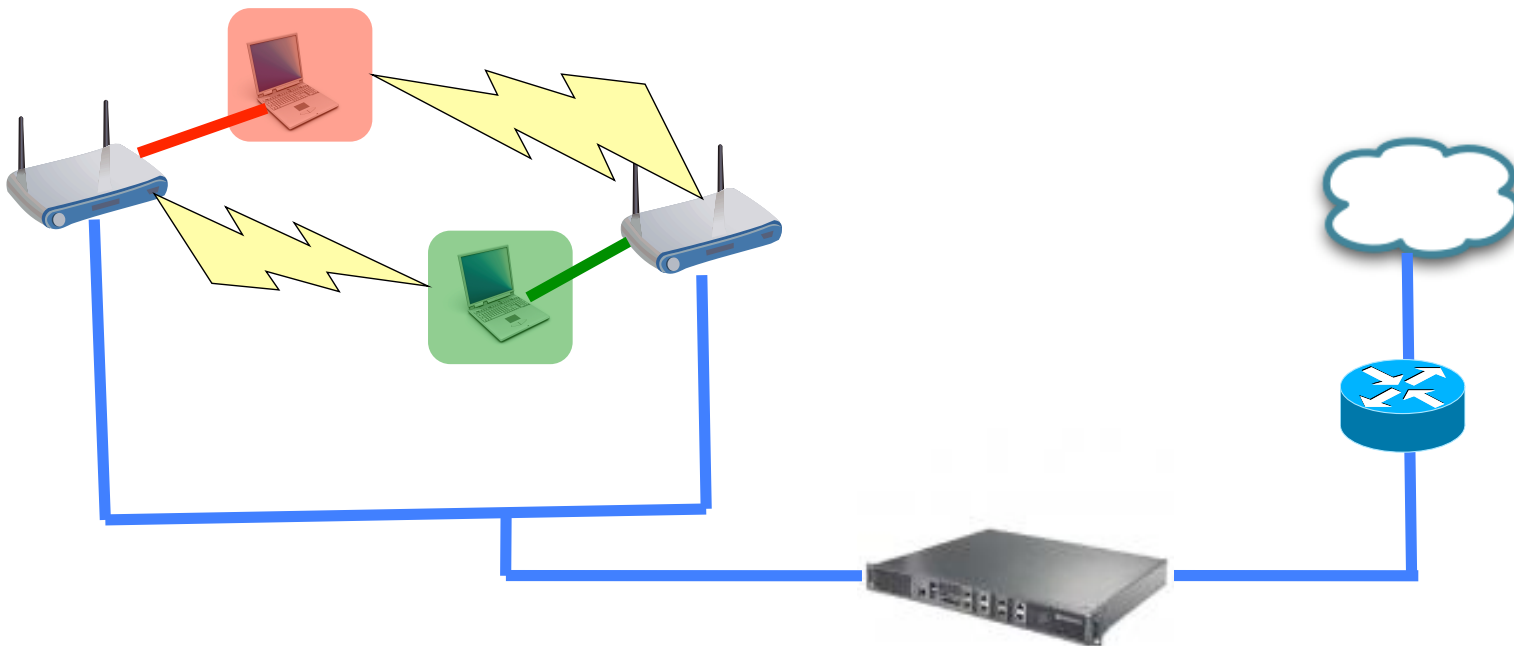
Outline

- A naïve attempt at centralized scheduling (DET)
- Our hybrid centralized scheduler (CENTAUR)
- Evaluation
- Related Work
- Summary

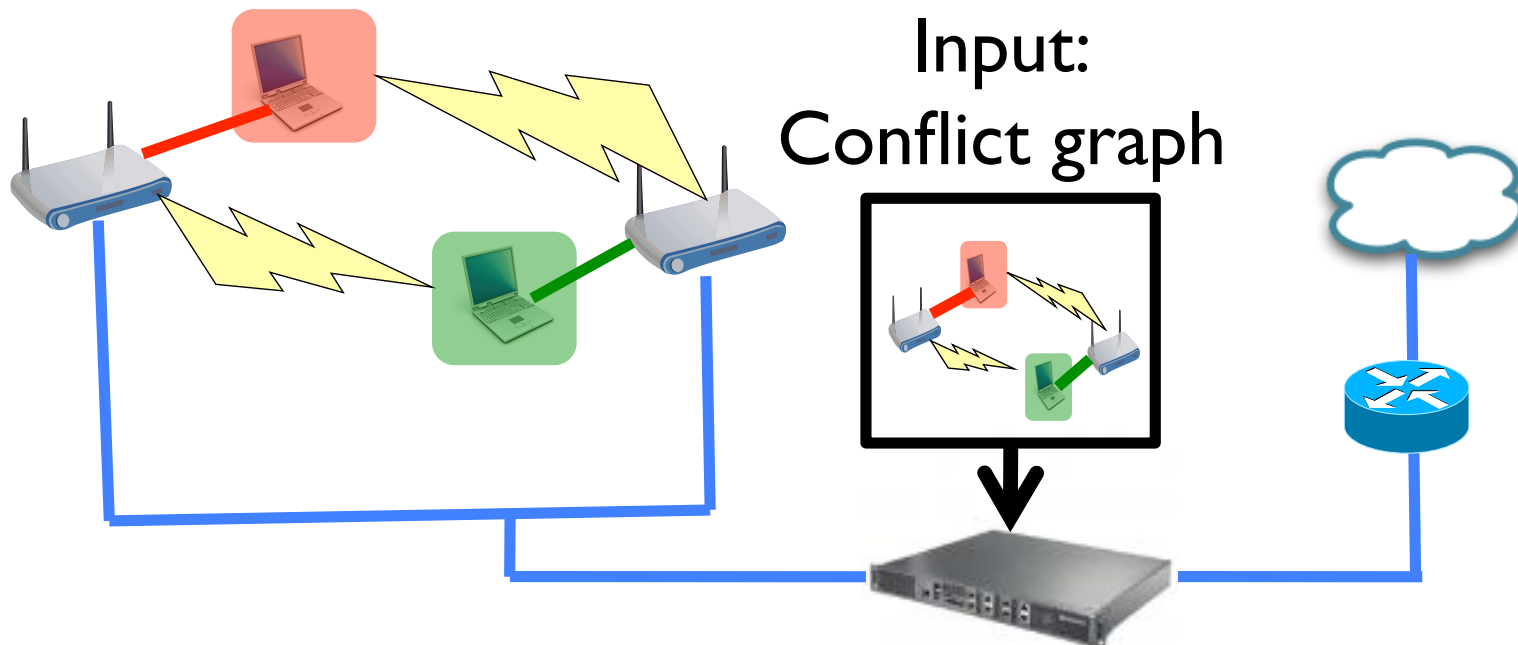
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DET: A Simple Deterministic Scheduler

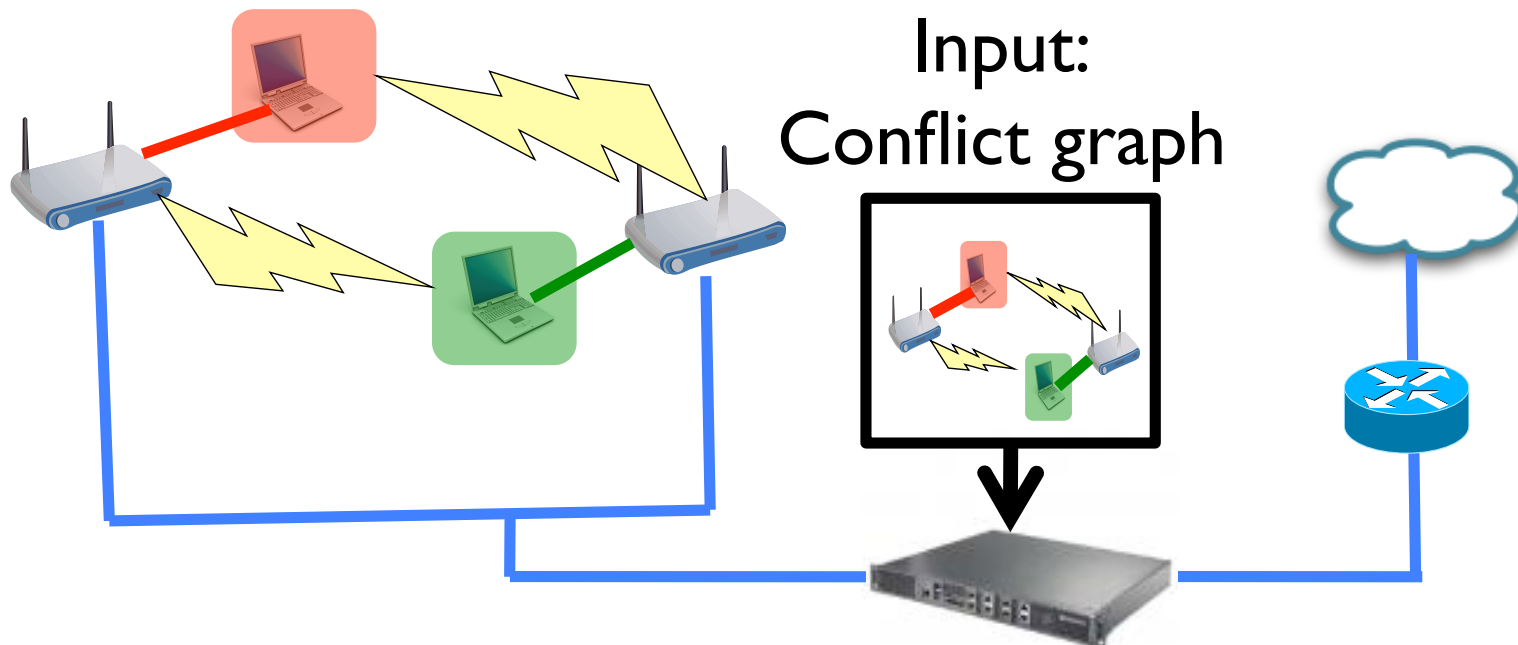


DET: A Simple Deterministic Scheduler



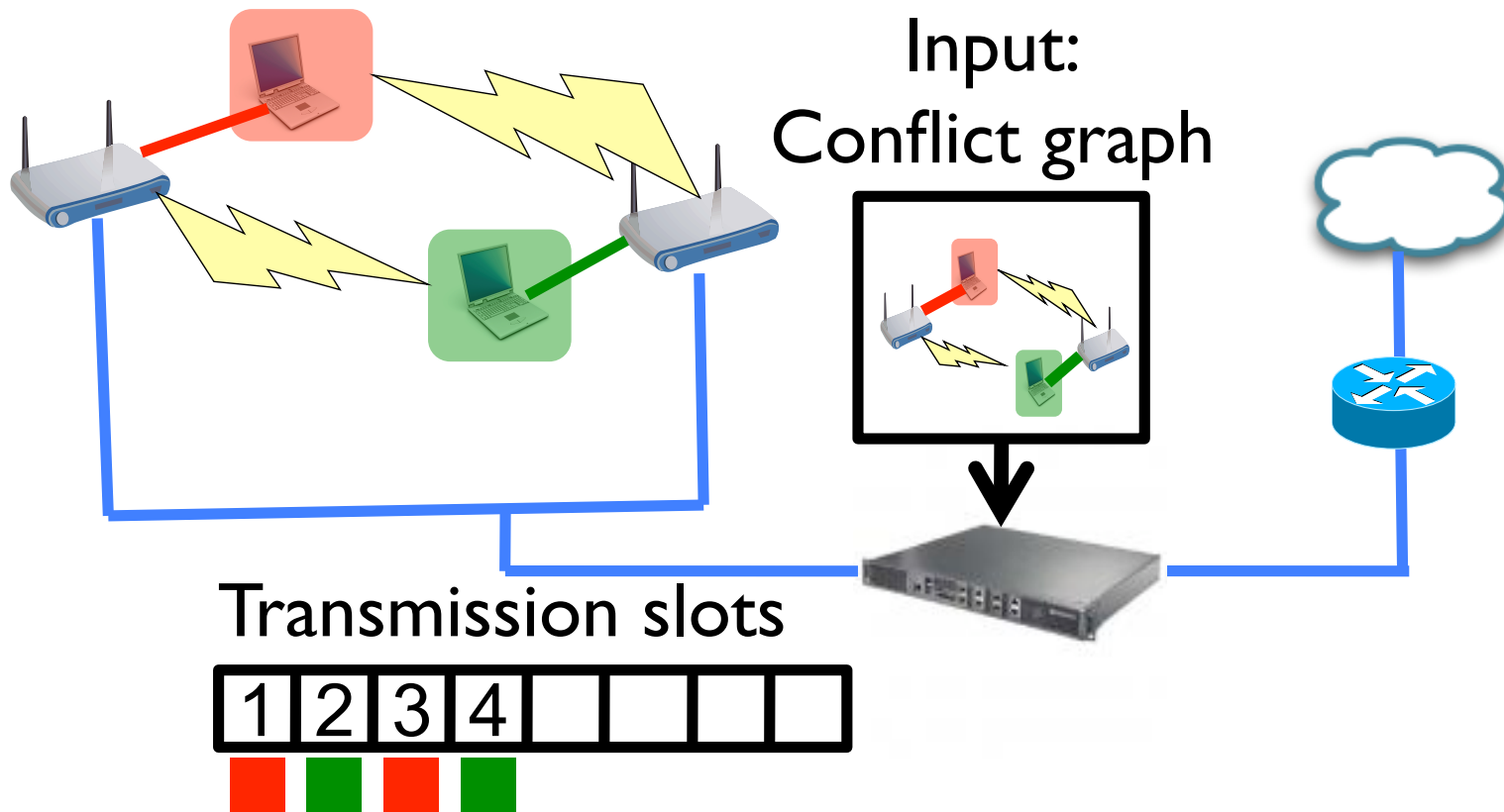
DET: A Simple Deterministic Scheduler

Goal:
Schedule each incoming downlink packet



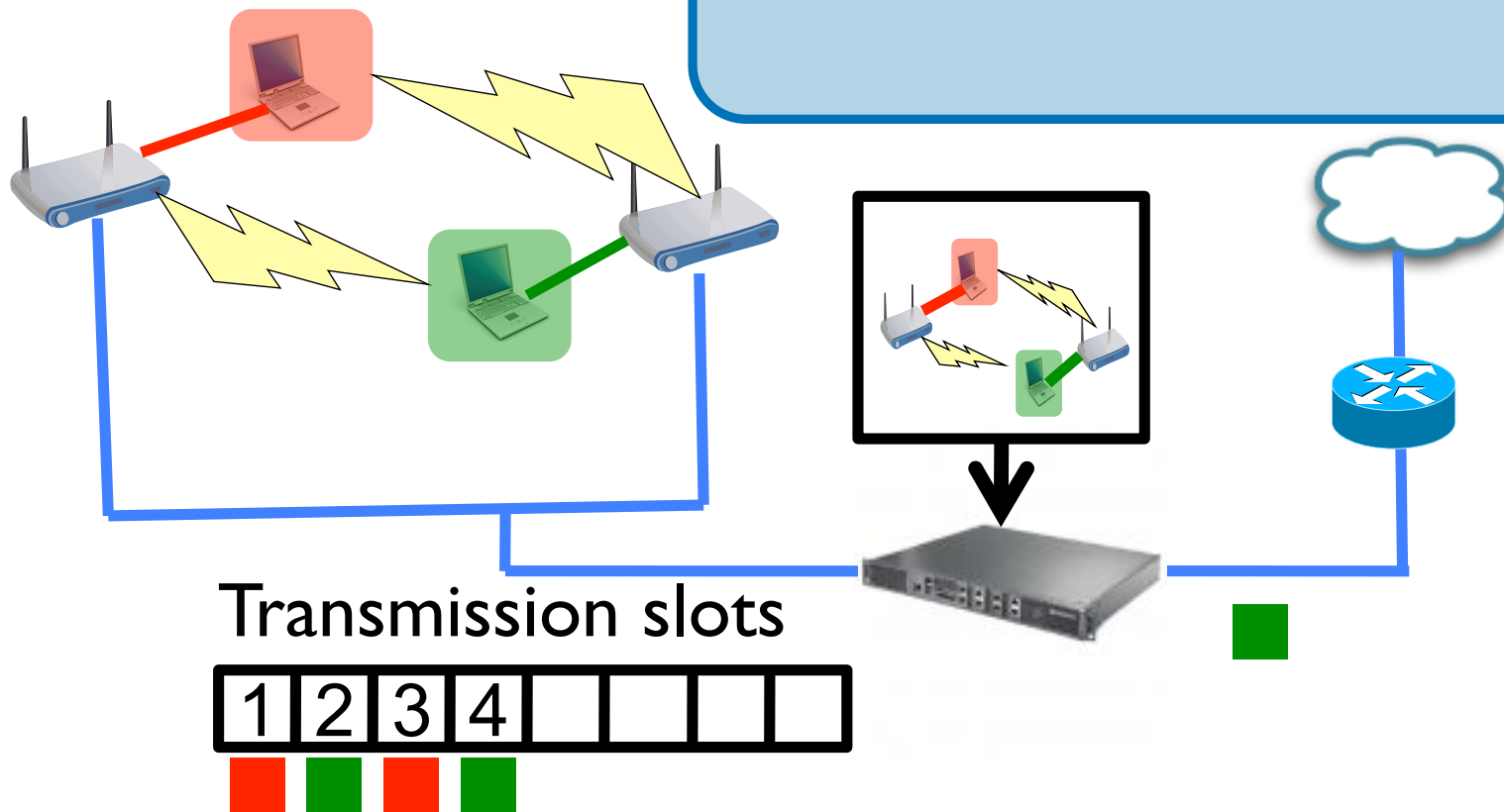
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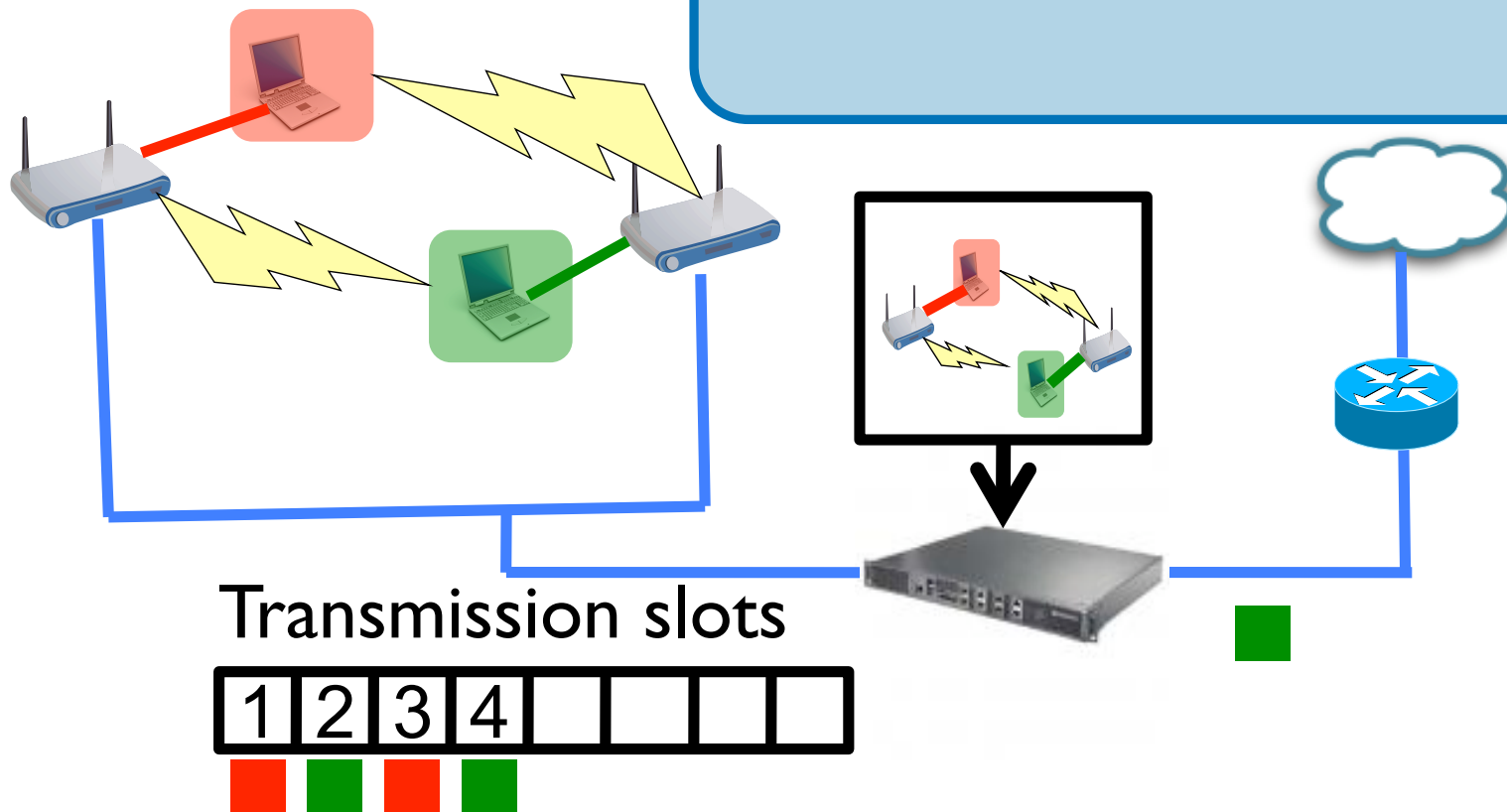
DET: A Simple Deterministic Scheduler

1. New packet arrives



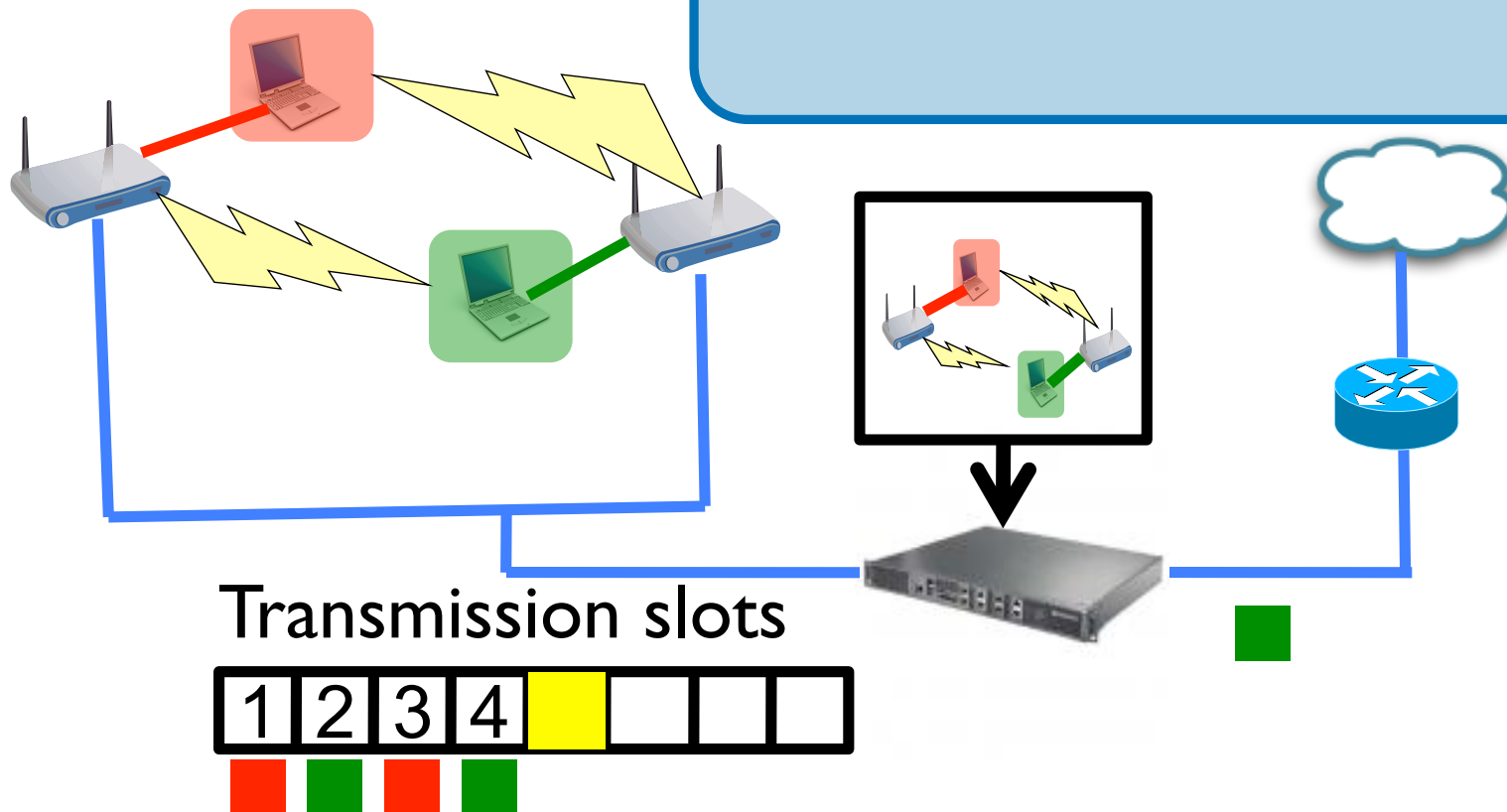
DET: A Simple Deterministic Scheduler

1. New packet arrives
2. Find the earliest conflict free slot



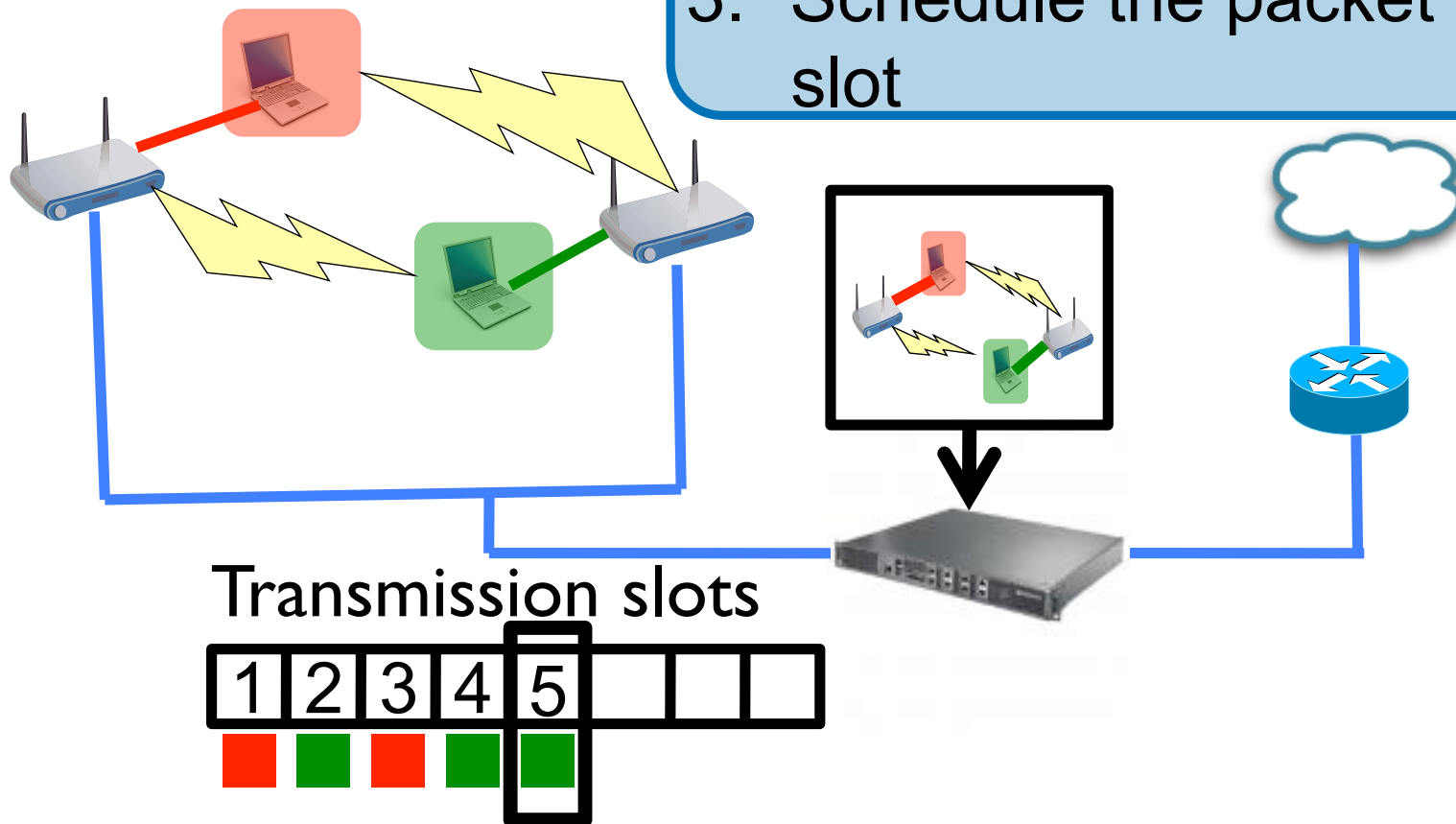
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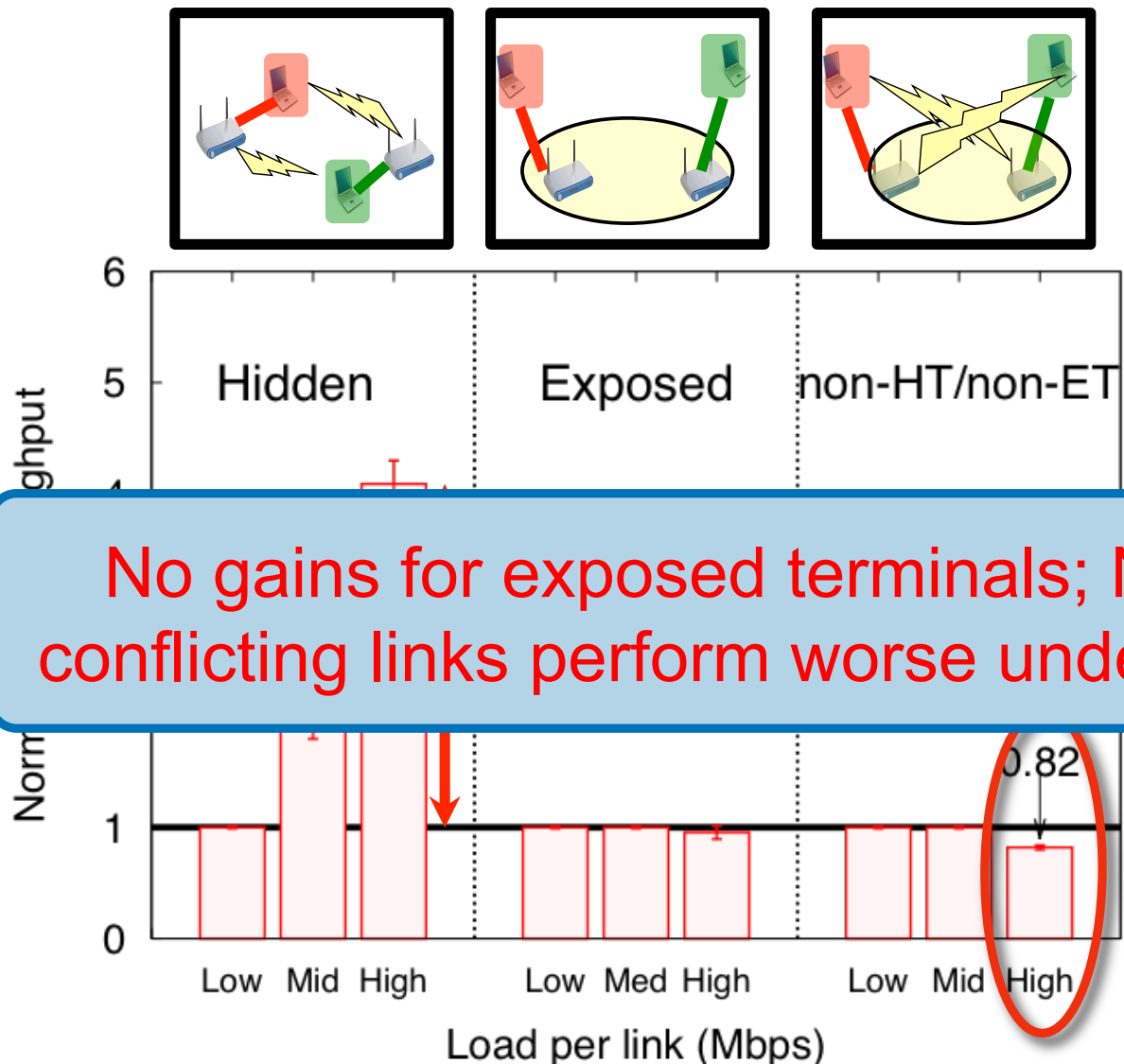


DET: A Simple Deterministic Scheduler

1. New packet arrives
2. Find the earliest conflict free slot
3. Schedule the packet in that slot



Performance of DET



No gains for exposed terminals; Non-conflicting links perform worse under load

Outline

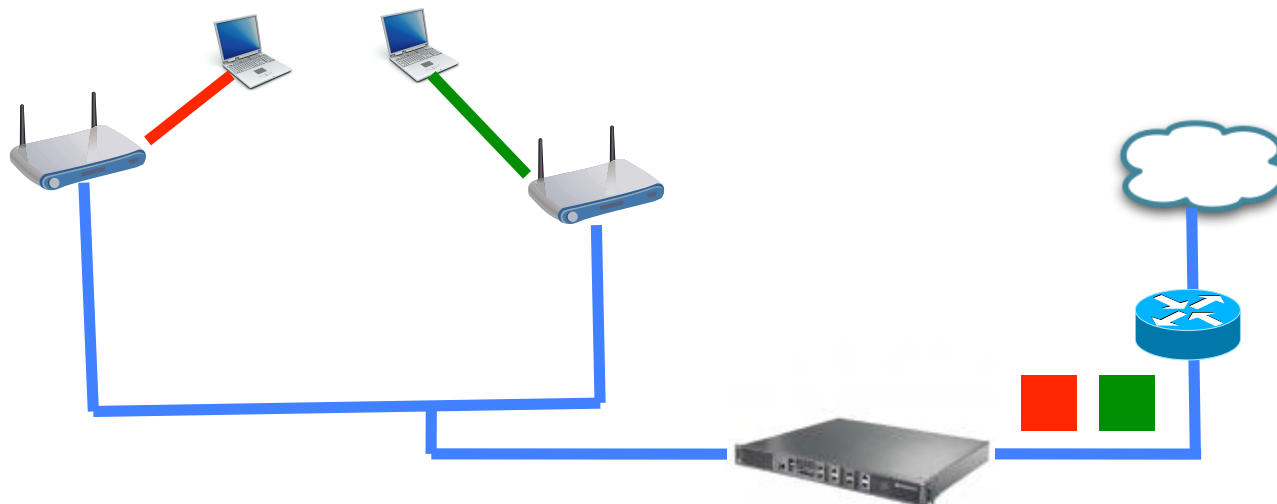
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Overview of CENTAUR

- Incorporate basic DET scheduler
- Tackle DET's shortcomings:
 - Amortize scheduling overhead
 - Improve performance for exposed links
 - Avoid degrading normal links
- Coexist with non-enterprise and uplink traffic

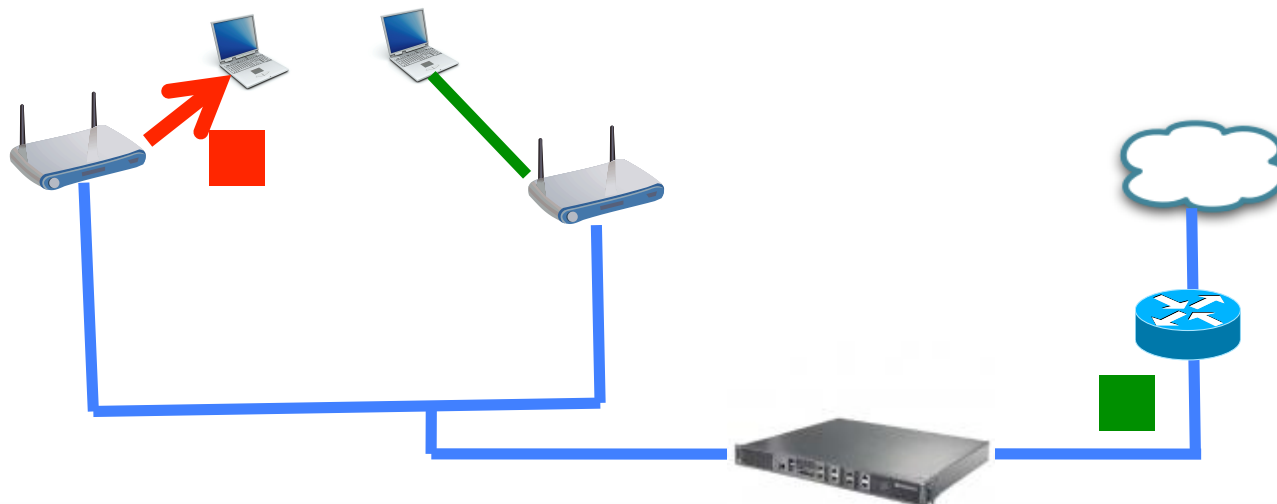
(1) Avoid Scheduling Overheads

- Problem: Per-packet scheduling performs poorly under high network loads
- Solution:
 - Schedule packets in batches (or epochs)



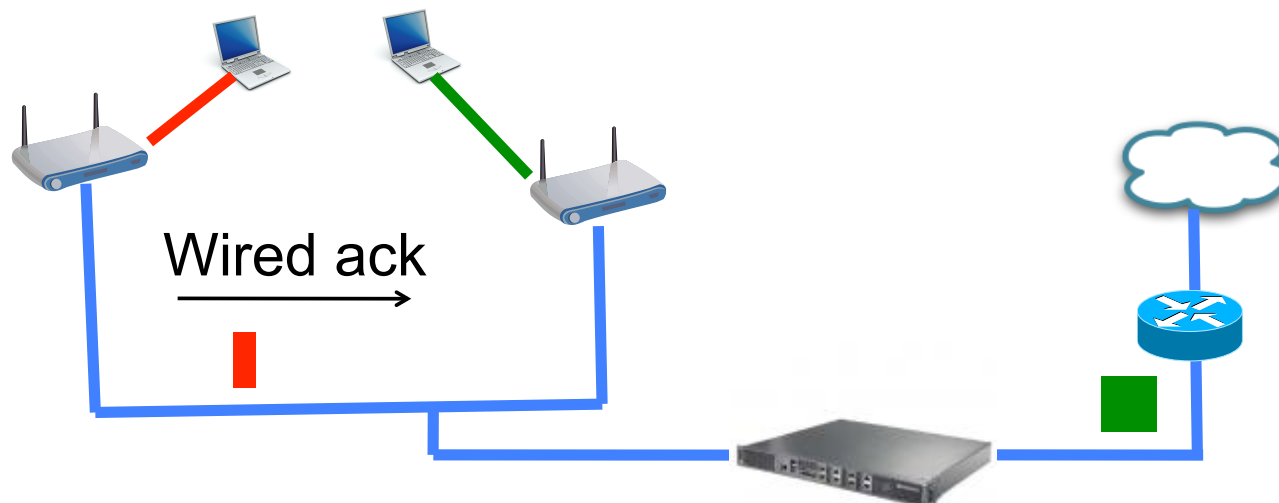
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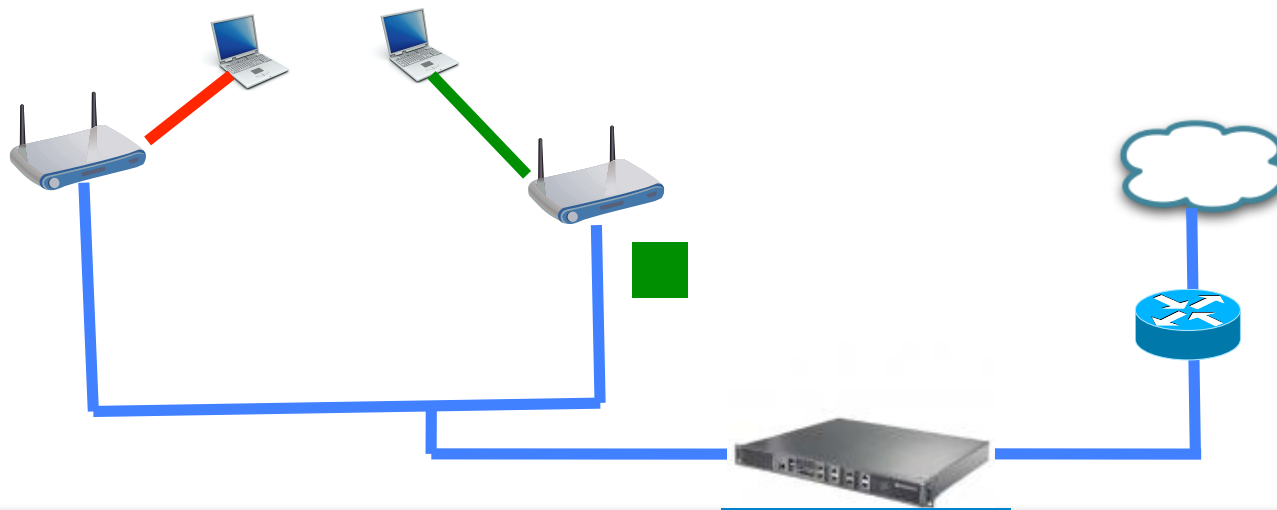
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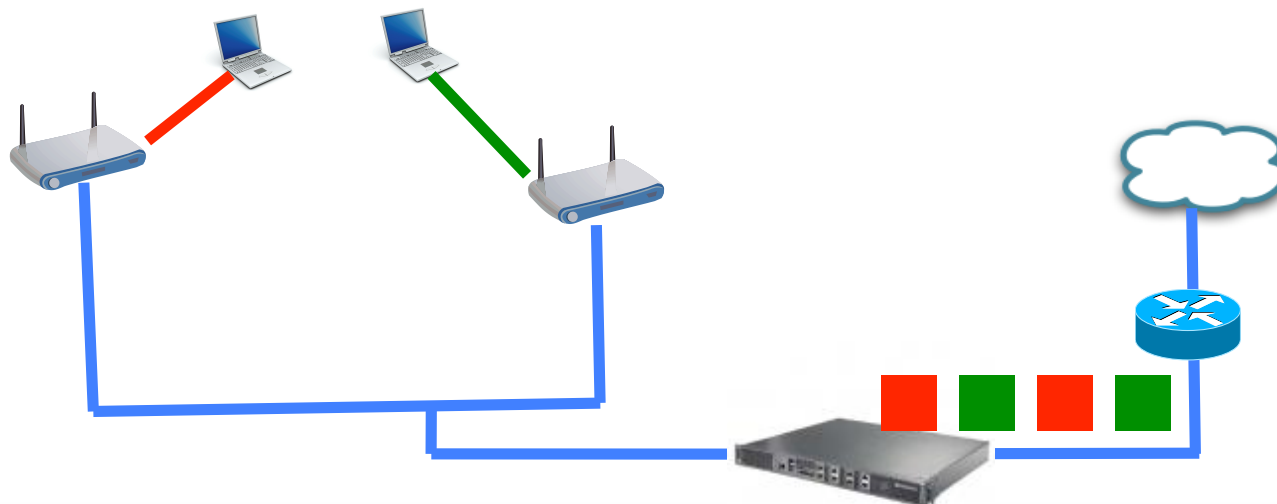
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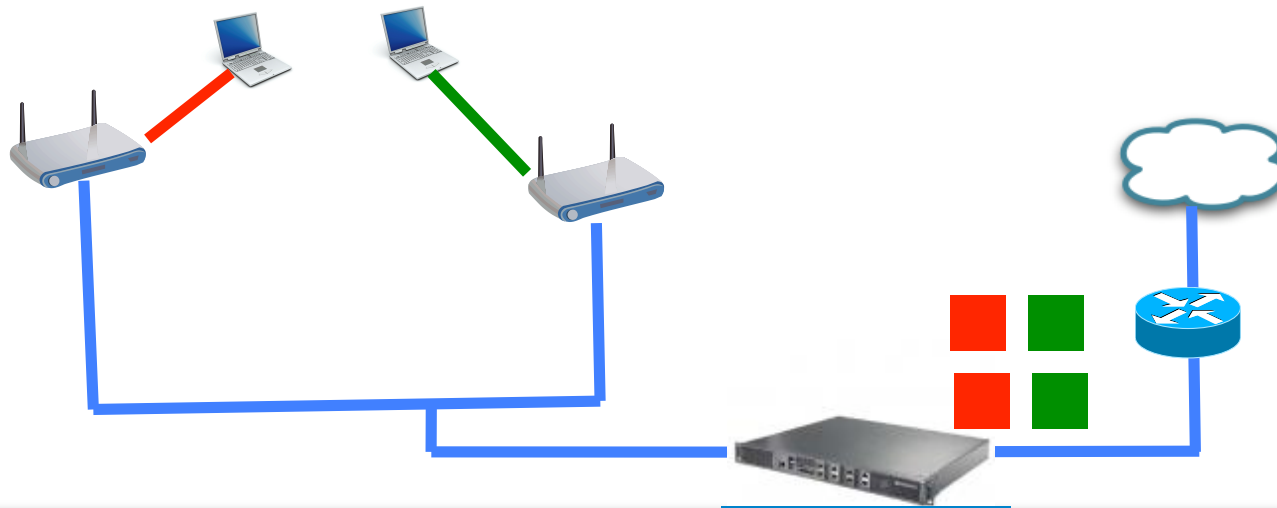
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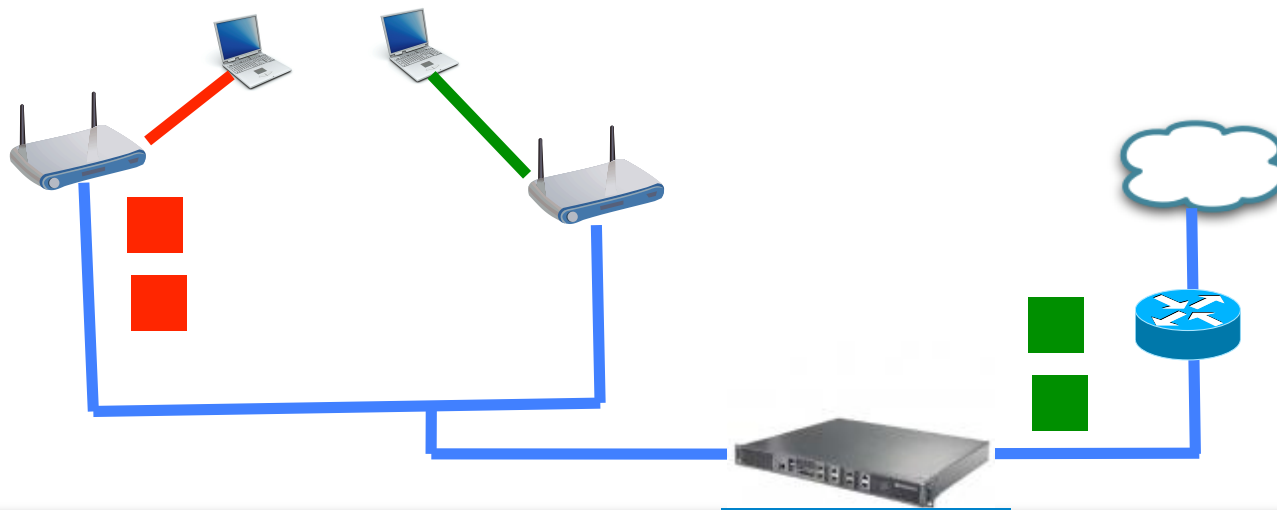
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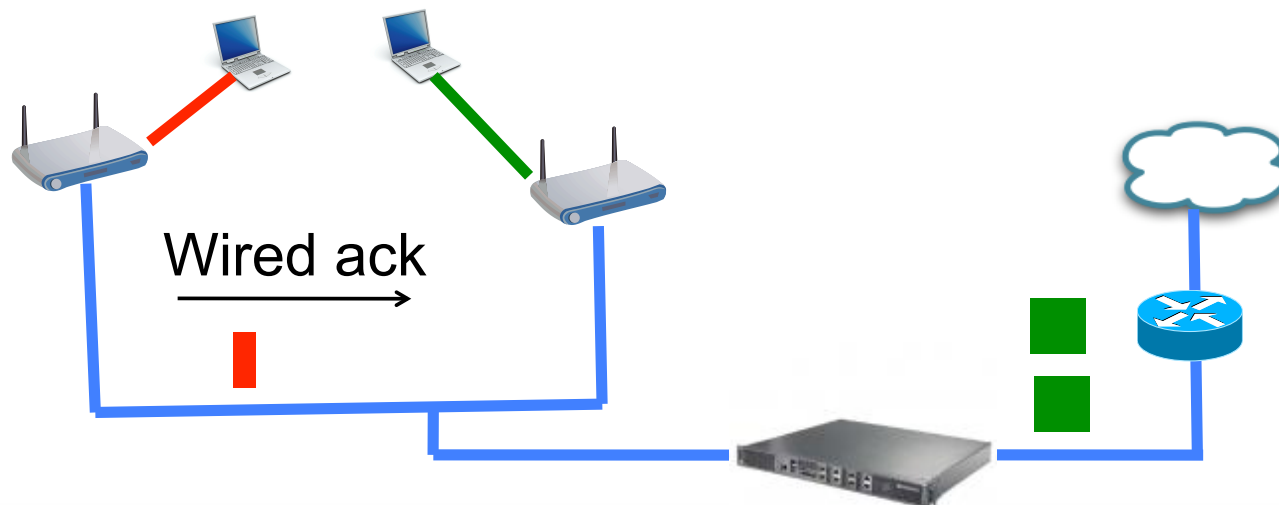
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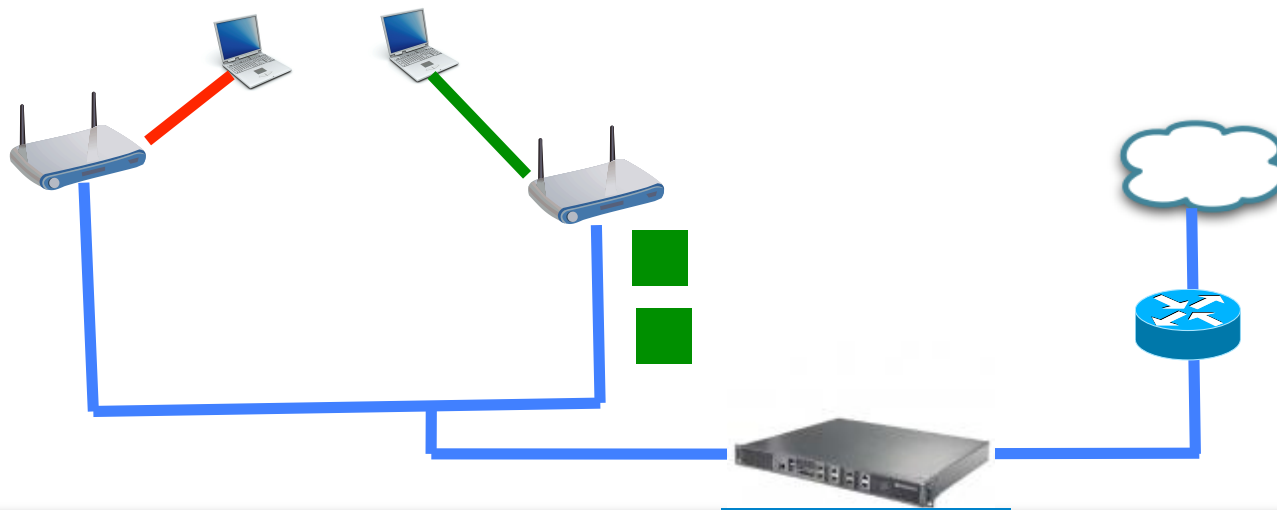
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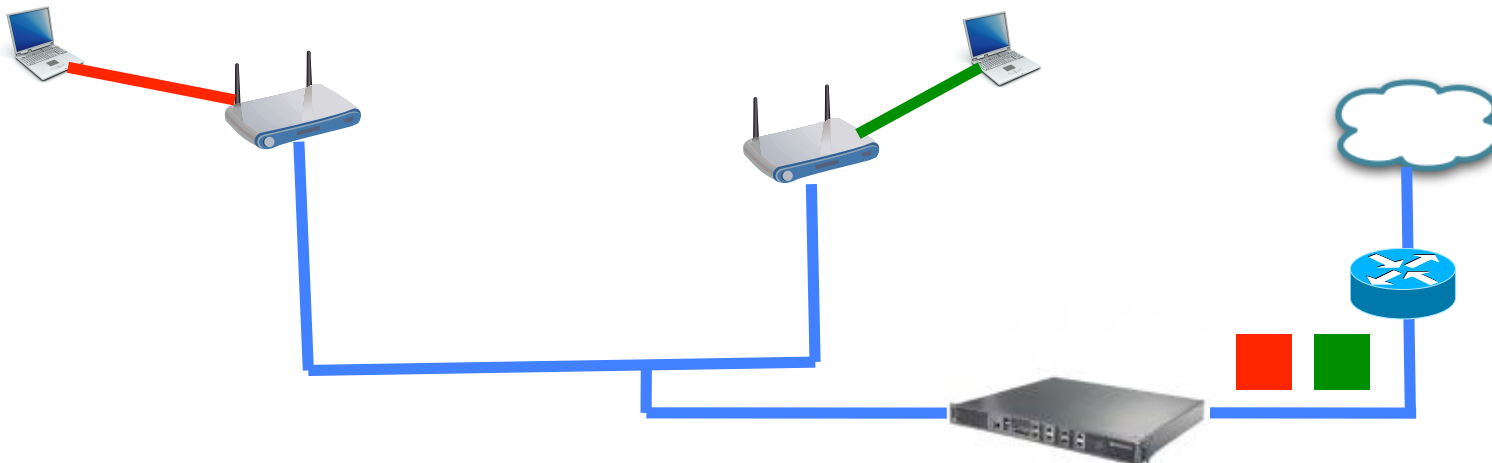
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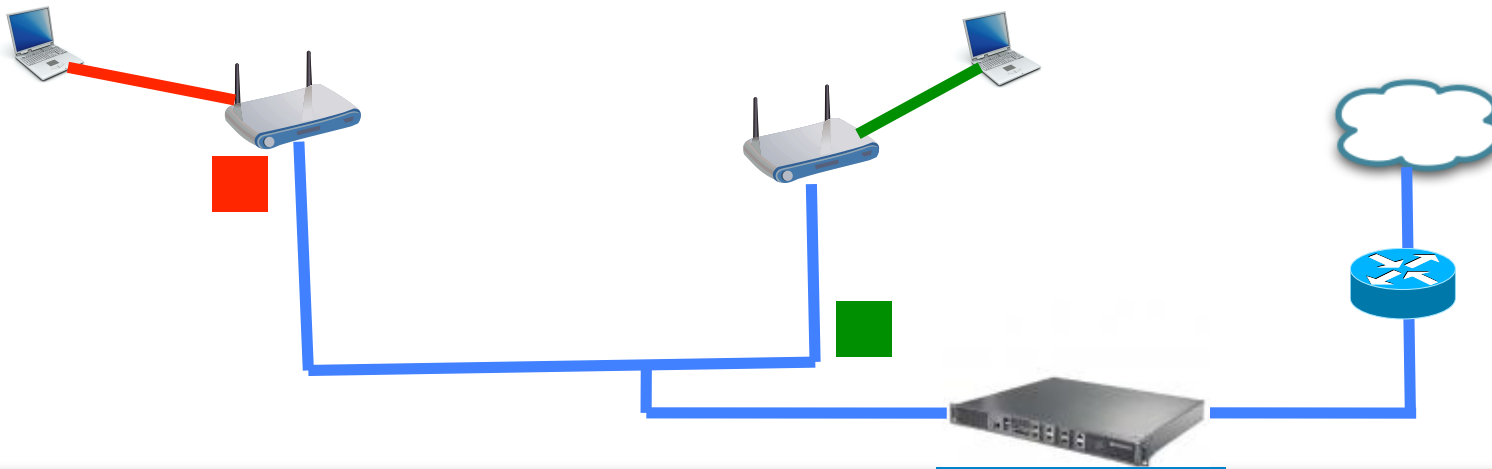
(2) Improve Exposed Terminals

- Problem: Exposed links can operate in parallel but don't due to *carrier-sensing*
- Solution:
 - Schedule batch of packets, fix backoff periods



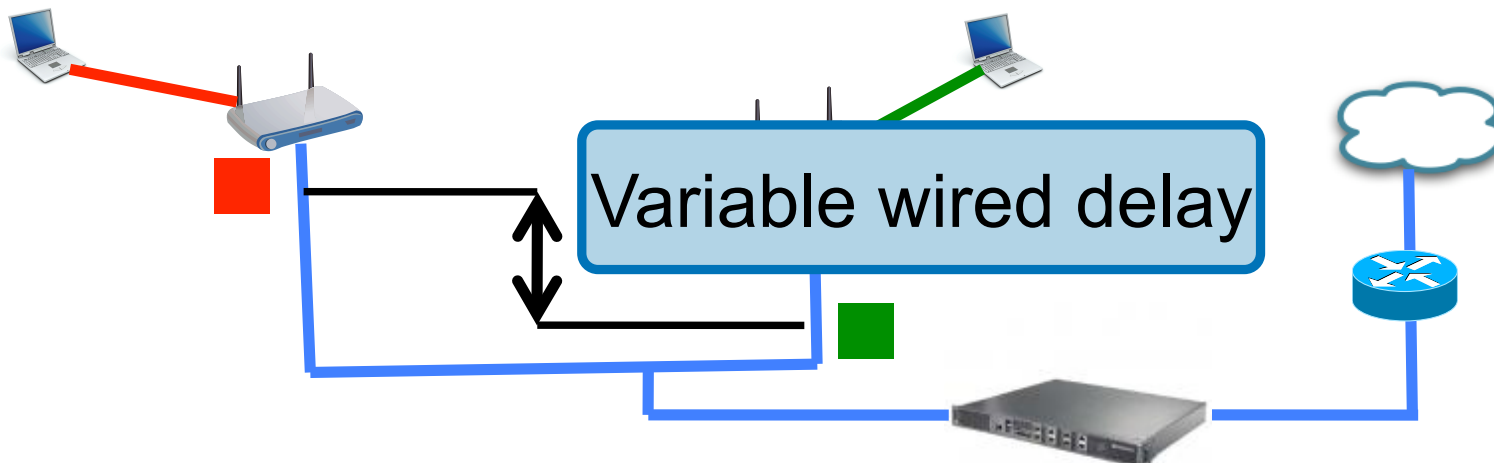
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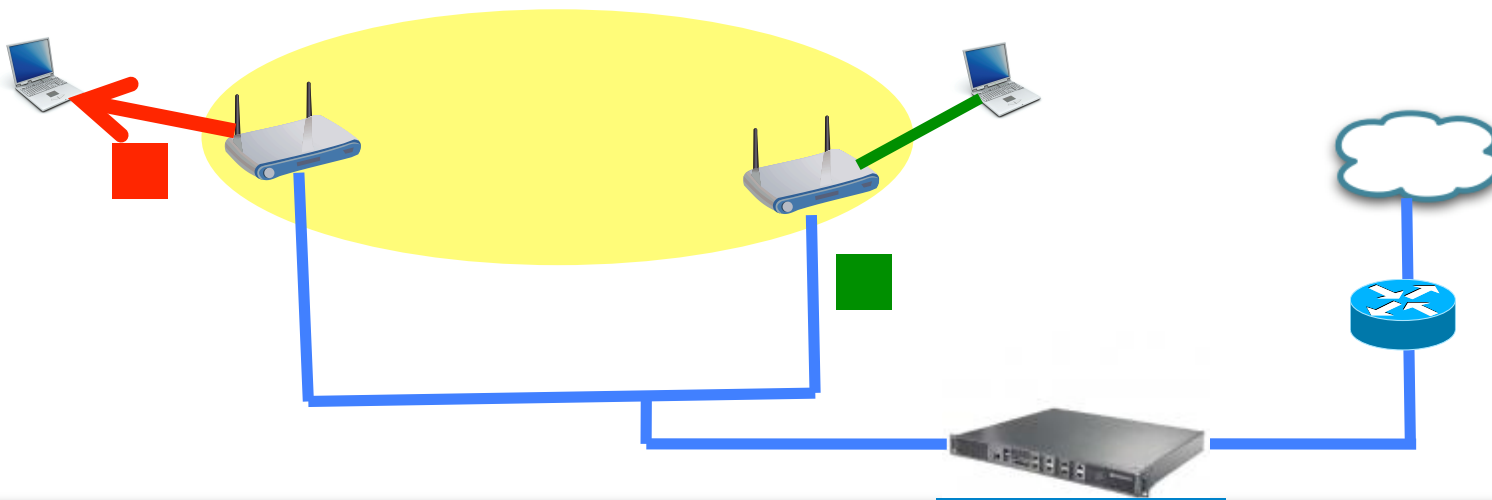
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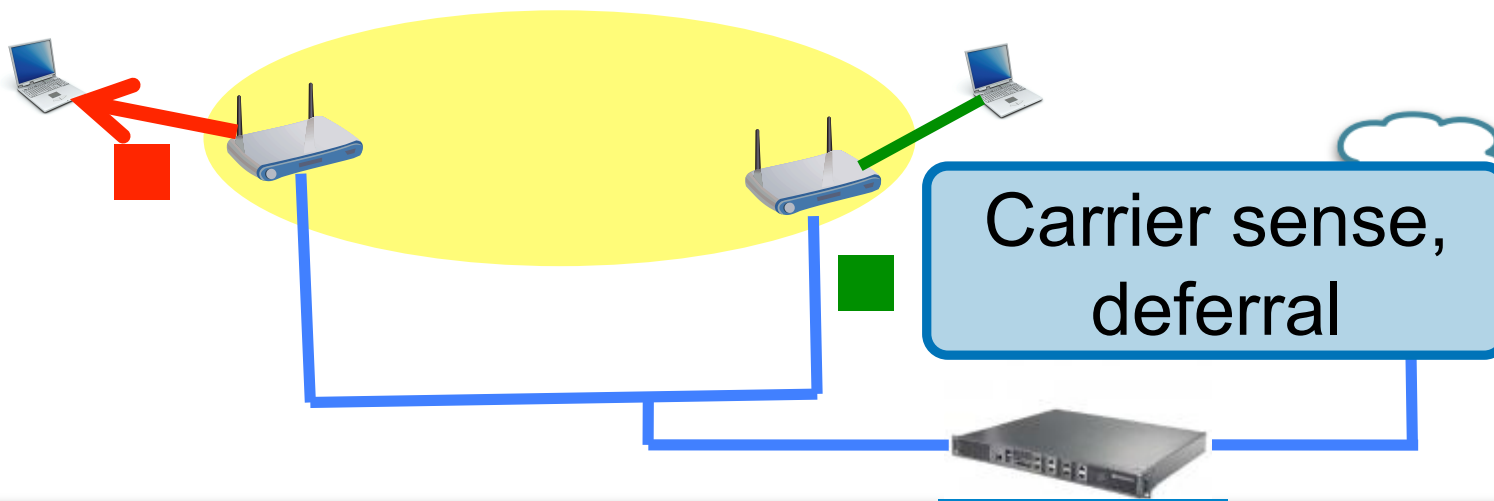
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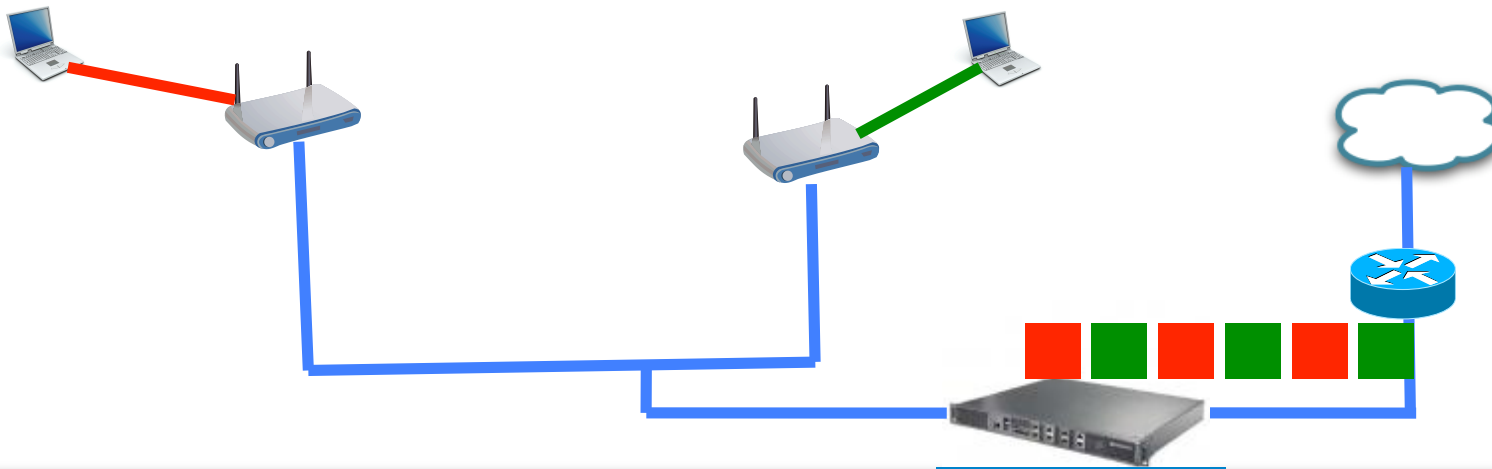
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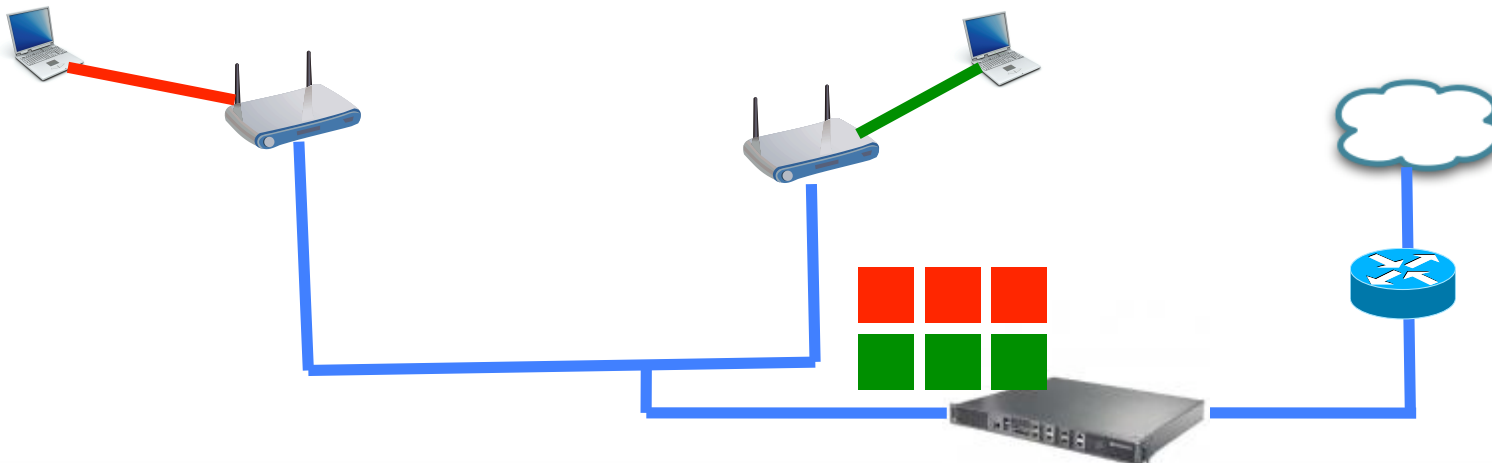
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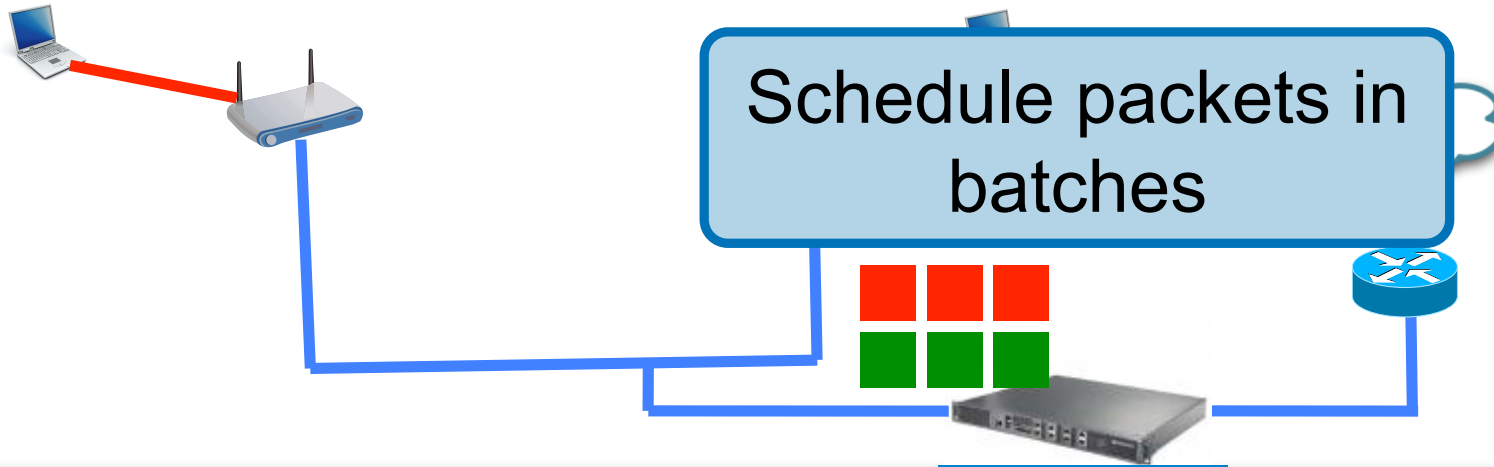
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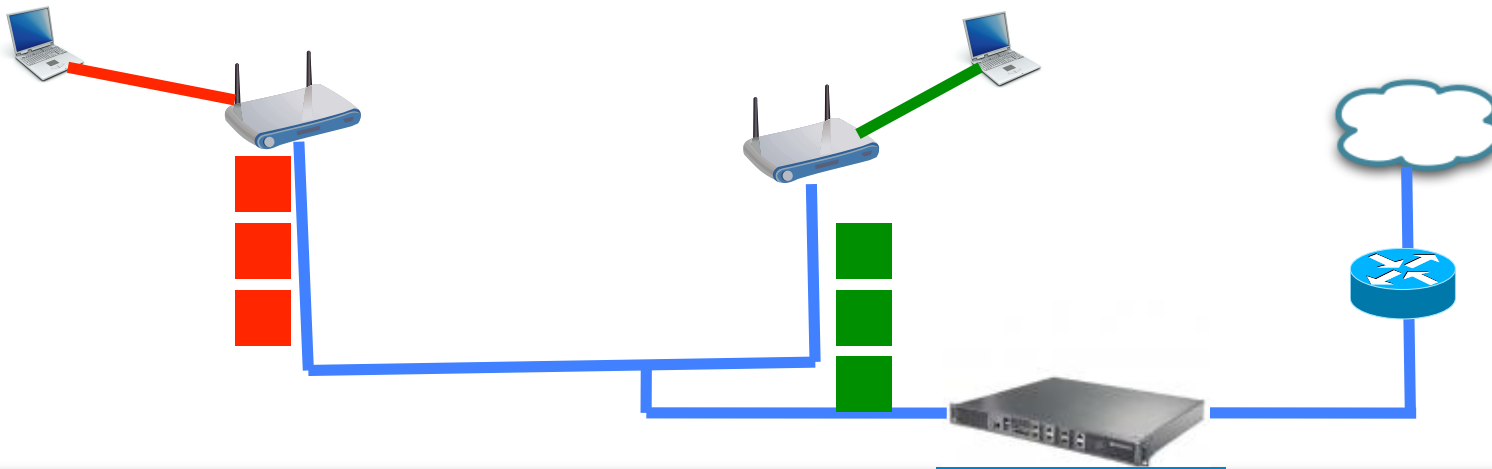
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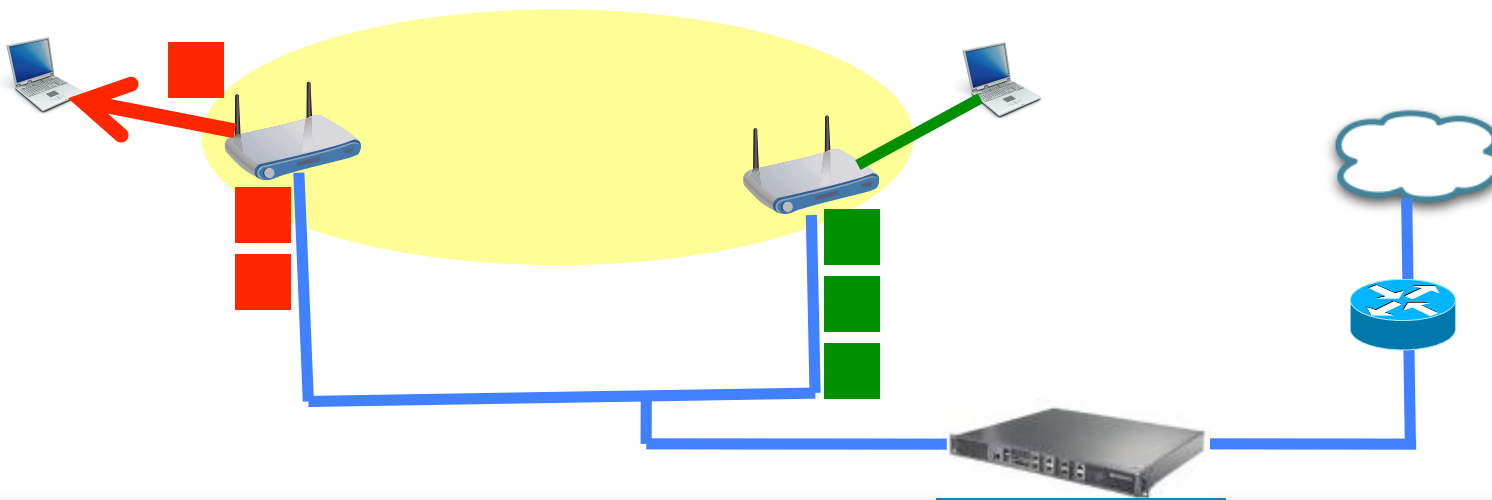
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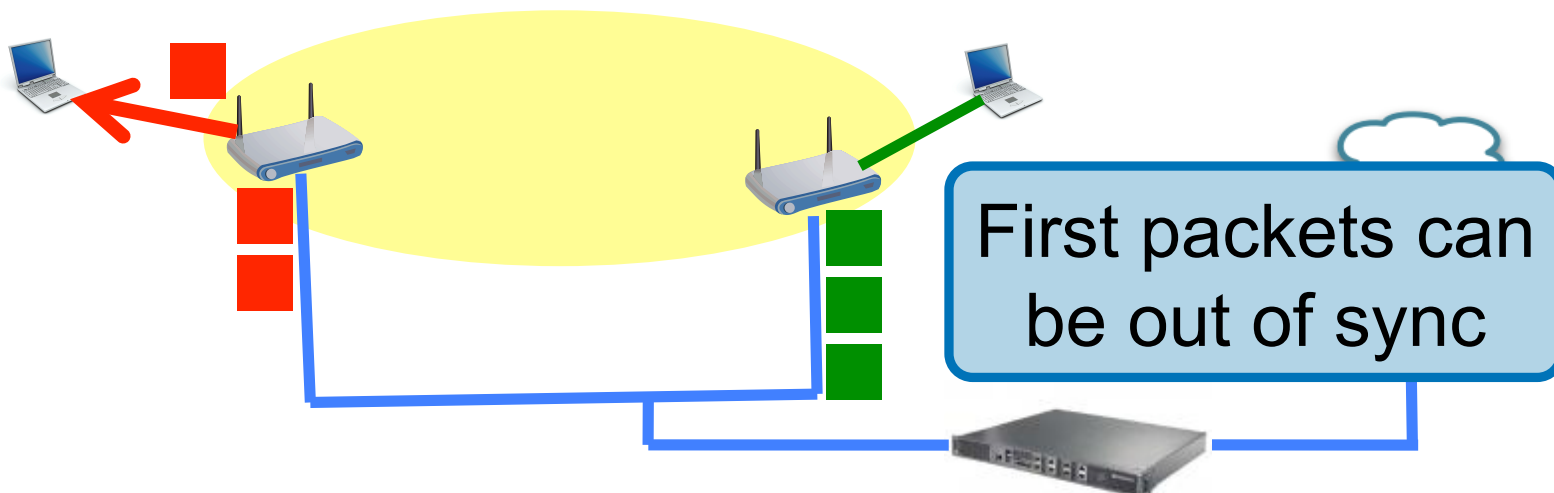
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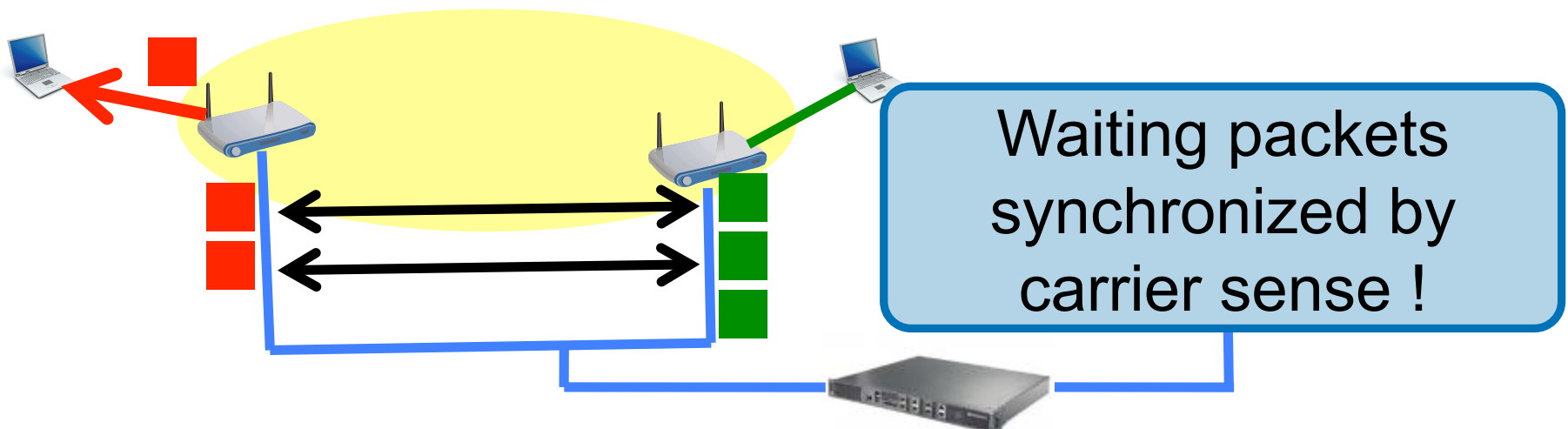
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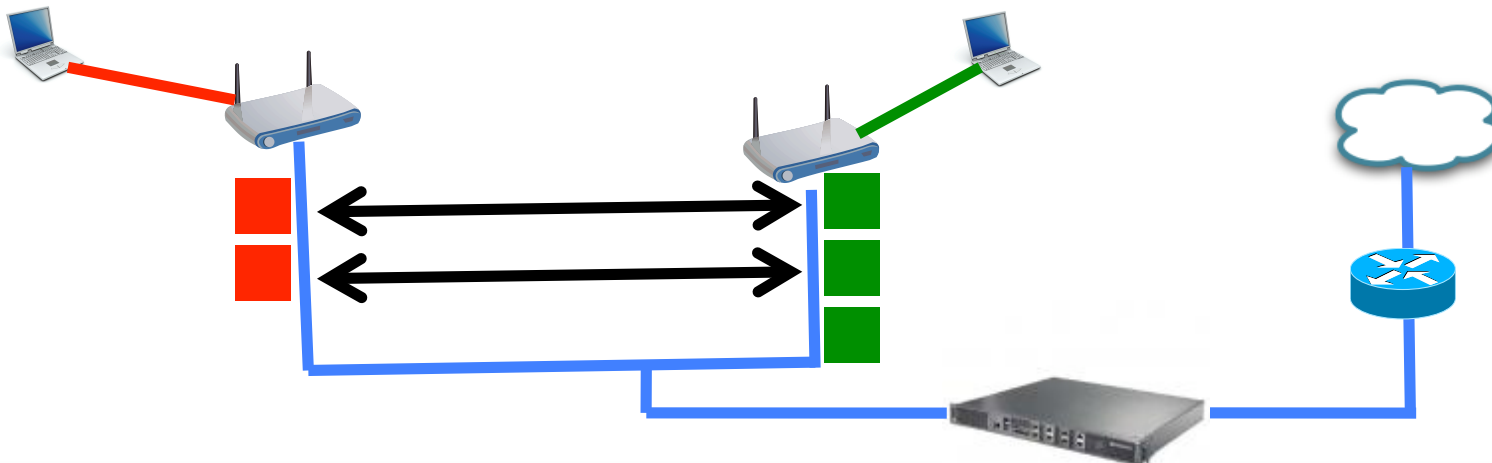
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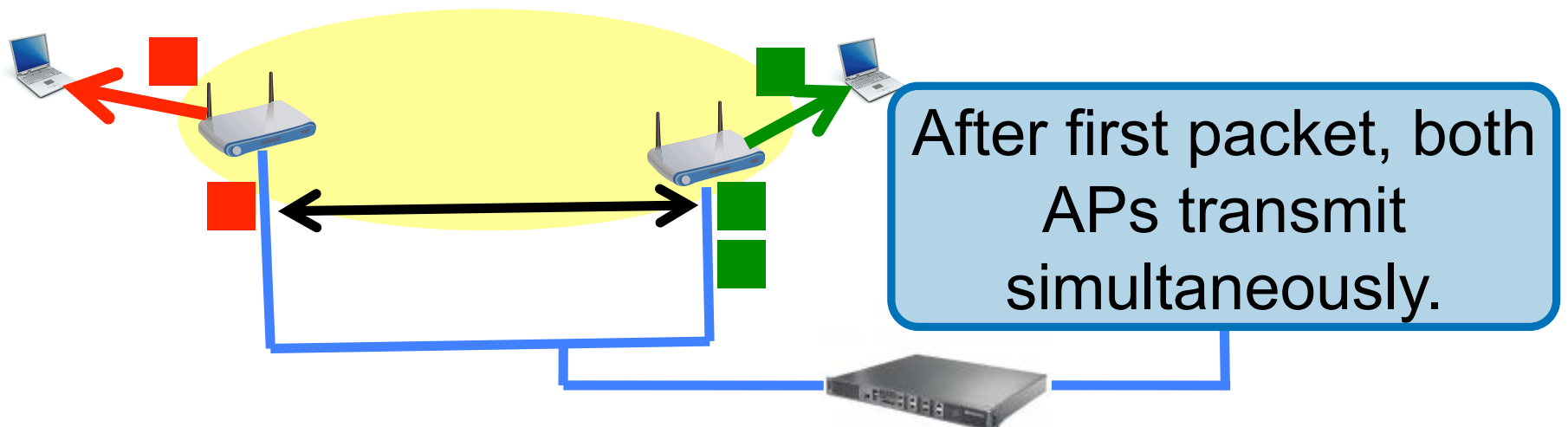
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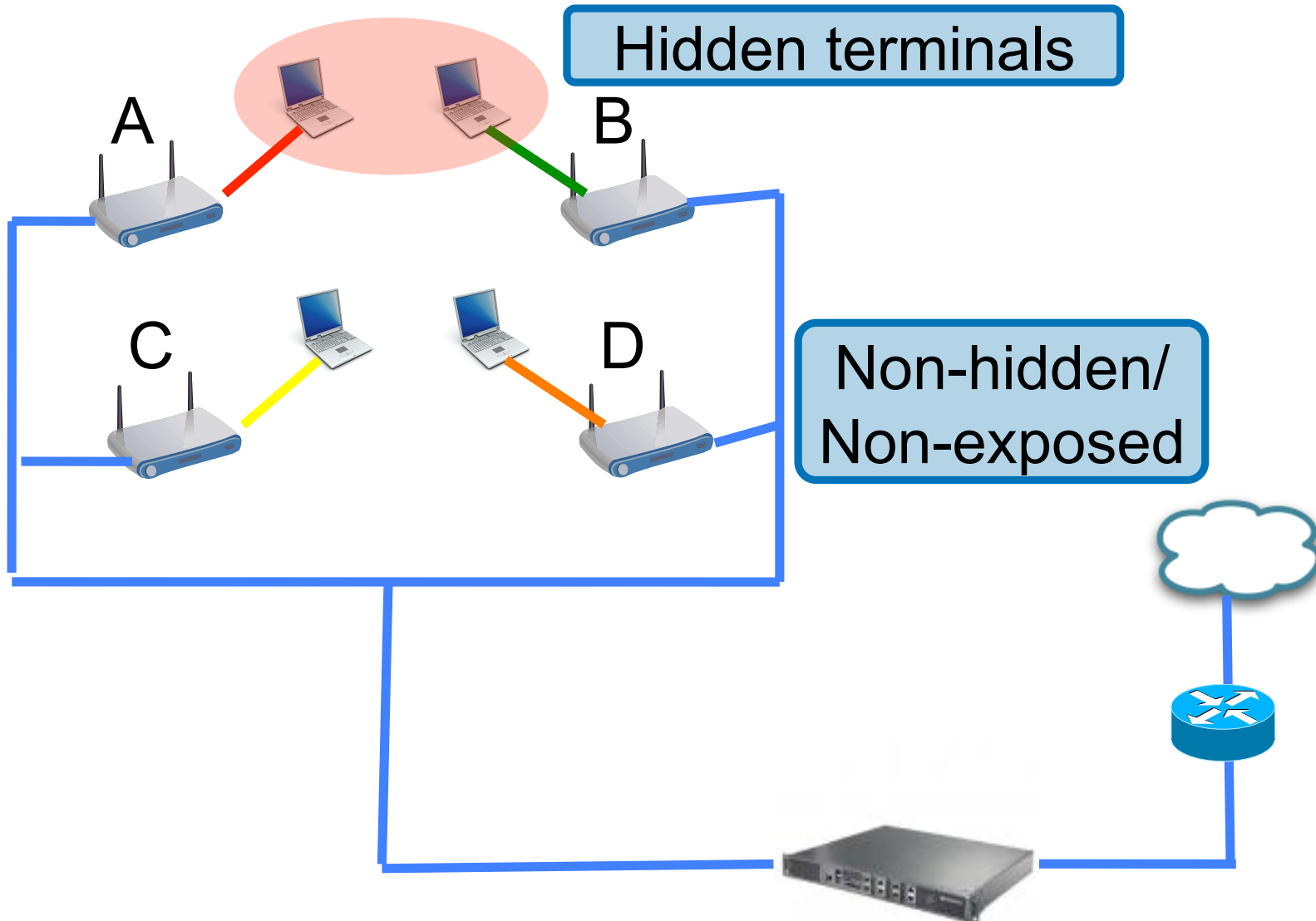


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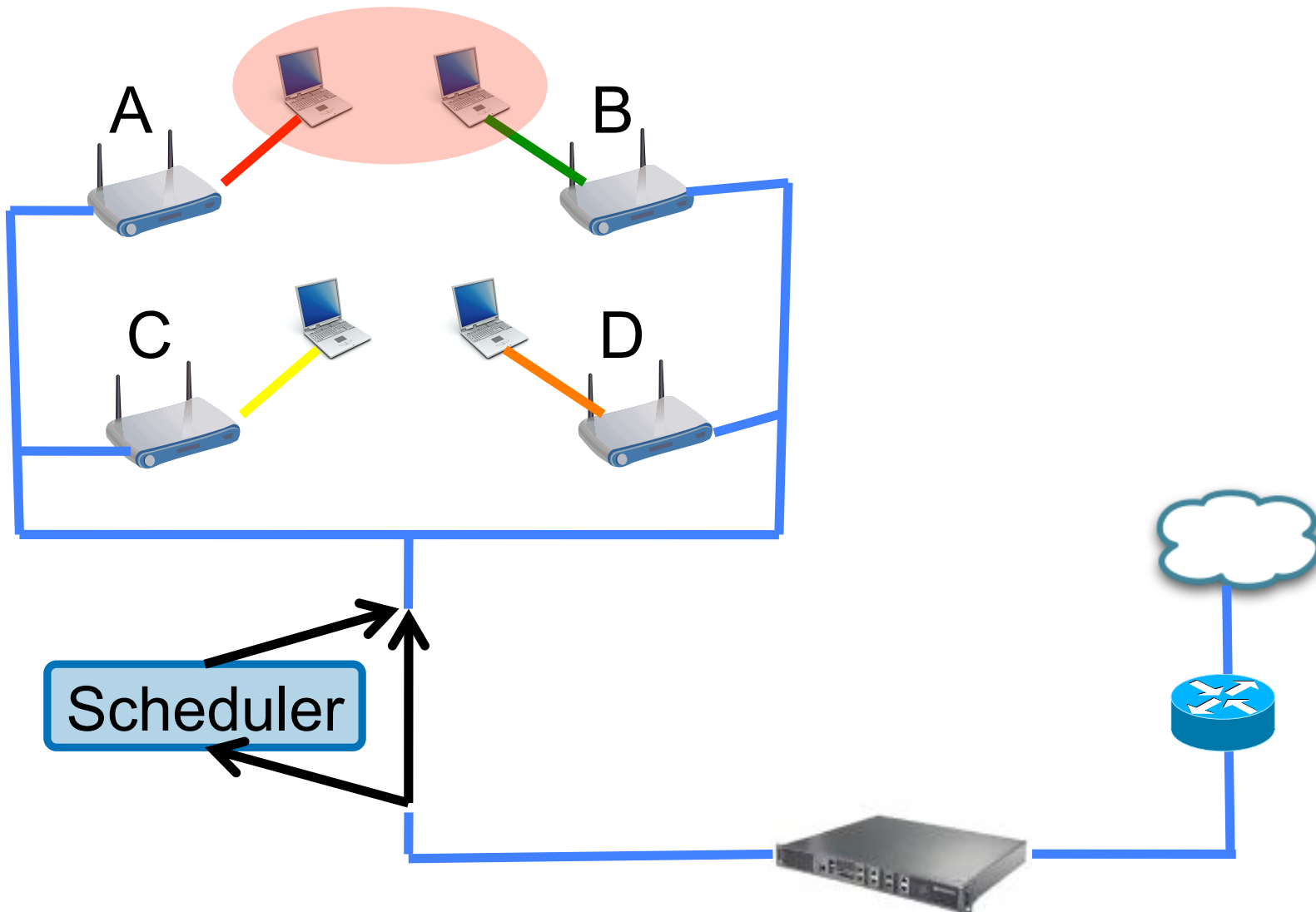
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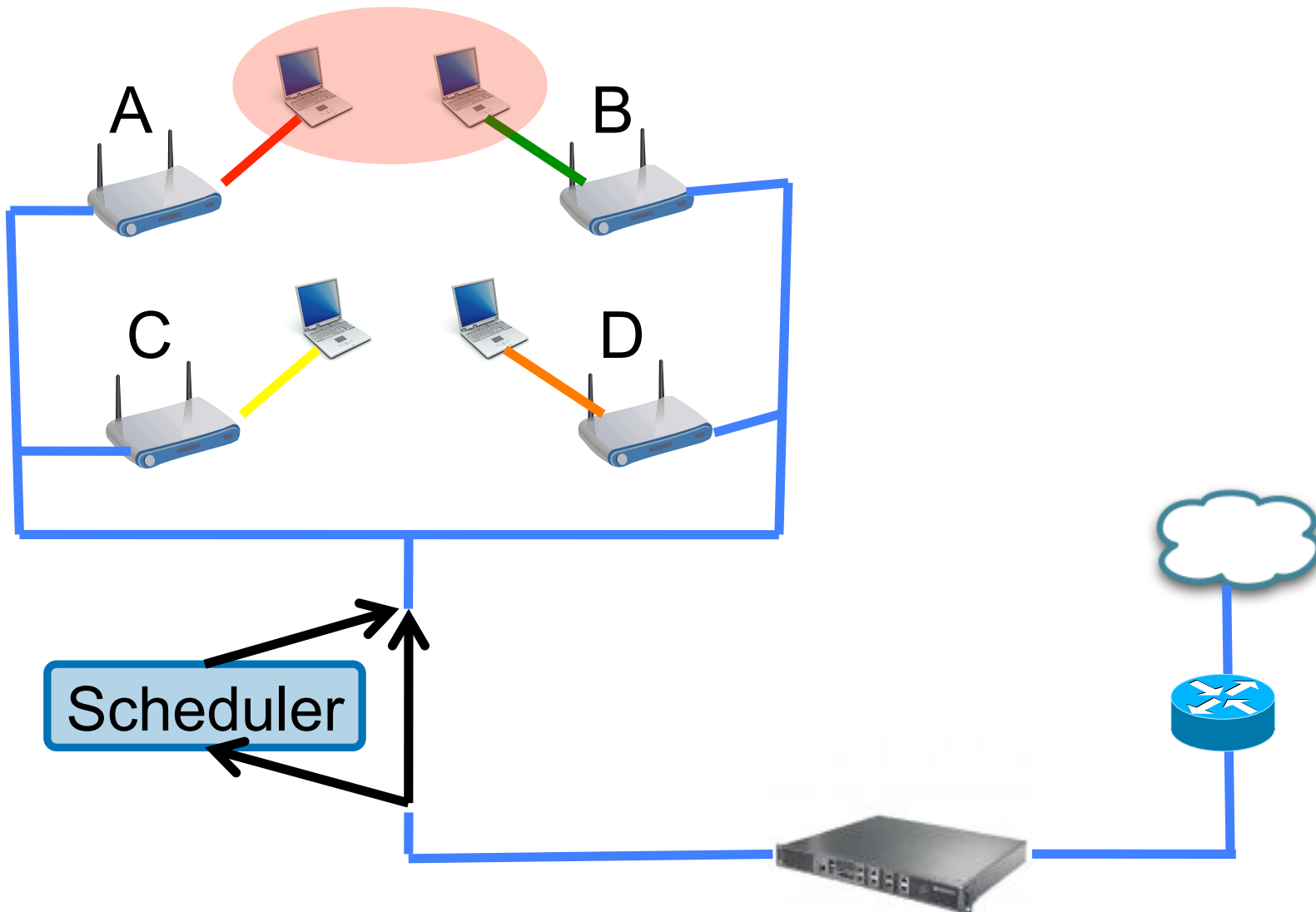
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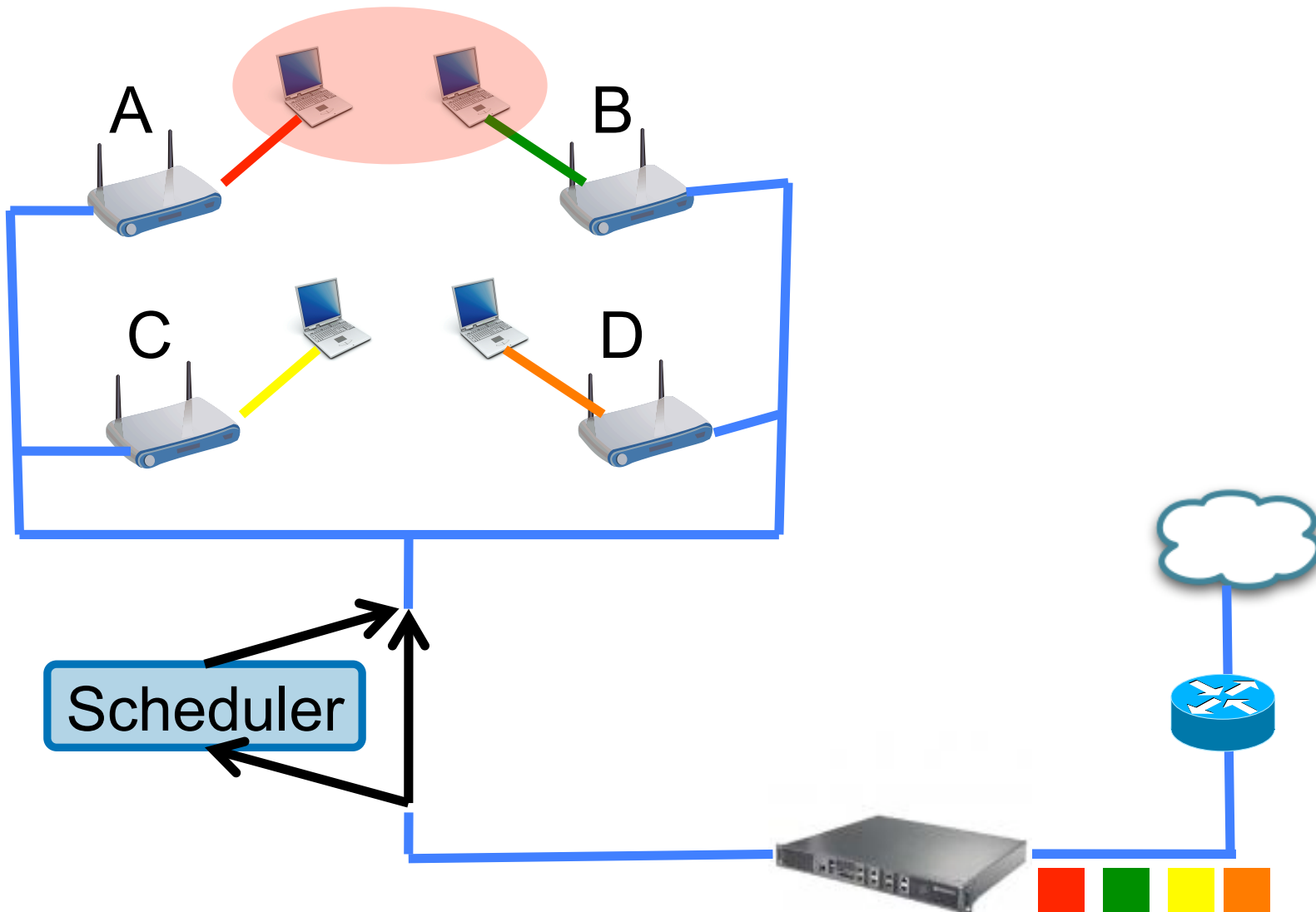
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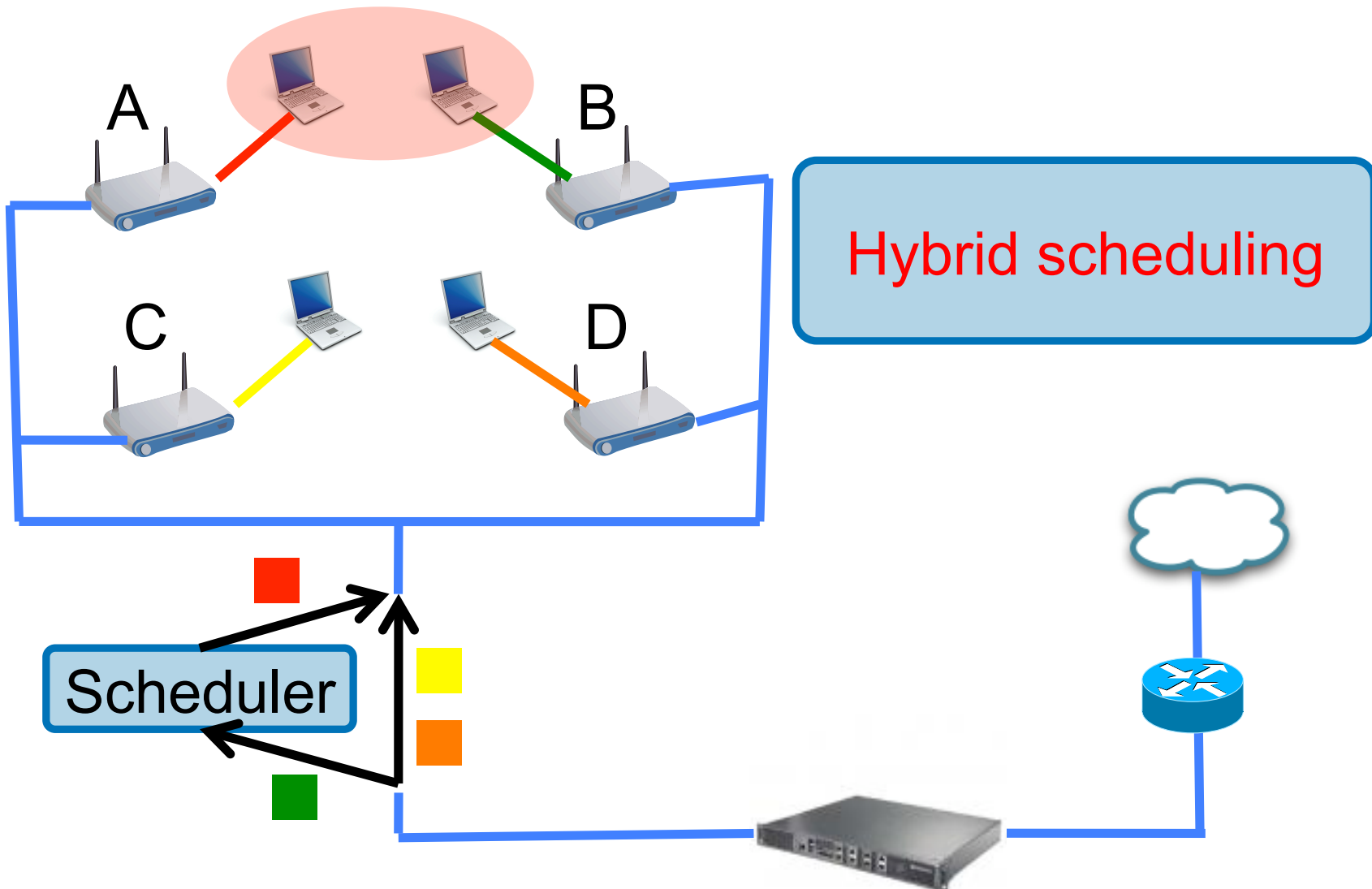
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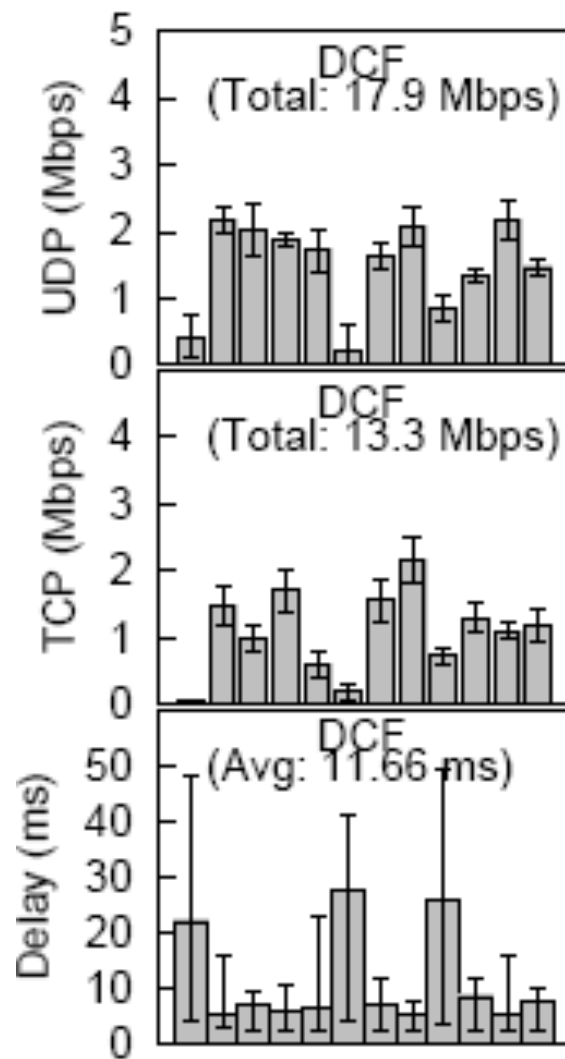
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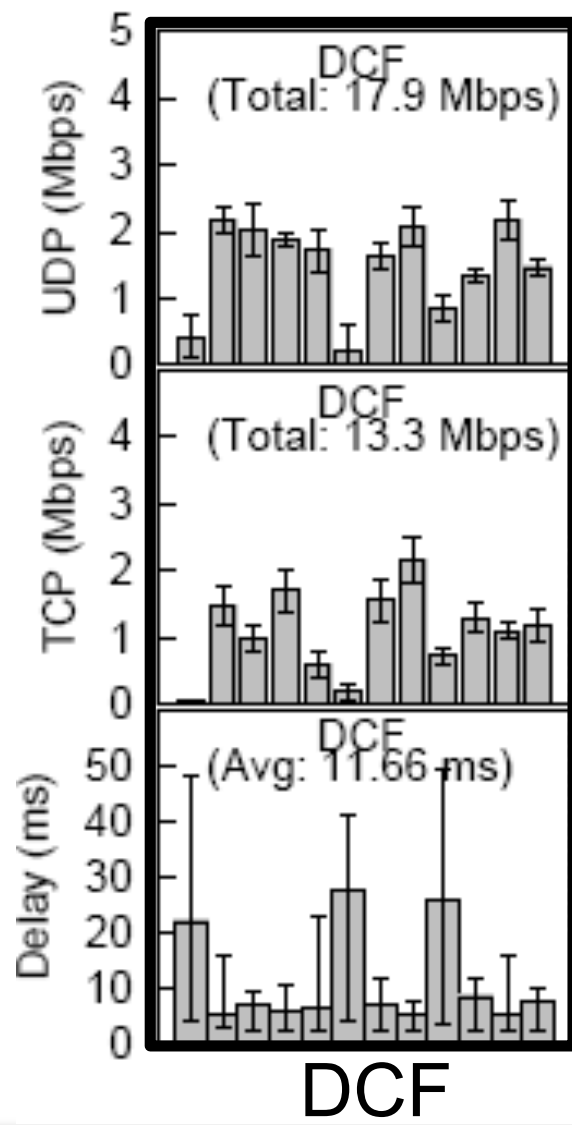
Large-Scale Experiments

- Platform:
 - Two WLAN testbeds in separate buildings
- Topology
 - Representative: 7 APs, 12 clients
- Traffic and metrics
 - UDP, TCP, VoIP, HTTP (real traces)
 - Throughput, delay, MOS, web transaction delay

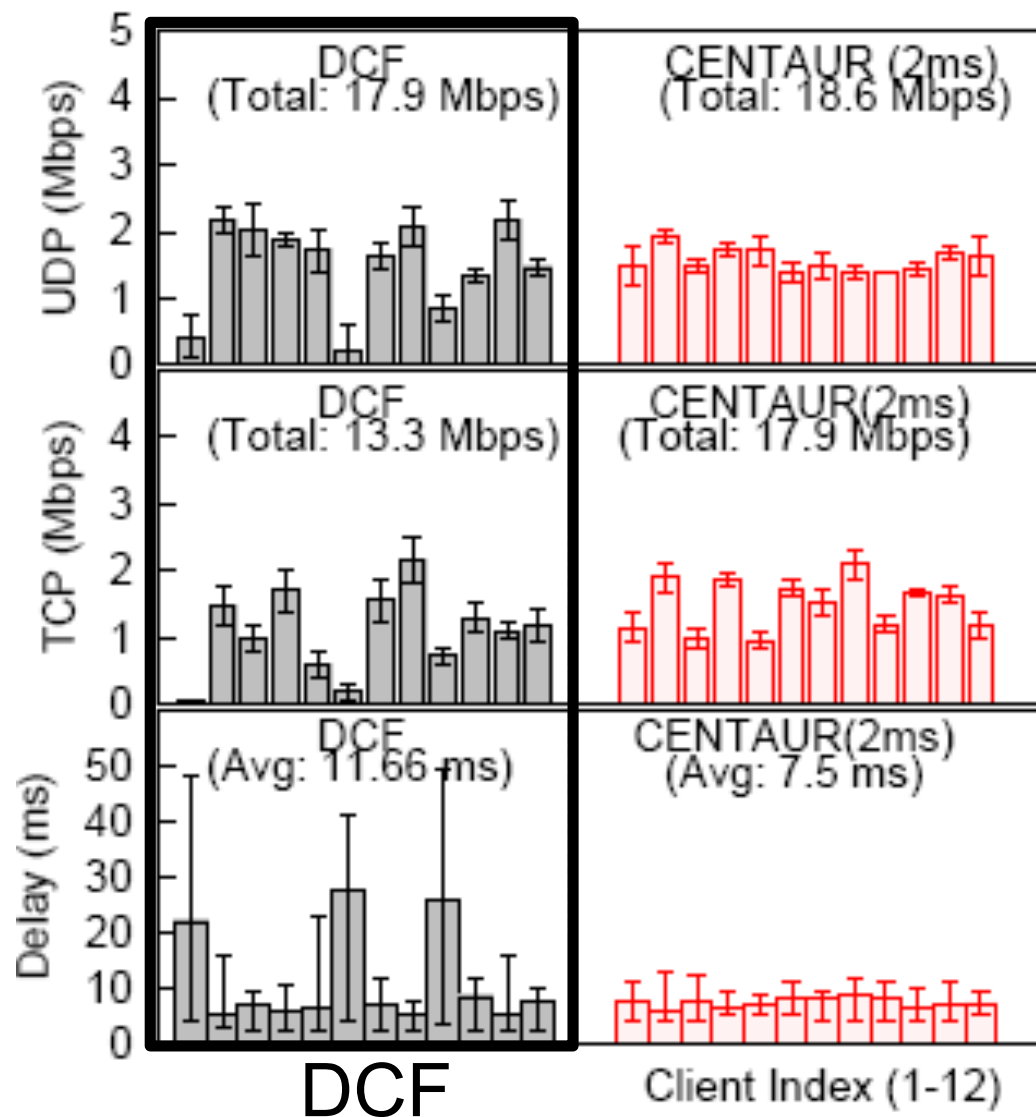
Result 1: UDP/TCP Performance



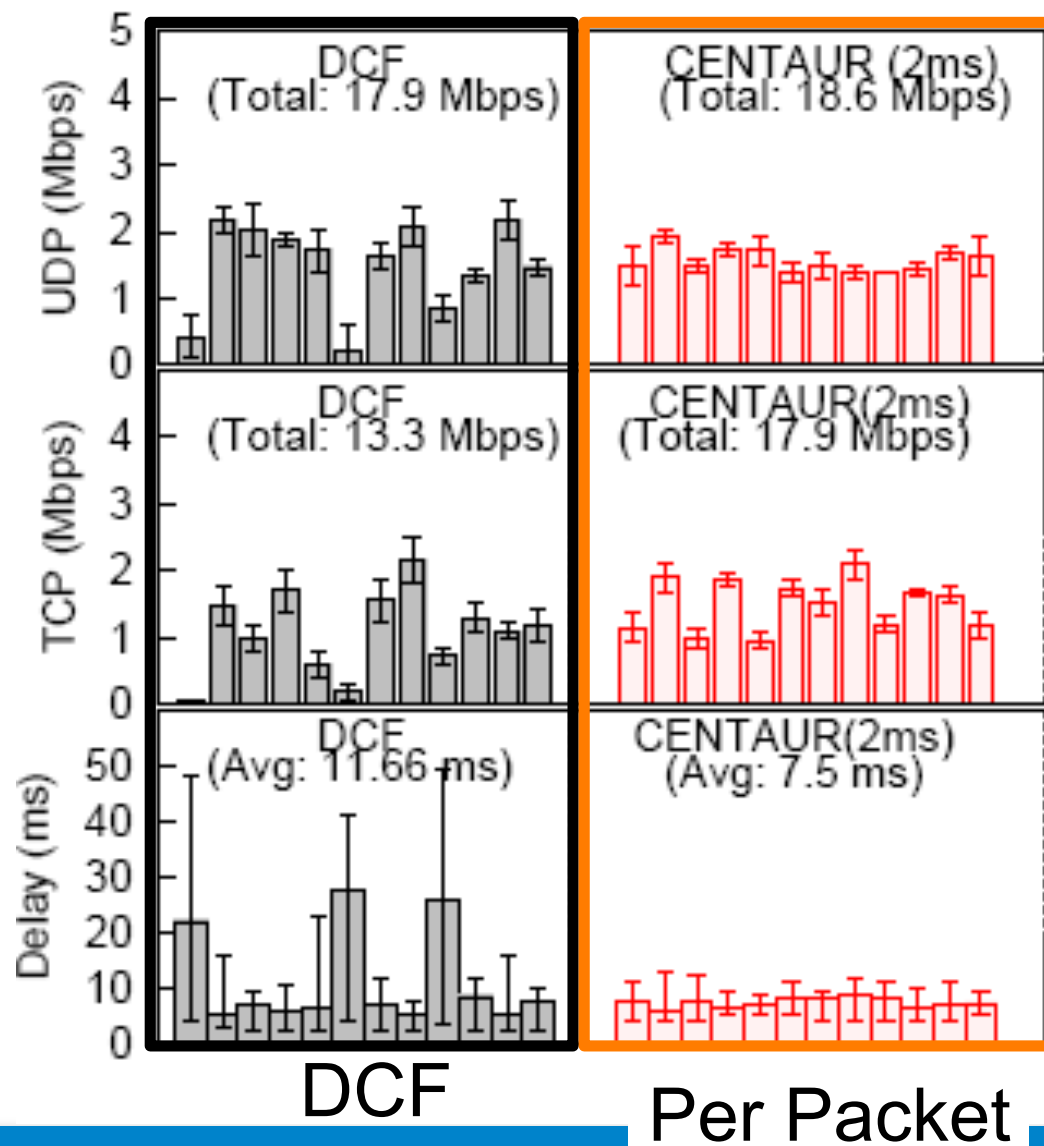
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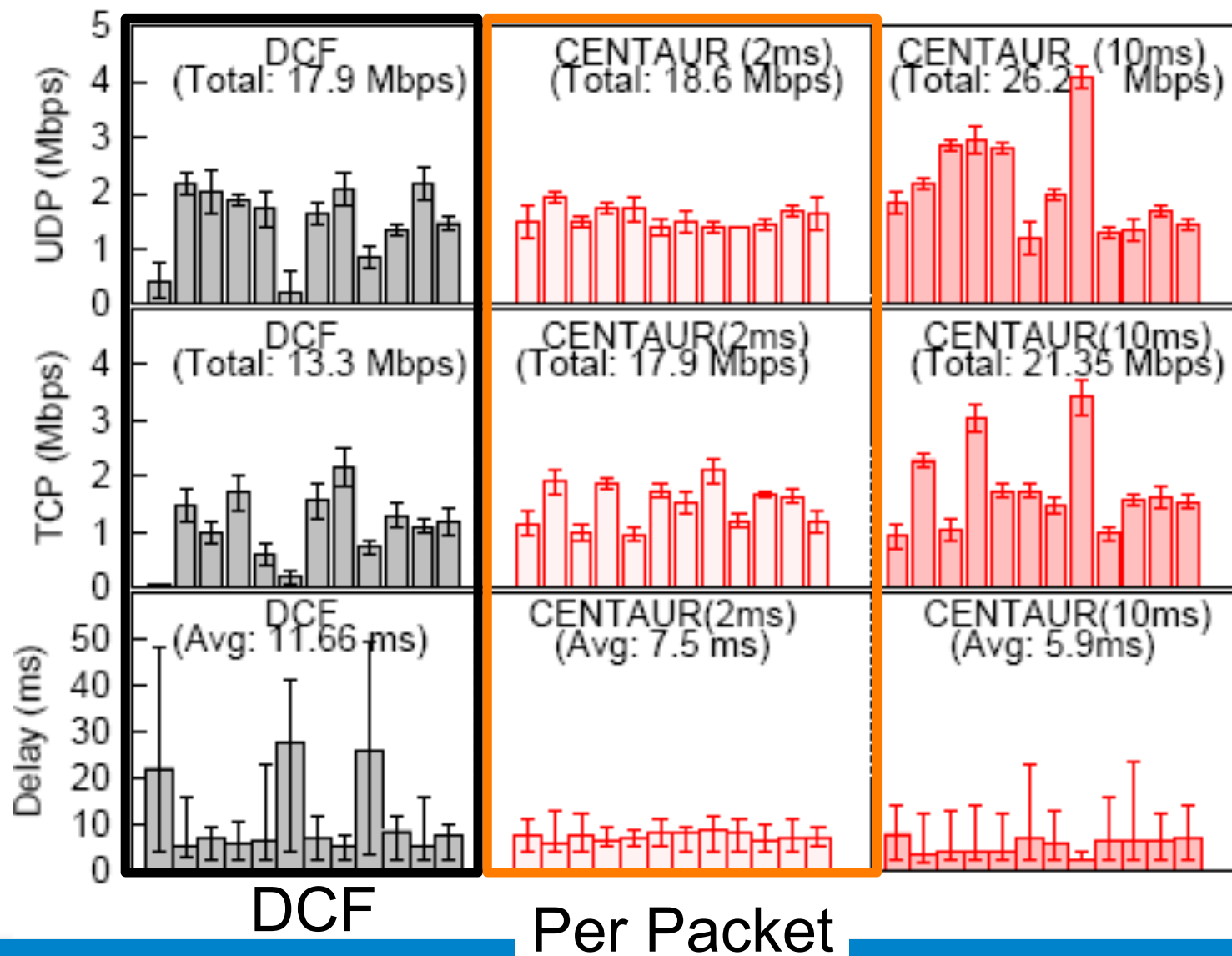
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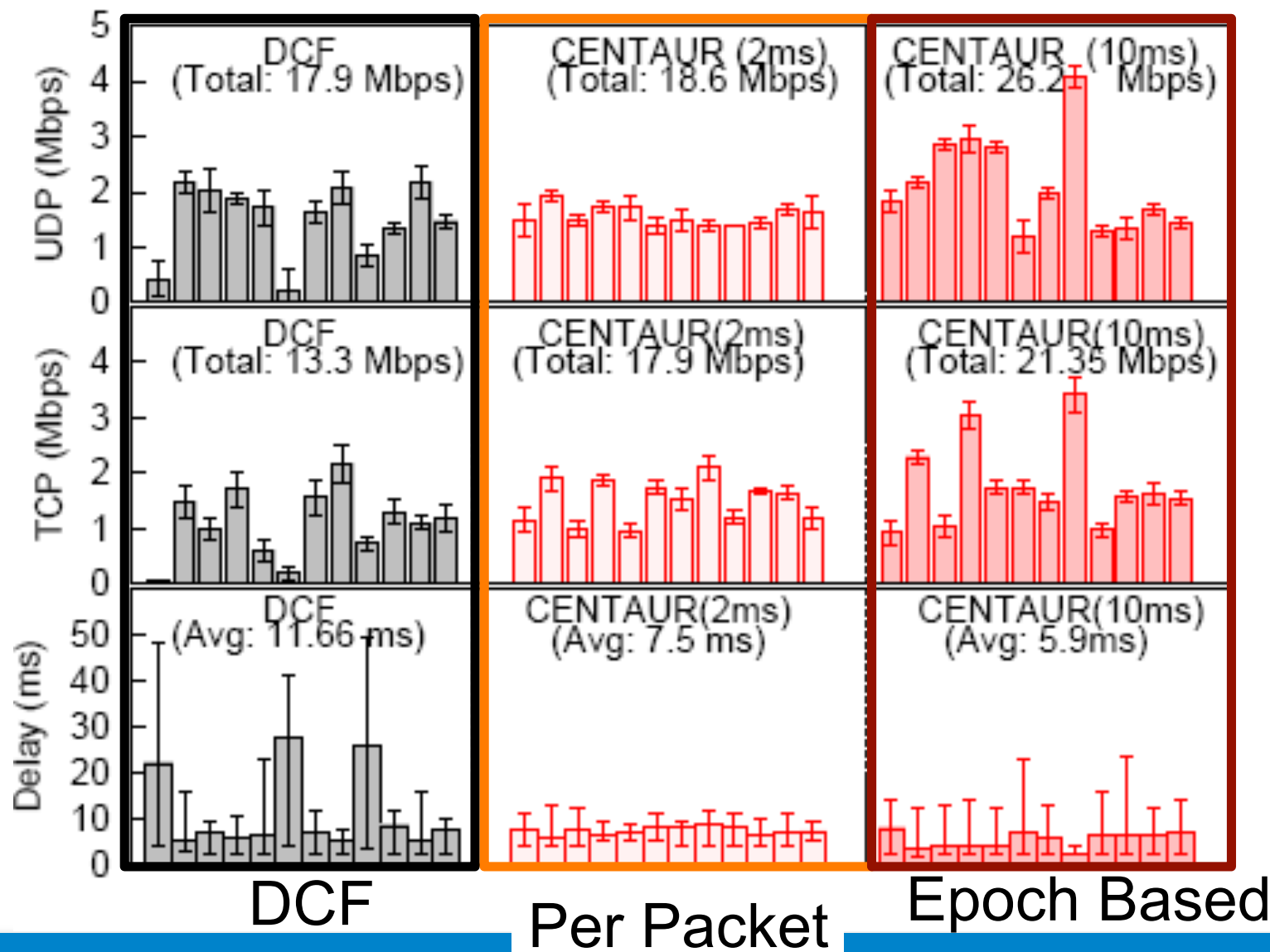
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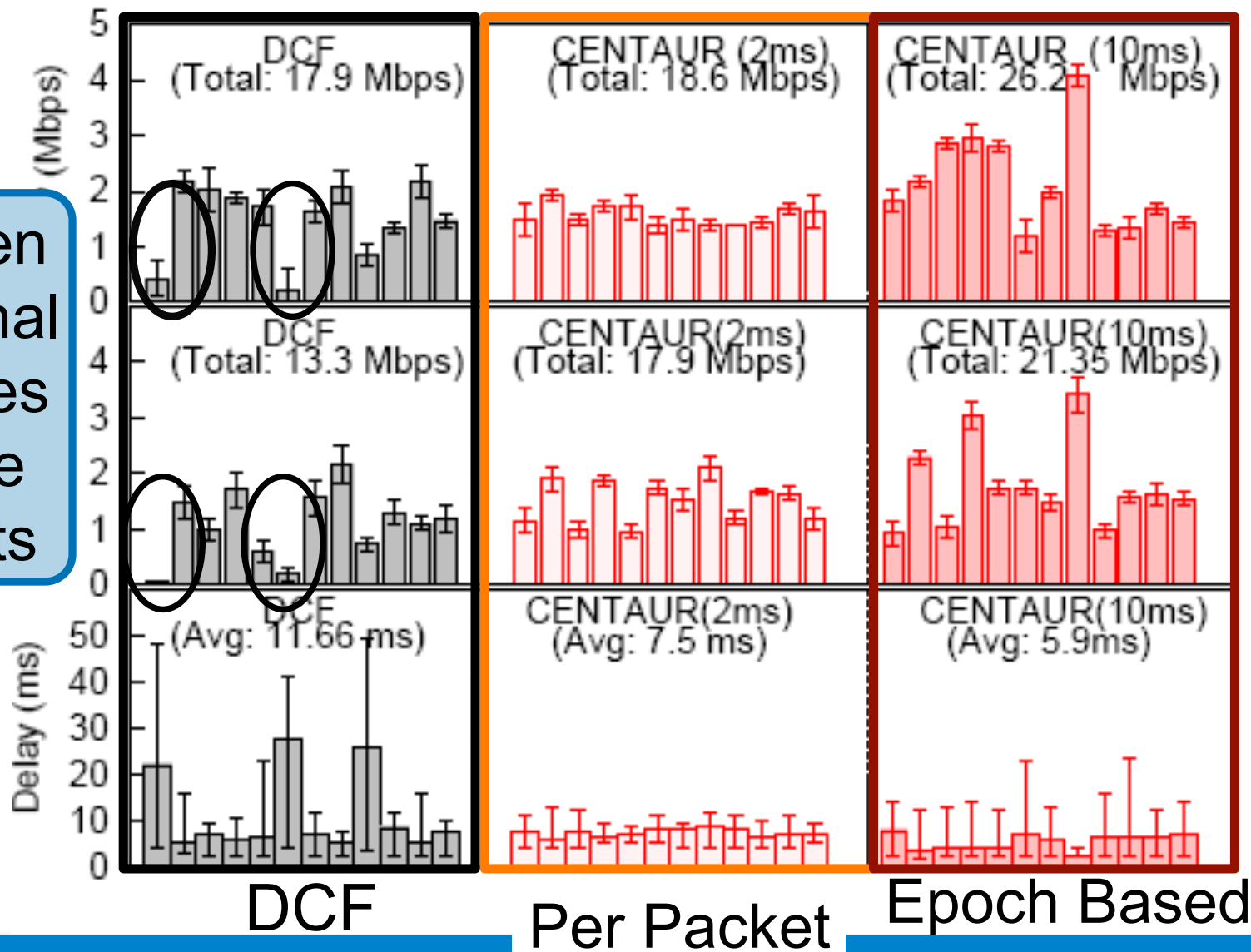


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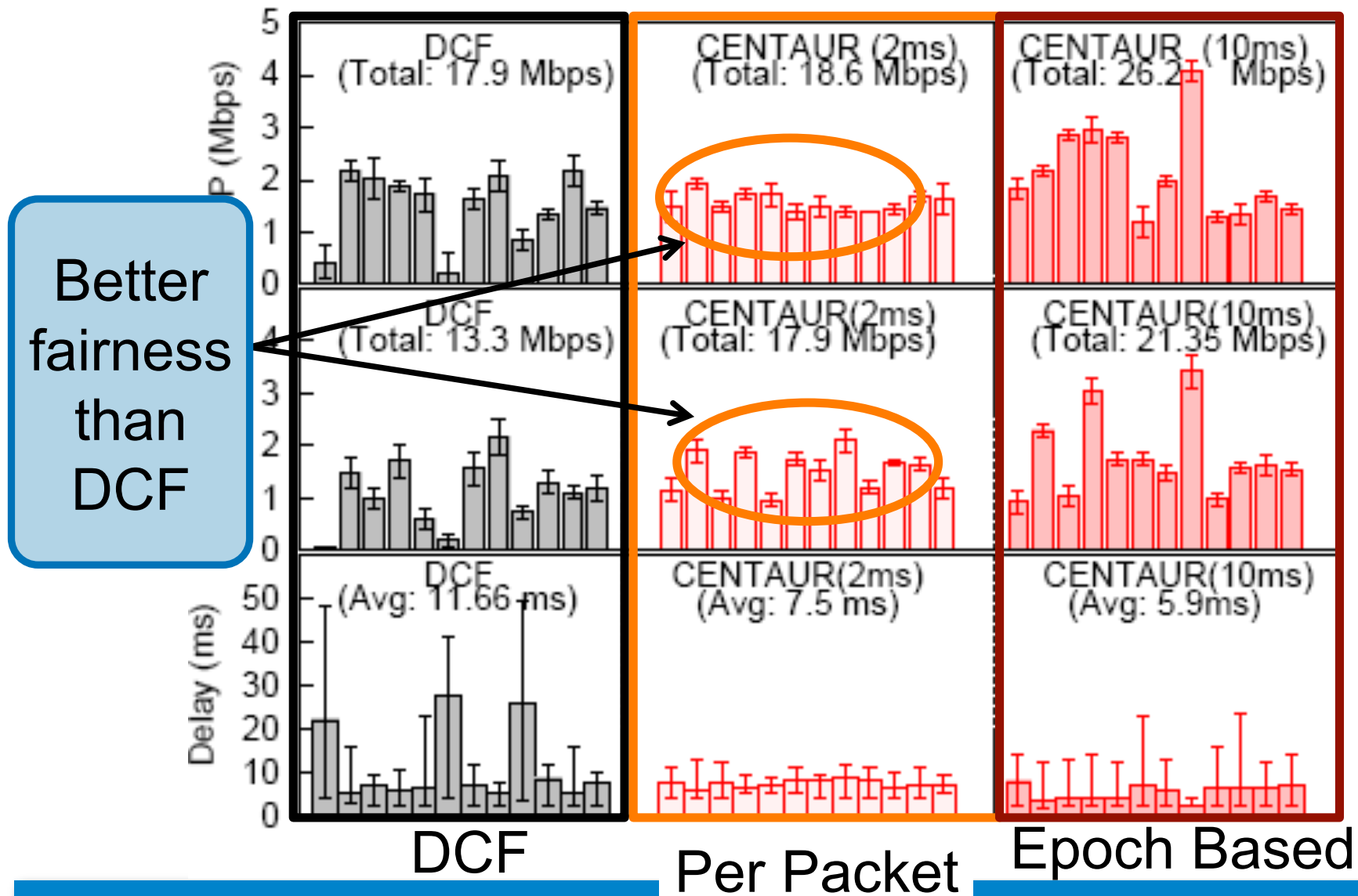


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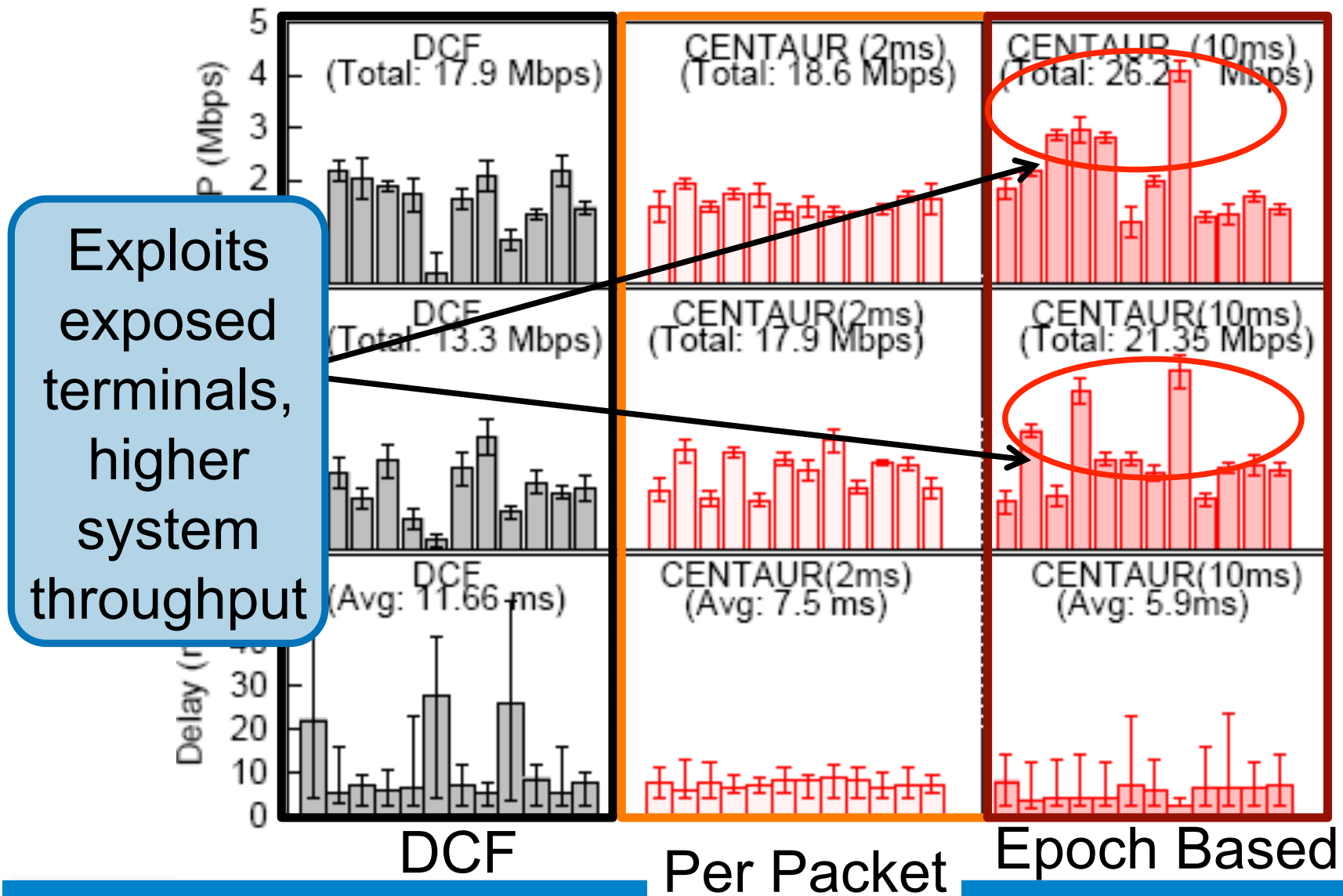
Hidden
terminal
starves
some
clients



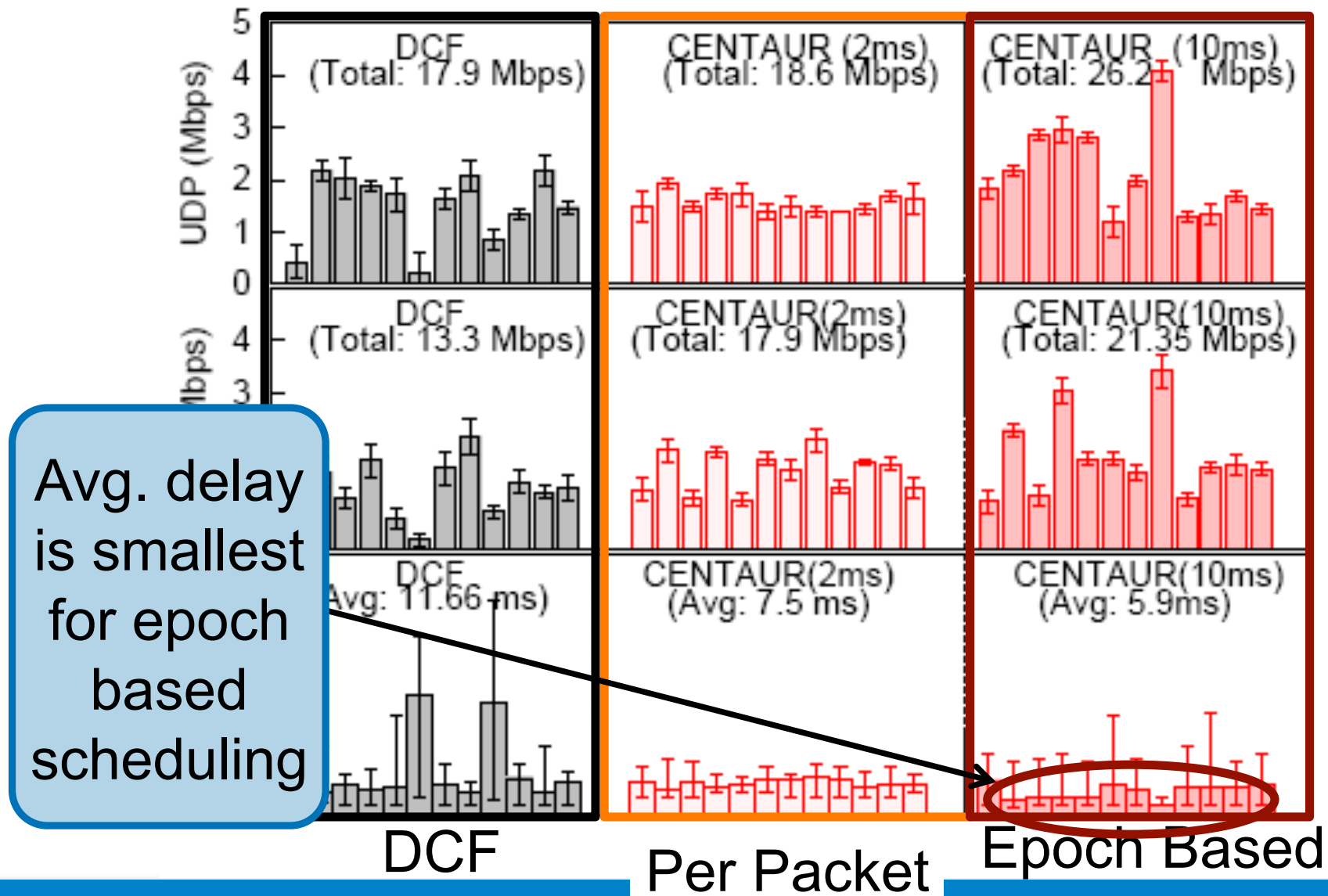
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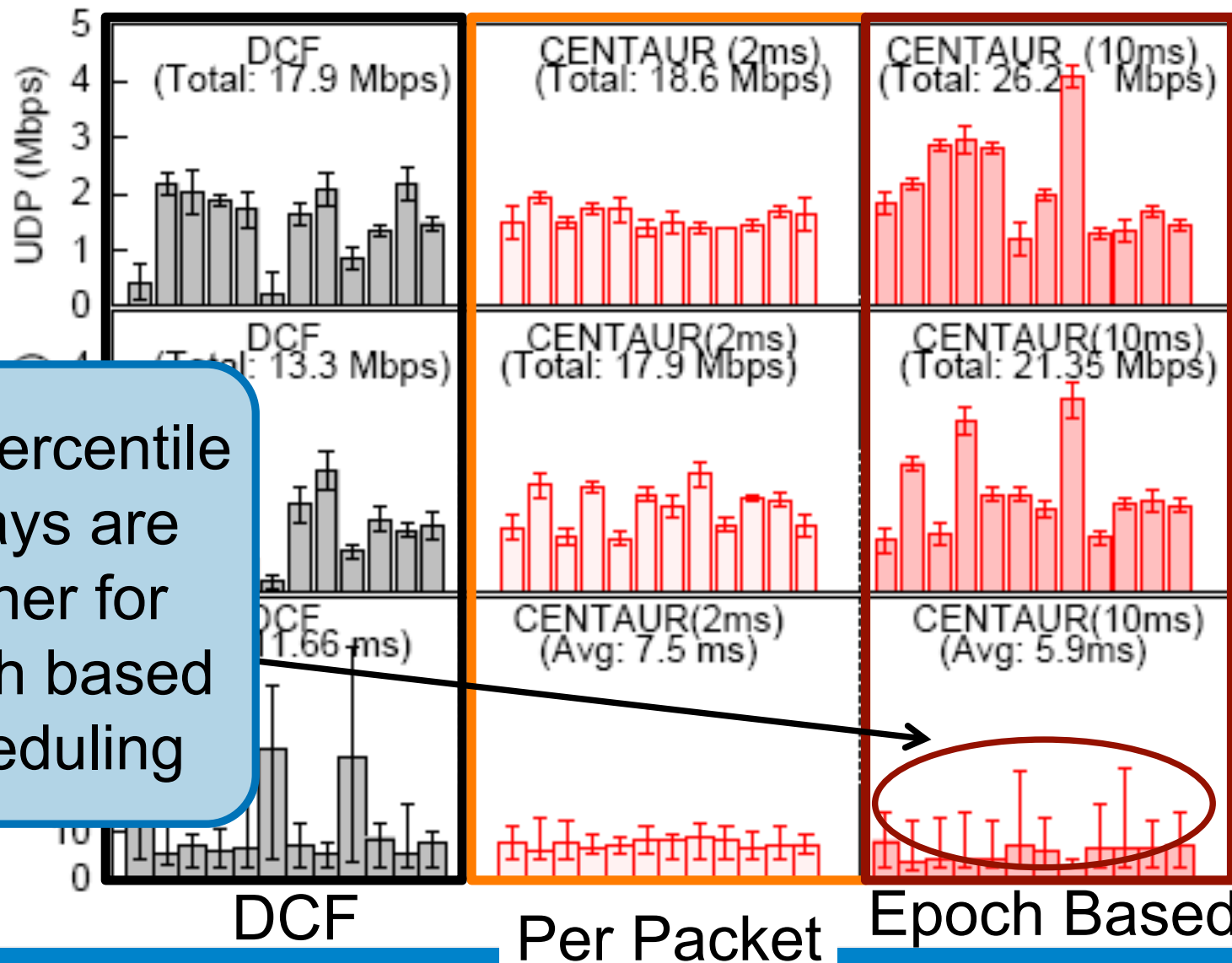
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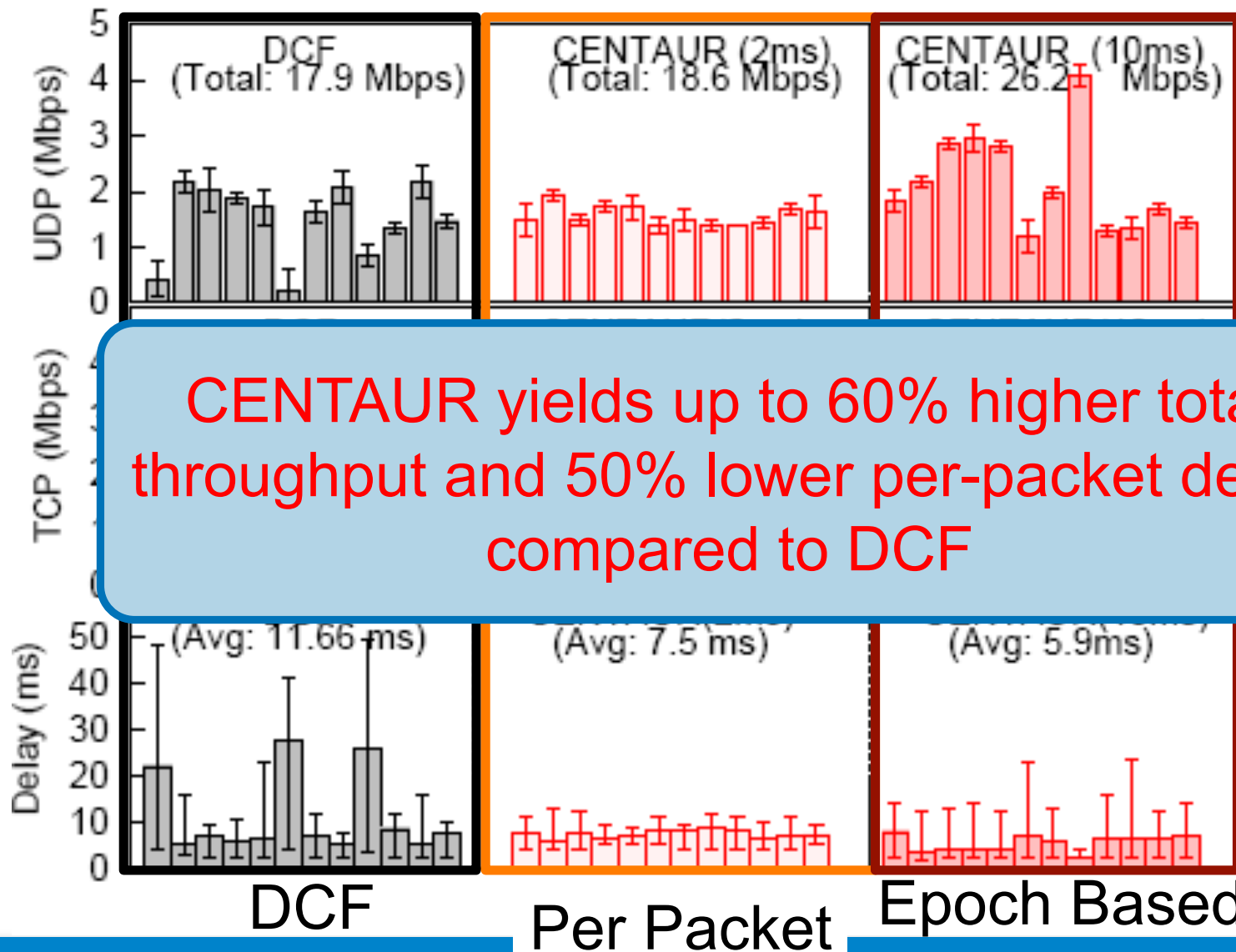
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Result 1: UDP/TCP Performance



Result 1: UDP/TCP Performance



More Results in Paper

- **Centaur micro-benchmarks:** performance for exposed and hidden terminals under centaur
- **Uplink traffic:** coexistence and persistent gains with different fractions of uplink traffic
- **Data rate:** robust to changes in data rate and ARF
- **Realistic HTTP traces:** significant reduction in web transaction delay
- **VoIP traffic:** better performance (*MOS*) for voice traffic with small epoch duration

Related Work

- Commercial WLAN offerings (Aruba, Meru)
- Theoretical formulations (Vaidya '00, Kanodia '01)
- Epoch based scheduling (Kompella '05, 802.11n/e)
- Interference mitigation (CMAP, SIC, Shuffle)

Summary

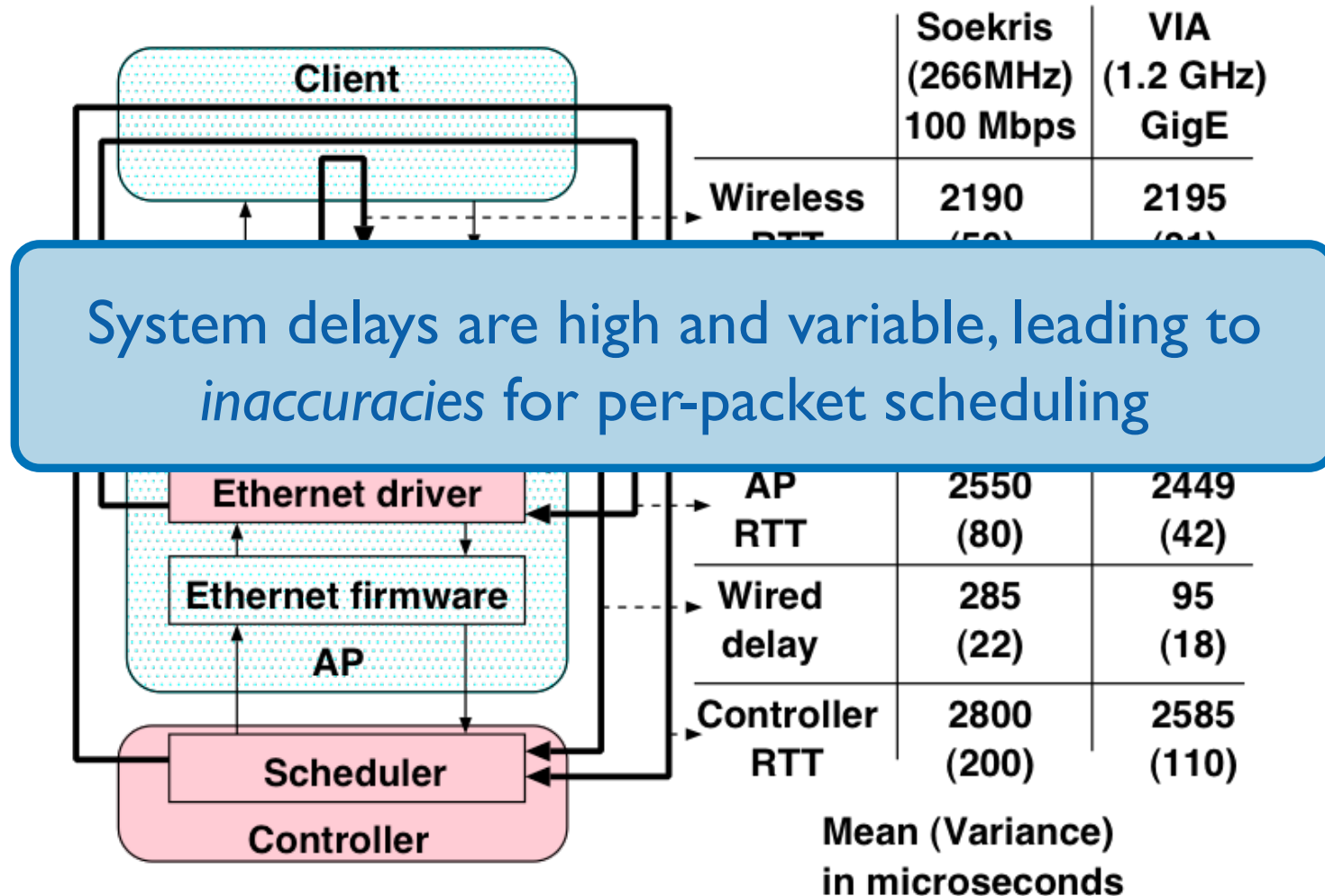
- Interference a growing problem in enterprises
- Careful design of a centralized data plane provides substantial performance gains
- CENTAUR implements a hybrid data path to improve aggregate performance without client modifications
- CENTAUR does not disable carrier sense and co-exists with non-enterprise and uplink traffic

Future work

- Even more efficient conflict graph generation
- What if we were allowed client modifications ?

Questions ?

Characterizing System Latencies



Evaluation of Micro-Probing

Topologies	20 node	30 node
Bandwidth Tests	16.2 mins	1hr 11 mins
Micro-Probing	~4 secs	~11 secs

Can be computed in stages
with each instance taking ~
2.5ms

Result 2: Impact of Uplink Traffic

- Vary proportion of downlink/uplink traffic
- 6 different configurations
- 80/20 → 40/60 (downlink/uplink)
- Results:
- Downlink: 1.6x → 6.8x gain in throughput
- Uplink: 1x → 1.18x gain in throughput

CENTAUR provides persistent gains for different proportions of uplink and downlink traffic load

Result 3: Impact of Topology

- **Three** topologies
- *Hidden Heavy* topology → 10 links
- *Exposed Heavy* topology → 6 links
- *Mixed Topology* → 12 links
- **Results:** Up to 50% gain in overall system throughput
- Up to **6x** gain for HT; Up to **1.7x** gain for ET

Improvements from using CENTAUR can be seen across many different network topologies

Result 2: Impact of Uplink Traffic

Downlink load (Mbps)	Uplink load (Mbps)	Downlink Throughput			Uplink Throughput median
		10%	50%	90%	
6	1.2	6.78×	1.48×	1.78×	1.15×
6	2.4	3.17×	1.37×	1.75×	1.04×
6	6	2.24×	1.21×	1.53×	1.01×
2.4	1.2	1.05×	1×	1×	1×
2.4	2.4	1.32×	1.11×	1.27×	1.06×
2.4	6	1.68×	1.21×	1.49×	1.18×

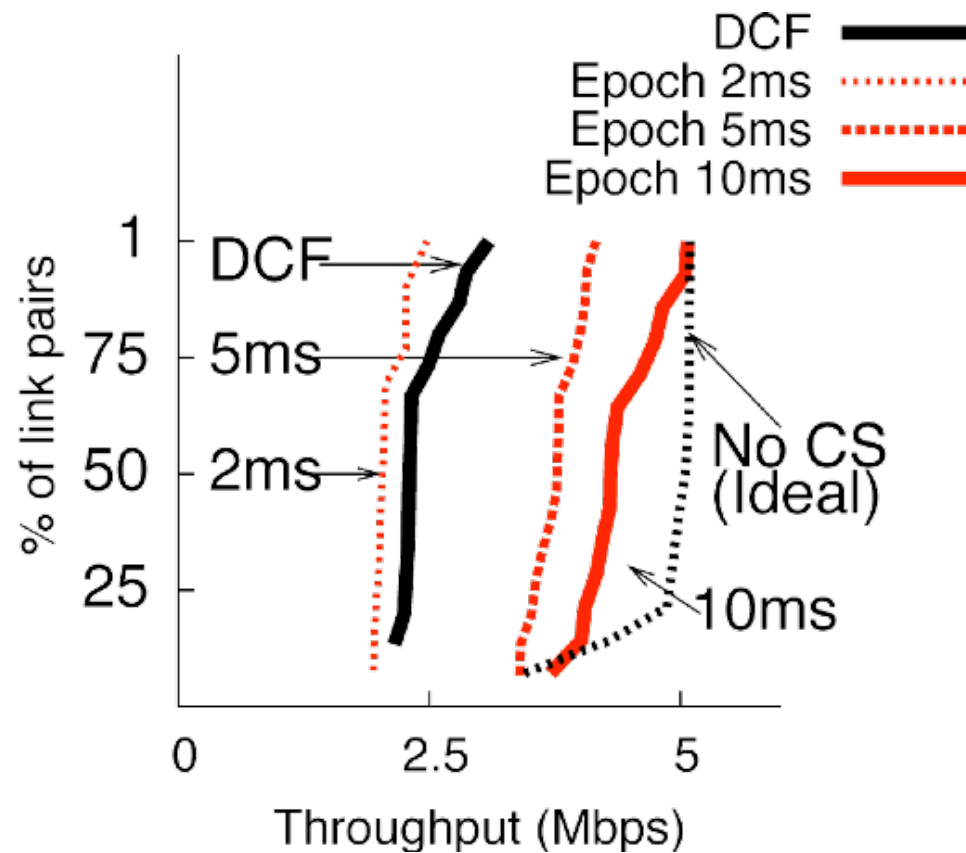
CENTAUR provides persistent gains for different proportions of uplink and downlink traffic load

Result 3: Impact of Topology

	Hidden-heavy (Testbed 2)	Exposed-heavy (Testbed 1)	Mixed (Testbed 1)
% HT	14%	0%	6.7%
% ET	0%	22%	10.2%
Overall Gains (HT/ET gains)	34.7% (HT: 6×)	47.2% (ET: 1.7×)	44% (HT: 3.2×, ET: 1.4×)

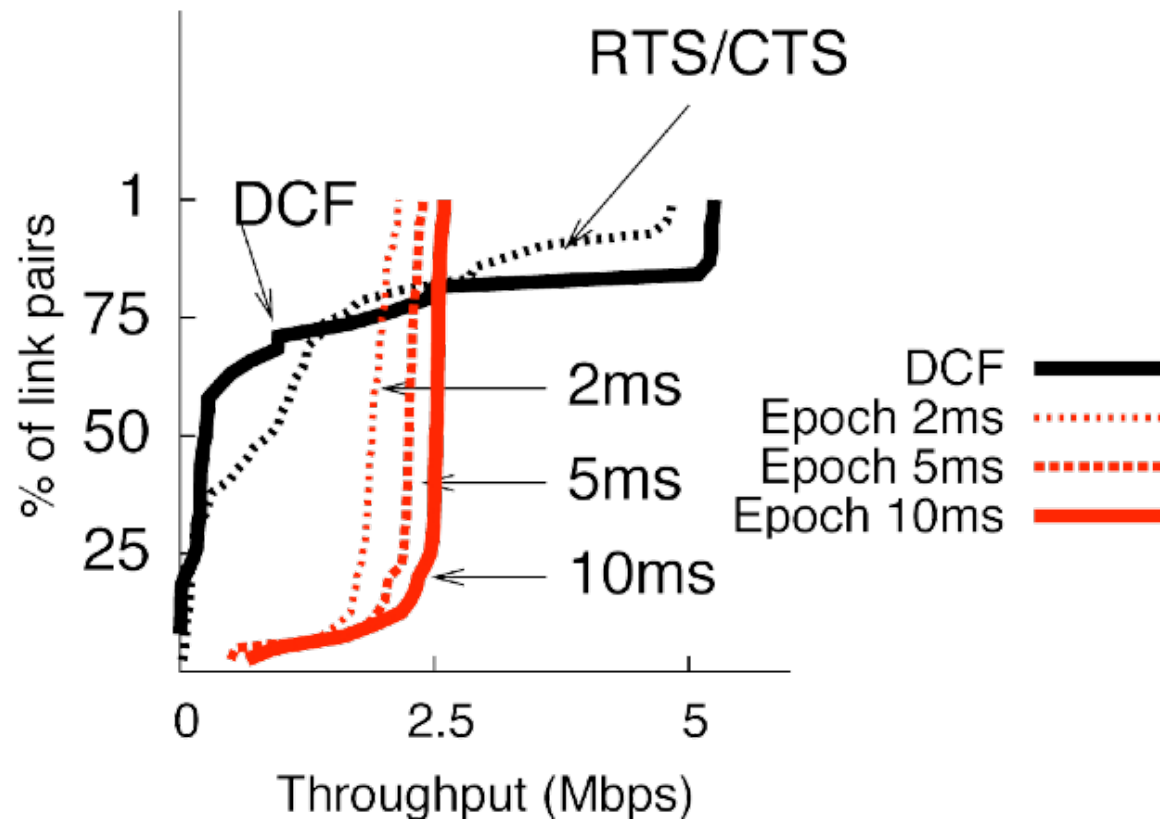
Improvements from using CENTAUR can be seen across many different network topologies

CENTAUR Micro-Benchmarks (I)



Exposed Terminals

CENTAUR Micro-Benchmarks (II)



Hidden Terminals

Result I: UDP/TCP Performance

Delay (ms)

CENTAUR yields up to 60% higher total throughput and 50% lower per-packet delay compared to DCF

Client Index (1-12)