

Midterm Examination

CS 525 - Fall 2010

Thursday, October 28, 2010, 7:15-9:15pm

Each question is worth the same number of points.

No electronic devices, notes, or books allowed, except that you may bring one standard-size sheet of paper, handwritten on both sides, into the test. **You need to give reasoning and justify all your answers**, citing the appropriate theorems where necessary.

1. (a) For the following choice of A and b , solve the system of equations $Ax = b$ by using tableaus. If there are multiple solutions, describe the full solution set.

8

$$A = \begin{bmatrix} 1 & 3 & 4 \\ 2 & 6 & -1 \\ 0 & 2 & 2 \end{bmatrix}, \quad b = \begin{bmatrix} 7 \\ -4 \\ 2 \end{bmatrix}.$$

- (b) Using Jordan exchanges, find the inverse of this matrix. If the matrix has no inverse, find a dependency relationship between its rows.

12

$$C = \begin{bmatrix} 0 & 1 & -2 \\ 2 & 3 & -1 \\ 1 & -1 & 3 \end{bmatrix}.$$

2. Consider the following linear program:

$$\begin{array}{ll} \min & 2x_1 + x_2 \\ & x_1 - x_2 \geq 4, \\ \text{subject to} & 3x_1 + 2x_2 \geq 10, \\ & x_1, x_2 \geq 0. \end{array}$$

- 6 (a) Write down the dual of this problem.
 6 (b) Find solutions for the primal and dual.
 8 (c) Suppose the right-hand side of the first constraint is changed from 4 to 10. Without performing any additional simplex iterations or referring to the tableau, give a lower bound on the optimal primal objective for the modified problem. Explain.

3. Solve the following linear program, using any of techniques we have learnt about in class (or any combination thereof). If it is infeasible, say so. If it is unbounded, give a direction of unboundedness. If there are multiple solutions, describe the full set of solutions.

schene II: 5 min $6x_1 + 6x_2 + 2x_3$
 dual simplex: 5 $4x_1 - 5x_2 + x_3 \geq 10,$
 solution: 4 subject to $3x_1 - 3x_2 + x_3 = 5,$
 multiple solutions: 6 $x_1, x_2 \geq 0, x_3 \text{ free.}$

4. Consider the following linear program:

$$\begin{aligned} \min_{x_1, x_2, \dots, x_n} & -x_1 - x_2 - x_3 - \dots - x_n \\ \text{subject to} & a_1x_1 + a_2x_2 + a_3x_3 + \dots + a_nx_n \leq 1, \\ & x_1, x_2, x_3, \dots, x_n \geq 0. \end{aligned}$$

where a_1, a_2, \dots, a_n are (strictly) positive constants.

- 4 (a) Write down the dual of this problem.
 4 (b) Find the solution of the dual, by inspection. Is this solution unique?
 4 (c) Write down the KKT conditions for this problem.
 4 (d) By using the KKT conditions (or by any other means), find a solution of the original (primal) linear program. (It is sufficient to write down just one solution.)
 4 (e) How do the primal and dual solutions change if one or more of the constants $a_i, i = 1, 2, \dots, n$ is allowed to be zero?

1 (a).

$$\begin{array}{l}
 y_1 = \\
 y_2 = \\
 y_3 =
 \end{array}
 \begin{array}{|ccc|c}
 x_1 & x_2 & x_3 & 1 \\
 \hline
 1 & 3 & 4 & -7 \\
 2 & 6 & -1 & 4 \\
 0 & 2 & 2 & -2
 \end{array}$$

→

$$\begin{array}{l}
 x_1 = \\
 y_2 = \\
 y_3 =
 \end{array}
 \begin{array}{|ccc|c}
 y_1 & x_2 & x_3 & 1 \\
 \hline
 1 & -3 & -4 & 7 \\
 2 & 0 & -9 & 18 \\
 0 & 2 & 2 & -2
 \end{array}$$

↓

$$\begin{array}{l}
 x_1 = \\
 y_2 = \\
 x_3 =
 \end{array}
 \begin{array}{|ccc|c}
 y_1 & x_2 & y_3 & 1 \\
 \hline
 1 & 1 & -2 & 3 \\
 2 & 9 & -\frac{9}{2} & 9 \\
 0 & -1 & \frac{1}{2} & 1
 \end{array}$$

←

$$\begin{array}{l}
 x_1 = \\
 x_2 = \\
 x_3 =
 \end{array}
 \begin{array}{|ccc|c}
 y_1 & y_2 & y_3 & 1 \\
 \hline
 & \frac{1}{9} & & 2 \\
 -\frac{2}{9} & \frac{1}{9} & \frac{1}{2} & -1 \\
 & -\frac{1}{9} & & 2
 \end{array}$$

solution $x = \begin{pmatrix} 2 \\ -1 \\ 2 \end{pmatrix}$

(b)

$$\begin{array}{l}
 y_1 = \\
 y_2 = \\
 y_3 =
 \end{array}
 \begin{array}{|ccc|}
 x_1 & x_2 & x_3 \\
 \hline
 0 & 1 & -2 \\
 2 & 3 & -1 \\
 1 & -1 & 3
 \end{array}$$

→

$$\begin{array}{l}
 x_2 = \\
 y_2 = \\
 y_3 =
 \end{array}
 \begin{array}{|ccc|}
 x_1 & y_1 & x_3 \\
 \hline
 0 & 1 & 2 \\
 2 & 3 & 5 \\
 1 & -1 & 1
 \end{array}$$

↓

$$\begin{array}{l}
 x_2 = \\
 y_2 = \\
 x_1 =
 \end{array}
 \begin{array}{|ccc|}
 y_3 & y_1 & x_3 \\
 \hline
 0 & 1 & 2 \\
 2 & 5 & 3 \\
 1 & 1 & -1
 \end{array}$$

←

$$\begin{array}{l}
 x_2 = \\
 x_3 = \\
 x_1 =
 \end{array}
 \begin{array}{|ccc|}
 y_3 & y_1 & y_2 \\
 \hline
 \frac{4}{3} & -\frac{1}{3} & \frac{2}{3} \\
 \frac{2}{3} & -\frac{5}{3} & \frac{1}{3} \\
 \frac{7}{3} & \frac{8}{3} & -\frac{1}{3}
 \end{array}$$

↓ permute columns

$$\begin{array}{l}
 x_2 = \\
 x_3 = \\
 x_1 =
 \end{array}
 \begin{array}{|ccc|}
 y_1 & y_2 & y_3 \\
 \hline
 -\frac{1}{3} & \frac{7}{3} & -\frac{4}{3} \\
 -\frac{5}{3} & \frac{1}{3} & -\frac{1}{3} \\
 \frac{8}{3} & -\frac{1}{3} & \frac{5}{3}
 \end{array}$$

→ permute rows

$$\begin{array}{l}
 x_1 = \\
 x_2 = \\
 x_3 =
 \end{array}
 \begin{array}{|ccc|}
 y_1 & y_2 & y_3 \\
 \hline
 \frac{8}{3} & -\frac{1}{3} & \frac{5}{3} \\
 -\frac{1}{3} & \frac{7}{3} & -\frac{4}{3} \\
 -\frac{5}{3} & \frac{1}{3} & -\frac{1}{3}
 \end{array}$$

$$C^{-1} = \begin{pmatrix} \frac{8}{3} & -\frac{1}{3} & \frac{5}{3} \\ -\frac{1}{3} & \frac{7}{3} & -\frac{4}{3} \\ -\frac{5}{3} & \frac{1}{3} & -\frac{1}{3} \end{pmatrix}$$

2 (a) $\max 4u_1 + 10u_2$
 s.t. $u_1 + 3u_2 \leq 2$
 $-u_1 + 2u_2 \leq 1$
 $u_1, u_2 \geq 0$

(b)

		$u_3 =$	$u_4 =$	$w =$
		x_1	x_2	1
$-u_1$	$x_3 =$	①	-1	-4
$-u_2$	$x_4 =$	3	2	-10
1	$z =$	2	1	0

↓ dual simplex
 chooses (1,1) first

		$u_1 =$	$u_4 =$	$w =$
		x_3	x_2	1
$-u_3$	$x_1 =$	1	1	4
$-u_2$	$x_4 =$	3	5	2
1	$z =$	2	3	8

solution: primal $x^* = \begin{pmatrix} 4 \\ 0 \end{pmatrix}$ dual $u^* = \begin{pmatrix} 2 \\ 0 \end{pmatrix}$

(c). changing b_1 from 4 to 10 does not change dual constraints, so u^* is feasible for the modified dual, with objective $10u_1^* + 10u_2^* = 20$.
 By weak duality, this is a lower bound on the optimal objective of the modified primal

(d) ~~x^* is still feasible for the modified primal, and u^* is still feasible for the modified dual.
 The modified primal objective at x^* is still 8
 The modified dual objective at u^* is still 8
 Hence by strong duality, x^* solves modified primal and u^* solves modified dual.~~

	x_1	x_2	x_3	1
$x_4 =$	4	-5	1	-10
$x_5 =$	3	-3	1	-5
$z =$	6	6	2	0

Scheme II

	x_1	x_2	x_3	1
$x_4 =$	1	-2	1	-5
$x_3 =$	-3	3	1	5
$z =$	0	12	2	10

remove $x_5 = 0$
move x_3 to bottom

	x_1	x_2	1
$x_4 =$	1	-2	-5
$z =$	0	12	10
$x_3 =$	-3	-3	5

dual simplify
pivot on (1,1)

	x_1	x_2	1
$x_1 =$	1	2	5
$z =$	0	12	10
$x_3 =$	-3	-3	-10

Solution: $x = \begin{pmatrix} 5 \\ 0 \\ -10 \end{pmatrix}$

non unique can get additional solutions by letting $x_4 = \alpha \geq 0$

$$x = \begin{pmatrix} 5 \\ 0 \\ -10 \end{pmatrix} + \begin{pmatrix} 1 \\ 0 \\ -3 \end{pmatrix} \alpha$$

4. (a) First change to standard form:

$$\begin{aligned} \min \quad & -x_1 - x_2 \dots - x_n \\ \text{s.t.} \quad & -a_1 x_1 - a_2 x_2 \dots - a_n x_n \geq -1 \\ & x_1, x_2, \dots, x_n \geq 0 \end{aligned}$$

Dual is

$$\begin{aligned} \max \quad & u \\ \text{s.t.} \quad & -a_i u \leq -1, \quad i=1, \dots, n \\ & u \geq 0 \end{aligned}$$

or equivalently,

$$\begin{aligned} \min \quad & u \\ \text{s.t.} \quad & a_i u \geq 1 \quad (i=1, \dots, n) \\ & u \geq 0 \end{aligned}$$

(b) solution $u^* = \frac{1}{\min_{i=1, \dots, n} (a_i)}$

yes, it is unique

(c) KKT: $0 \leq x_i \perp -1 + a_i u \geq 0 \quad (i=1, \dots, n)$
 $0 \leq 1 - a_1 x_1 - \dots - a_n x_n \perp u \geq 0$

(d) Let $\mathcal{J} = \{j \mid j \in \arg \min_i a_i\}$, that is, \mathcal{J} is the set of indices j such that a_j is a minimum of $\{a_1, a_2, \dots, a_n\}$

then for $j \notin \mathcal{J}$ we have $-1 + a_j u > 0 \Rightarrow x_j^* = 0$ while for $j \in \mathcal{J}$ we may have $x_j^* > 0$.

Since $u^* > 0$, we have $\sum_{i \in \mathcal{J}} a_i x_i^* = \sum_{i \in \mathcal{J}} a_i x_i^* = 1$.

these conditions together define x^* .

One such solution would be: for some $j \in \mathcal{J}$,

$$\begin{aligned} x_j^* &= \frac{1}{a_j} \\ x_i^* &= 0, \quad (i \neq j) \end{aligned}$$

(e) if any $a_i = 0$, the dual is infeasible, while the primal is clearly still feasible ($x=0$ is a feasible point). Hence, by strong duality the primal is unbounded. (We can also see this by inspection: for any i such that $a_i = 0$, we can let $x_i \uparrow \infty$ without affecting feasibility, while driving the objective to $-\infty$.)