Symbiosis: The Art of Application and Kernel Cache Cooperation

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Database Systems and Storage Engines

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Large scale database systems MongoDB / MySQL / CockroachDB



Database Systems and Storage Engines

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Large scale database systems MongoDB / MySQL / CockroachDB Storage engines within DB to interact with filesystems

RocksDB / WiredTiger / InnoDB



Caching in Storage Engine

Cache within the application An implicit layer of kernel caching





Caching in Storage Engine

Cache within the application An implicit layer of kernel caching Stores compressed data





Caching in Storage Engine

Forms a special two-layer cache structure Sharing of memory quota
Optimal cache partitioning is important and not trivial Yields great benefit
Depends on various factors



Optimal Cache Partitioning is Hard





Performance with **small** working set size

Performance with **large** working set size

Optimal Cache Partitioning is Hard





large working set size

Optimal Cache Partitioning is Hard





The best partition could be arbitrary percentage in the middle



Symbiosis





Symbiosis

Embedded in the storage engine





Symbiosis

- Embedded in the storage engine
- Optimizes cache partitioning automatically





Symbiosis

Embedded in the storage engine Optimizes cache partitioning automatically Adapts cache sizes to dynamic workloads





Integrated into production systems with <1000 LoC LevelDB, RocksDB, WiredTiger

Performance improvements

1.5x on average for read-heavy workloads

Online cache simulation with high accuracy and negligible overhead ~0.1% space overhead and ~1% time overhead



Outline



Overview of Symbiosis

Key Challenge

Simulate accurately with low overhead

Optimization Techniques

Incremental reuse of a single ghost cache

Misalignment-aware sampling

Guard against unmodeled

Evaluation

Static workloads

Dynamic (changing) workloads

Symbiosis - Overview



Detects workload change

Symbiosis - Overview



Detects workload change

Simulates multiple candidate cache size configurations





Detects workload change

Simulates multiple candidate cache size configurations

Finds the configuration with the best expected performance





Detects workload change

Simulates multiple candidate cache size configurations

Finds the configuration with the best expected performance

Applies the best configuration if it yields enough benefit



Key Challenge



How to simulate accurately?

How to simulate with negligible overhead?

Key Challenge



How to simulate accurately? Tension! How to simulate with negligible overhead?

Key Challenge



How to simulate accurately? Tension! How to simulate with negligible overhead?

Optimization Techniques Incremental reuse of a single ghost cache Misalignment-aware sampling Guard against unmodeled

Classical Solution - Ghost Cache Simulation



Ghost cache - maintain only cache access metadata Useful for cache statistics analysis



Ghost Cache Simulation - Single Layer



Use a single ghost cache instance Largely reduces space and time overhead Simulate caches of multiple sizes simultaneously When the eviction policy has stack property





Simultaneous simulation of caches with different sizes is not feasible The second layer sees different access patterns



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2 1 3 2 3 1 3 1 2 4 1st-layer access trace: 2 1 3 2 3 1 2 1 3 4 2 1 3 2 1 3 1 3 2 4 size = 1

1st-layer Ghost cache





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2 1 3 2 3 1 3 1 2 4 1st-layer access trace: 2 1 3 2 3 1 2 1 3 4 2 1 3 2 1 3 1 3 2 4 2 1 3 2 3 1 3 1 3 2 4 2 1 3 2 3 1 3 1 2 4 2 1 3 2 3 1 3 1 2 4 2 1 3 2 3 1 2 1 3 4 2 1 3 2 3 1 2 1 3 4 2 1 3 2 3 1 2 1 3 4 2 1 3 2 1 3 1 3 2 4 size = 1

1st-layer Ghost cache





Simultaneous simulation of caches with different sizes is not feasible The second layer sees different access patterns

size = 3 1st-layer access trace: 2132312134 2132131324 1st-layer Ghost cache 1 2 3

2nd-layer access trace: 4 3 4 2 4

Incremental reuse of a single ghost cache



Use a single ghost cache instance Largely reduces space and time overhead

Incremental reuse of a single ghost cache



Use a single ghost cache instance Largely reduces space and time overhead

Simulate candidates one at a time



Full app cache (for reuse)

Candidate Configuration 1

Candidate Configuration 2

Candidate Configuration 3

Incremental reuse of a single ghost cache



Use a single ghost cache instance Largely reduces space and time overhead

Simulate candidates one at a time





Use a single ghost cache instance Largely reduces space and time overhead Simulate candidates one at a time Minimize warm-up before each simulation Utilize the stack property for the first layer Simulate in the order of increasing application cache size





Use a single ghost cache instance Largely reduces space and time overhead Simulate candidates one at a time Minimize warm-up before each simulation Utilize the stack property for the first layer Simulate in the order of increasing application cache size





Key-space sampling Sample targets in key space and all accesses to the targets





Key-space sampling Samples targets in key space and all accesses to the targets Works in normal multi-layer caching





Misalignment - different caching units The unit of app cache: app-defined blocks The unit of kernel cache: pages One unit in the 1st layer corresponds to multiple units in the 2nd layer





Key-space sampling with misalignment





Key-space sampling with misalignment Hits in the 2nd layer due to accessed neighboring blocks in the 1st layer





Key-space sampling with misalignment Hits in the 2nd layer due to accessed neighboring blocks in the 1st layer Contiguous targets not all sampled so hits are not counted Loses the effects of accessing contiguous pages



Key-space sampling with misalignment Much lower kernel cache hit rate



Waldspurger et al. Efficient MRC Construction with SHARDS. 2015 (original key-space sampling)



Misalignment-aware Sampling

Sample contiguous targets together to keep the effect of misalignment





Misalignment-aware Sampling

Sample contiguous targets together to keep the effect of misalignment In Symbiosis, we samples in groups of 32 and the sampling rate is 1/64 Misalignment-aware sampling is much more accurate





Some kernel features not modeled (e.g., read-ahead)

Some kernel features not modeled (e.g., read-ahead) When simulation is accurate





Some kernel features not modeled (e.g., read-ahead) When simulation is not accurate Observation: omitting such features usually results in a worse perf.





Some kernel features not modeled (e.g., read-ahead) Observation: omitting such features usually results in a worse perf. Policy: only adapts when predicted perf. is better than current perf.



Evaluation - Static Workloads



Experiments

Performance with various workloads and environments

Factors		Presented Space	# of parameters
Workloads	Data Set Size	5, 2.5, 1.67, 1.24, 1	4
	Access Pattern	uniform, zipfian, hotspot{30,20,10}	5
Software	Compression Lib.	snappy (default), zstd	2
	Storage Engine	LevelDB, RocksDB, WiredTiger	3
Hardware	CPU Frequency	CPU1: 2.9GHz, CPU2: 2.0GHz	2
	Device Latency	SSD1: ~10us, SSD2: ~70us	2

Results

Improves performance in 98% of 190+ workloads by 1.5x on average

Evaluation - Dynamic Workloads



Experiments

Performance with abrupt and gradual workload changes

Performance with real world workloads

Overhead and tail latency analysis

Results

Symbiosis properly reacts to all of 38 workloads

Symbiosis incurs 0.1% space overhead and 1% time overhead



Consists of 2 static workloads Workload 1 is very skewed, while workload 2 is less skewed Data set fits in memory when compressed, but not when uncompressed





Workload 1 100% app cache performs better due to high skewedness



Workload 1 Symbiosis performs best after simulation

Workload 1



- 100% Kernel Cache
- 100% App Cache
- Symbiosis

Workload 1 Symbiosis incurs negligible overhead during simulation

Workload 1



- 100% Kernel Cache
- 100% App Cache
- Symbiosis

Workload 1



- 100% Kernel Cache
- 100% App Cache
- Symbiosis

Workload 1



- 100% Kernel Cache
- 100% App Cache
- Symbiosis

Workload 1 Simulation result: giving ~40% of memory to app cache is the best



- 100% Kernel Cache
- 100% App Cache
- Symbiosis

Workload 2 100% kernel cache performs similarly



Workload 2 100% app cache performs worse due to less skewedness



Workload 2 Symbiosis maintains the best performance after simulation



Workload 2 Symbiosis maintains the best performance after simulation



Workload 2 Symbiosis maintains the best performance after simulation



Conclusion



Symbiosis: dynamic cache size adjustment via online simulation

Integrated into production systems within 1000 LoC LevelDB, RocksDB, WiredTiger

- Adapts to various workloads and reacts to workload changes Performs well on 190+ static workloads and 38 dynamic workloads
- Delivers excellent performance with negligible overhead
 - 1.5x gain on average for read-heavy workloads
 - 0.1% space overhead and 1% time overhead

Conclusion



Simulation-based online performance tuning A tuning approach that does not rely on machine learning Deep understanding and careful optimizations are necessary Potential parameter tuning beyond cache sizes See the paper for: Offline simulation study Detailed implementation and evaluation

Symbiosis source code

https://github.com/daiyifandanny/Symbiosis

