



# CS 540 Introduction to Artificial Intelligence

## **Search III: Advanced Search**

Yudong Chen  
University of Wisconsin-Madison

**Nov 18, 2021**

# Announcements

- **Homeworks:**
  - HW8 due Tuesday.
  
- **Class roadmap:**
  - Today: Search III
  - Tuesday before and after Thanksgiving: Game (uses Search)

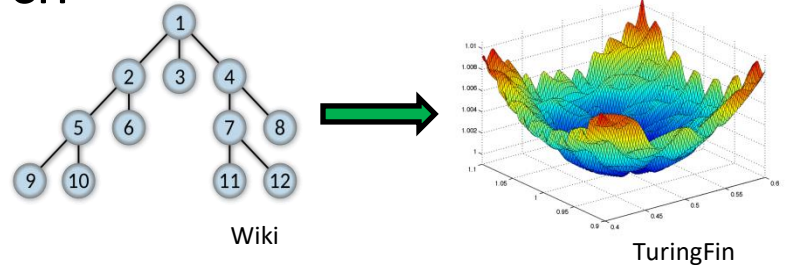
# Outline

- Advanced Search & Hill-climbing
  - More difficult problems, basics, local optima, variations
- Simulated Annealing
  - Basic algorithm, temperature, tradeoffs
- Genetic Algorithms
  - Basics of evolution, fitness, natural selection

# Search vs. Optimization

Before: wanted a **path** from start state to goal state

- Uninformed search, informed search



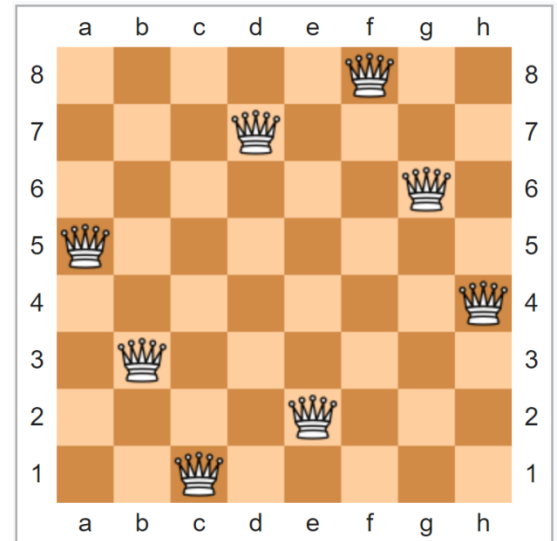
**New setting:** optimization

- States  $s$  have values  $f(s)$
- Want:  $s$  with optimal value  $f(s)$  (i.e, **optimize** over states)
- Challenging setting: **too many states** for previous search approaches, but maybe not a continuous function for SGD.

# Examples: $n$ Queens

A classic puzzle:

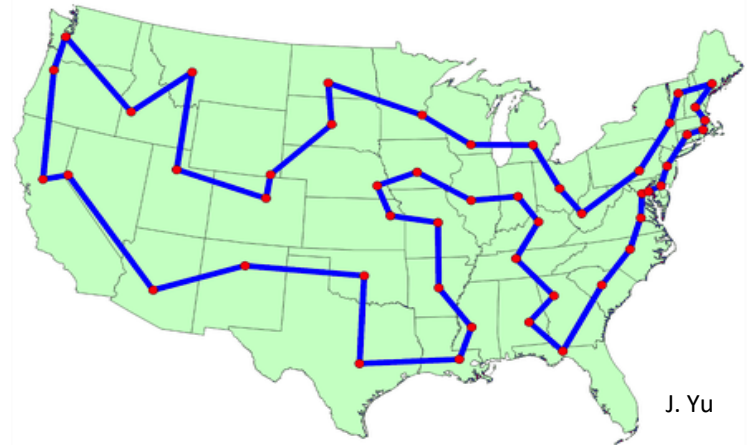
- Place 8 queens on a 8 x 8 chessboard so that no two have same row, column, or diagonal.
- Can generalize to  $n \times n$  chessboard.
- What are states  $s$ ? Values  $f(s)$ ?
  - State: configuration of the board
  - $f(s)$ : # of conflicting queens



# Examples: TSP

Famous graph theory problem.

- Get a graph  $G = (V, E)$ . **Goal:** a path that visits each node exactly once and returns to the initial node (a **tour**).
  - State: a particular tour (i.e., ordered list of nodes)
  - $f(s)$ : total weight of the tour (e.g., total miles traveled)



# Examples: Satisfiability

## Boolean satisfiability (e.g., 3-SAT)

- Recall our logic lecture. Conjunctive normal form

$$(A \vee \neg B \vee C) \wedge (\neg A \vee C \vee D) \wedge (B \vee D \vee \neg E) \wedge (\neg C \vee \neg D \vee \neg E) \wedge (\neg A \vee \neg C \vee E)$$

- Goal: find if satisfactory assignment exists.
- State: assignment to variables

–  $f(s)$ : # satisfied clauses

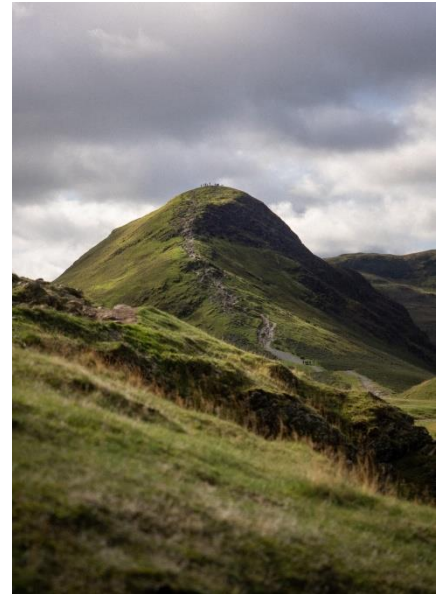
$$\begin{array}{l}
 R(x, a, d) \wedge R(y, b, d) \wedge R(a, b, e) \wedge R(c, d, f) \wedge R(z, c, 0) \\
 \hline
 R(0, a, d) \wedge R(0, b, d) \wedge R(a, b, e) \wedge R(c, d, f) \wedge R(0, c, 0) \\
 R(0, a, d) \wedge R(0, b, d) \wedge R(a, b, e) \wedge R(c, d, f) \wedge R(1, c, 0) \\
 R(0, a, d) \wedge R(1, b, d) \wedge R(a, b, e) \wedge R(c, d, f) \wedge R(0, c, 0) \\
 R(0, a, d) \wedge R(1, b, d) \wedge R(a, b, e) \wedge R(c, d, f) \wedge R(1, c, 0) \\
 R(1, a, d) \wedge R(0, b, d) \wedge R(a, b, e) \wedge R(c, d, f) \wedge R(0, c, 0) \\
 R(1, a, d) \wedge R(0, b, d) \wedge R(a, b, e) \wedge R(c, d, f) \wedge R(1, c, 0) \\
 R(1, a, d) \wedge R(1, b, d) \wedge R(a, b, e) \wedge R(c, d, f) \wedge R(0, c, 0) \\
 R(1, a, d) \wedge R(1, b, d) \wedge R(a, b, e) \wedge R(c, d, f) \wedge R(1, c, 0)
 \end{array}$$

$$\begin{array}{l}
 R(\neg x, a, b) \wedge R(b, y, c) \wedge R(c, d, \neg z) \\
 \hline
 R(1, a, b) \wedge R(b, 0, c) \wedge R(c, d, 1) \\
 R(1, a, b) \wedge R(b, 0, c) \wedge R(c, d, 0) \\
 R(1, a, b) \wedge R(b, 1, c) \wedge R(c, d, 1) \\
 R(1, a, b) \wedge R(b, 1, c) \wedge R(c, d, 0) \\
 R(0, a, b) \wedge R(b, 0, c) \wedge R(c, d, 1) \\
 R(0, a, b) \wedge R(b, 0, c) \wedge R(c, d, 0) \\
 R(0, a, b) \wedge R(b, 1, c) \wedge R(c, d, 1) \\
 R(0, a, b) \wedge R(b, 1, c) \wedge R(c, d, 0)
 \end{array}$$

# Hill Climbing

One approach to such optimization problems.

- Basic idea: move to a neighbor with a better  $f(s)$
- **Q:** how do we define **neighbor**?
  - Not as obvious as our successors in search
  - Problem-specific
  - As we'll see, needs a careful choice

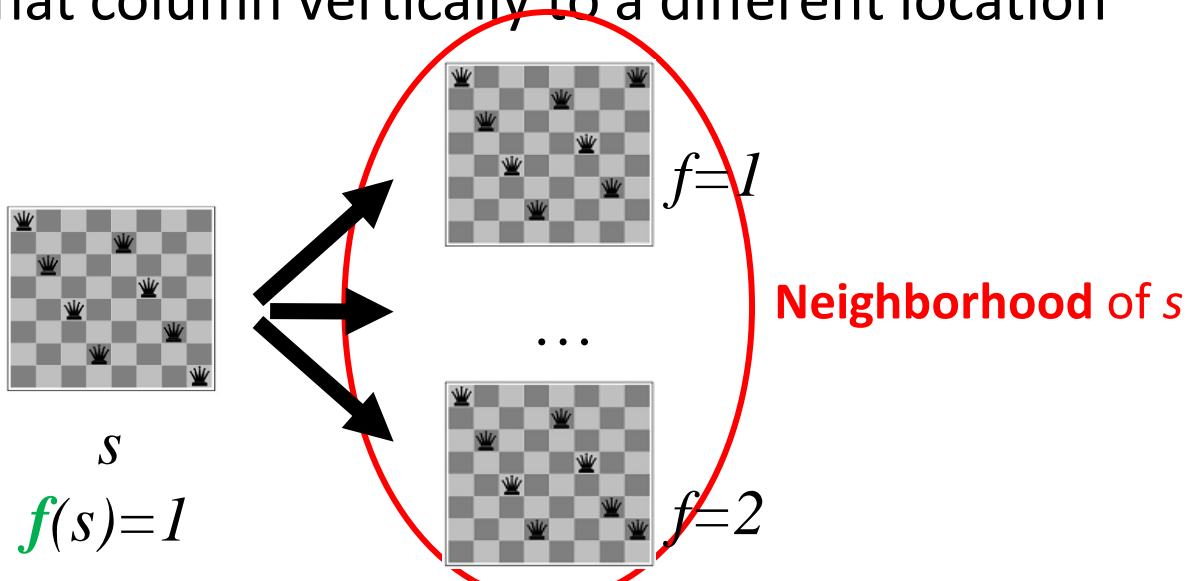




# Defining Neighbors: n Queens

In n Queens, a simple possibility:

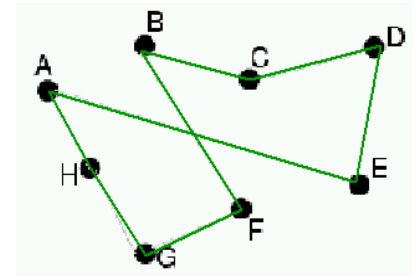
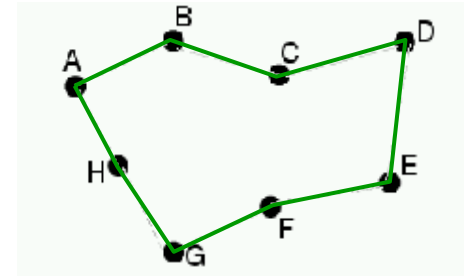
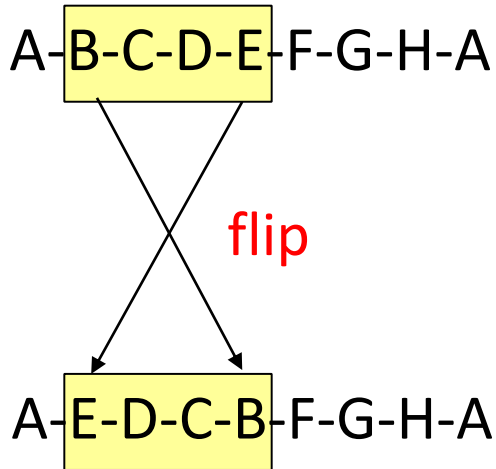
- Look at the **most-conflicting column** (ties? right-most one)
- Move queen in that column vertically to a different location



# Defining Neighbors: TSP

For TSP, can do something similar:

- Define neighbors by small changes
- Example: 2-change: A-E and B-F



# Defining Neighbors: SAT

For Boolean satisfiability,

- Define neighbors by flipping one assignment of one variable

Starting state: TFTTT

(A=**F**, B=F, C=T, D=T, E=T)

(A=T, B=**T**, C=T, D=T, E=T)

(A=T, B=F, C=**F**, D=T, E=T)

(A=T, B=F, C=T, D=**F**, E=T)

(A=T, B=F, C=T, D=T, E=**F**)

$A \vee \neg B \vee C$

$\neg A \vee C \vee D$

$B \vee D \vee \neg E$

$\neg C \vee \neg D \vee \neg E$

$\neg A \vee \neg C \vee E$

# Hill Climbing Neighbors

## Q: What's a **neighbor**?

- **Vague definition.** For a given problem structure, neighbors are states that can be produced by a small change
- **Tradeoff!**
  - Too small? Will get stuck.
  - Too big? Not very efficient
- Q: how to pick a neighbor? Greedy
- Q: terminate? When no neighbor has bigger value



# Hill Climbing Algorithm

## Pseudocode:

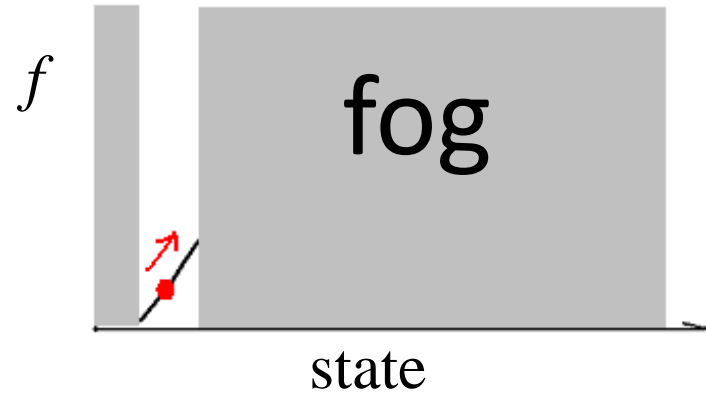
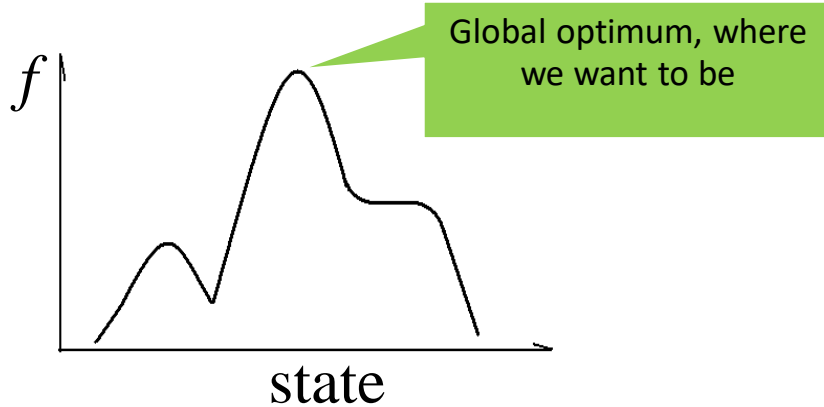
1. Pick initial state  $s$
2. Pick  $t$  in **neighbors**( $s$ ) with the largest  $f(t)$
3. if  $f(t) \leq f(s)$  THEN stop, return  $s$
4.  $s \leftarrow t$ . goto 2.

What could happen? **Local optima!**



# Hill Climbing: Local Optima

Q: Why is it called hill climbing?

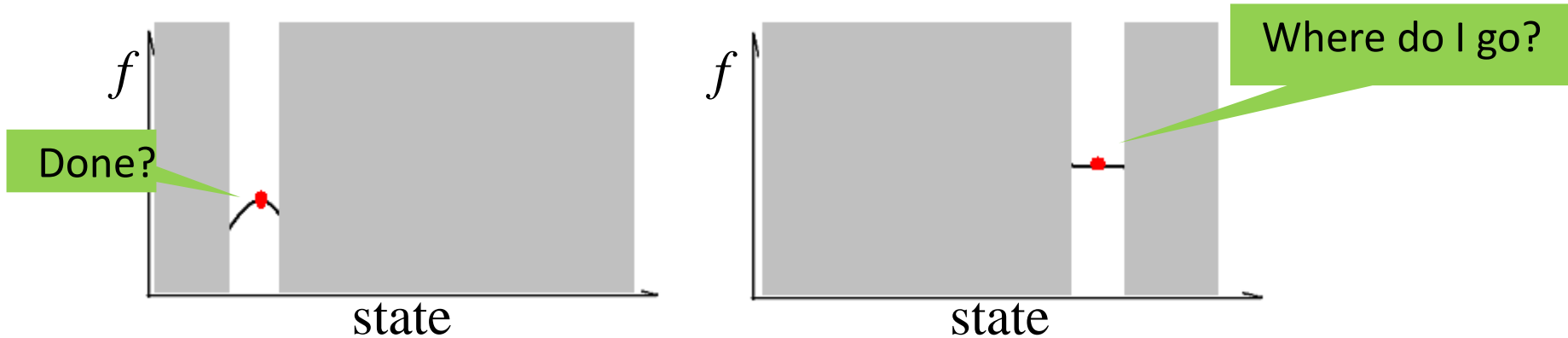


L: What's actually going on.

R: What we get to see.

# Hill Climbing: Local Optima

Note the **local optima**. How do we handle them?



# Escaping Local Optima

## Simple idea 1: random restarts

- Stuck: pick a random new starting point, re-run.
- Do  $k$  times, return best of the  $k$ .



## Simple idea 2: reduce greed

- “Stochastic” hill climbing: randomly select between neighbors
- Probability proportional to the value of neighbors



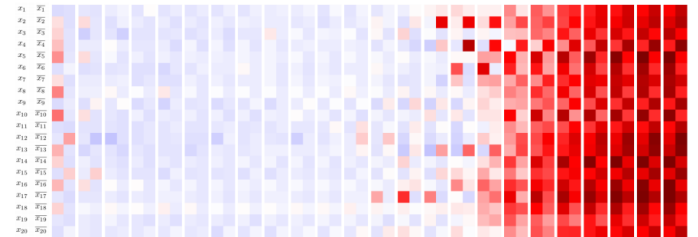
# Hill Climbing: Variations

**Q:** neighborhood too large?

- Generate a few random neighbors, **one at a time**. Take the better one.

**Q:** relax requirement to always go up?

- Often useful for harder problems
- 3SAT algorithm: Walk-SAT



# Break & Quiz

**Q 1.1:** Hill climbing and SGD are related by

- (i) Both head towards local optima
  - (ii) Both require computing a gradient
  - (iii) Both will find the global optimum for a convex problem (when minimizing)
- A. (i)
  - B. (i), (ii)
  - C. (i), (iii)
  - D. All of the above

# Simulated Annealing

A more sophisticated optimization approach.

- **Idea:** move quickly at first, then slow down



# Simulated Annealing

A more sophisticated optimization approach.

- **Idea:** move quickly at first, then slow down
- Pseudocode:

Pick initial state  $x$

For  $k = 0$  through  $k_{\max}$ :

    Reduce temperature  $T$

    Pick a random neighbour,  $y \leftarrow \text{neighbor}(x)$

    If  $f(y) \geq f(x)$ , then  $x \leftarrow y$

    Else, with prob.  $P(f(x), f(y), T)$  then  $x \leftarrow y$

**Output:** the final state  $x$

The interesting bit



# Simulated Annealing: Picking Probability

How do we pick probability  $P$ ?

- Decrease with gap  $|f(x) - f(y)|$
- Decrease with time  $k$

Pick initial state  $x$

For  $k = 0$  through  $k_{\max}$ :

    Reduce temperature  $T$

    Pick a random neighbour,  $y \leftarrow \text{neighbor}(x)$

    If  $f(y) \geq f(x)$ , then  $x \leftarrow y$

    Else, with prob.  $P(f(x), f(y), T)$  then  $x \leftarrow y$

**Output:** the final state  $x$

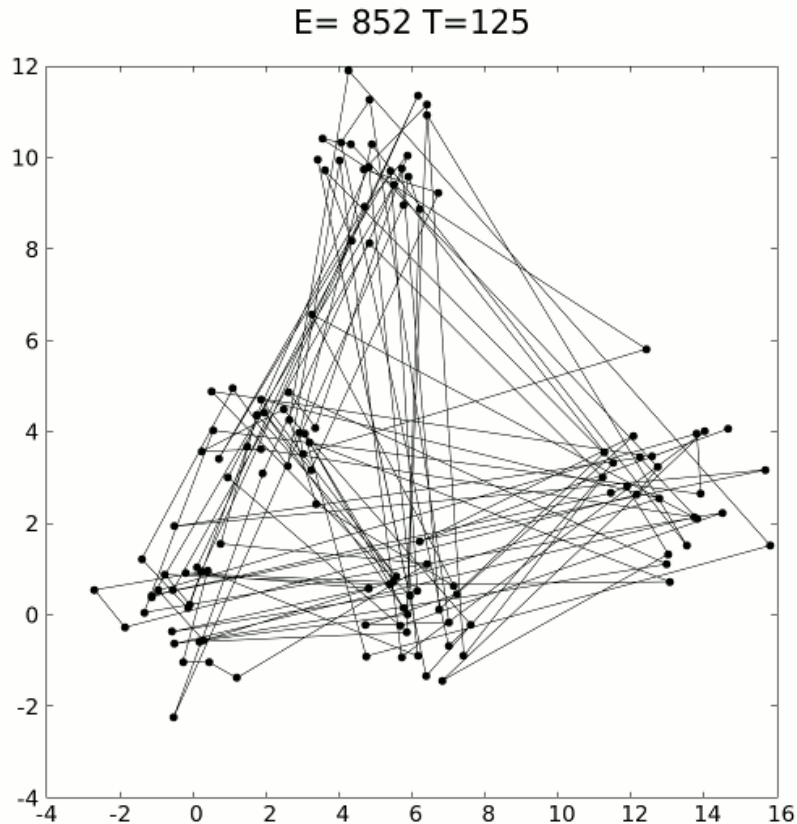
# Simulated Annealing: Picking Probability

How do we pick probability  $P$ ?

$$P(x, y, T) = \exp\left(-\frac{|f(x) - f(y)|}{T}\right)$$

- Decrease with gap  $|f(x) - f(y)|$
- Decrease with time  $k$
- Temperature  $T$  cools over time
  - High temperature, accept any  $y$
  - Low temperature, behaves like hill-climbing
  - Still,  $|f(x) - f(y)|$  plays a role: if big, replacement probability low.

# Simulated Annealing: Visualization



# Simulated Annealing: Picking Parameters

- Have to balance the various parts., e.g., cooling schedule.
  - Too fast: becomes hill climbing, stuck in local optima
  - Too slow: takes too long.
- Combines with variations (e.g., with random restarts)
  - Probably should try hill-climbing first though.
- Inspired by cooling of metals
  - We'll see one more alg. inspired by nature





# Break & Quiz

**Q 2.1:** Which of the following is likely to give the best cooling schedule for simulated annealing?

- A.  $\text{Temp}_{k+1} = \text{Temp}_k * 1.25$
- B.  $\text{Temp}_{k+1} = \text{Temp}_k$
- C.  $\text{Temp}_{k+1} = \text{Temp}_k * 0.8$
- D.  $\text{Temp}_{k+1} = \text{Temp}_k * 0.0001$

# Break & Quiz

**Q 2.2:** Which of the following would be better to solve with simulated annealing than A\* search?

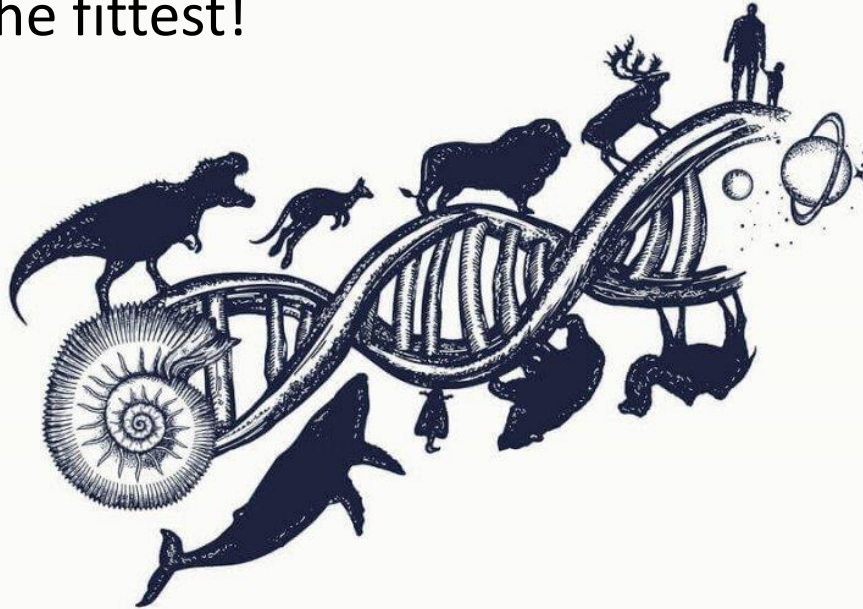
- i. Finding the smallest set of vertices in a graph that involve all edges
- ii. Finding the fastest way to schedule jobs with varying runtimes on machines with varying processing power
- iii. Finding the fastest way through a maze

- A. (i)
- B. (ii)
- C. (i) and (ii)
- D. (ii) and (iii)

# Genetic Algorithms

Another optimization approach based on nature

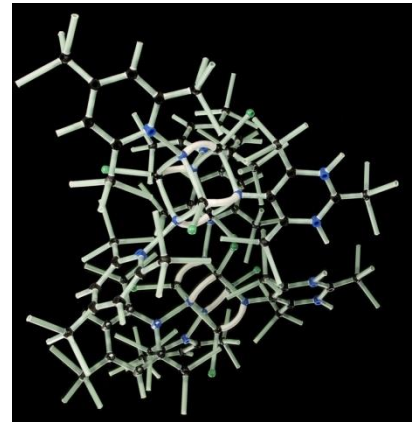
- Survival of the fittest!



# Evolution Review

Encode genetic information in DNA (four bases)

- A/C/T/G: nucleobases acting as symbols
- Two types of changes
  - Crossover: exchange between parents' codes
  - Mutation: rarer random process
    - Happens at individual level



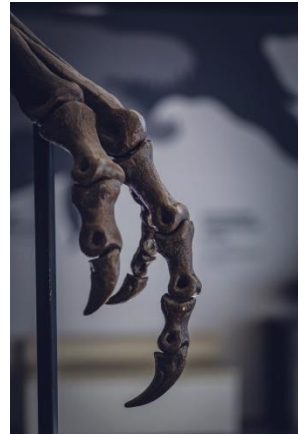
# Natural Selection

## Competition for resources

- Organisms better fit → better probability of reproducing
- Repeated process: fit become larger proportion of population

## Goal: use these principles for optimization

- New terminology: state  $s$  ‘**individual**’
- Value  $f(s)$  is now the ‘**fitness**’

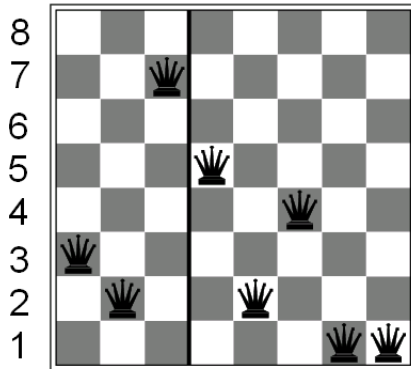


# Genetic Algorithms Setup I

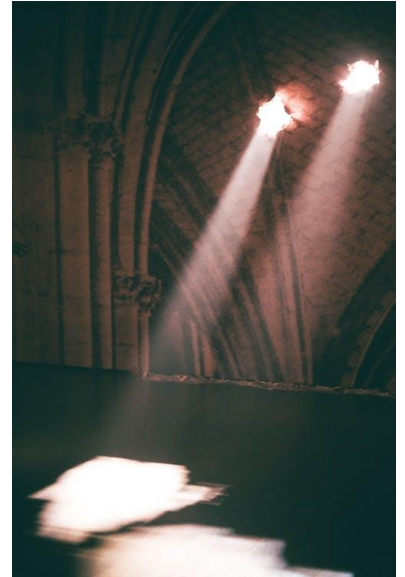
Keep around a fixed number of states/individuals

- A bit like beam search
- Call this the **population**

For our n Queens game example, an individual:



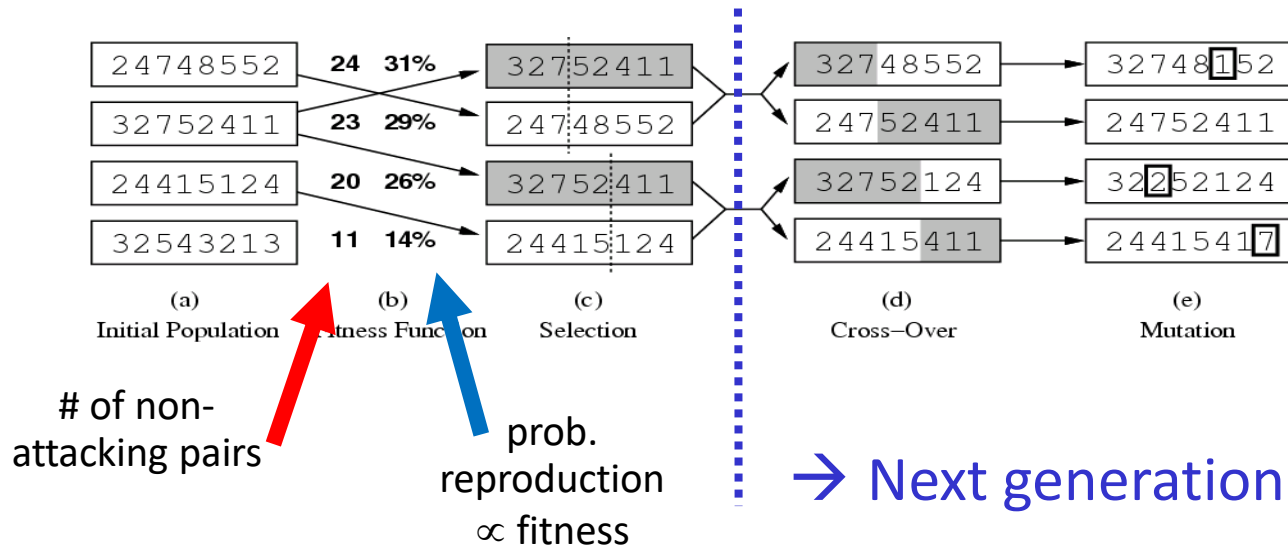
(3 2 7 5 2 4 1 1)



# Genetic Algorithms Setup II

Goal of genetic algorithms: optimize using principles inspired by mechanism for evolution

- E.g., analogous to **natural selection**, **cross-over**, and **mutation**



# Genetic Algorithms Pseudocode

Just one variant:

1. Let  $s_1, \dots, s_N$  be the current population
2. Let  $p_i = f(s_i) / \sum_j f(s_j)$  be the reproduction probability
3. for  $k = 1; k < N; k += 2$ 
  - parent1 = randomly pick according to  $p$
  - parent2 = randomly pick another
  - randomly select a crossover point, swap strings of parents 1, 2 to generate children  $t[k], t[k+1]$
4. for  $k = 1; k \leq N; k++$ 
  - Randomly mutate each position in  $t[k]$  with a small probability (mutation rate)
5. The new generation replaces the old:  $\{s\} \leftarrow \{t\}$ . Repeat



# Reproduction: Proportional Selection

Reproduction probability:  $p_i = f(s_i) / \sum_j f(s_j)$

- **Example:**  $\sum_j f(s_j) = 5+20+11+8+6=50$
- $p_1=5/50=10\%$

Individual	Fitness	Prob.
A	5	10%
B	20	40%
C	11	22%
D	8	16%
E	6	12%



# Example: Scheduling Courses

Let's run through an example:

- **5 courses: A,B,C,D,E**
- *3 time slots: Mon/Wed, Tue/Thu, Fri/Sat*
- Students wish to enroll in three courses
- Goal: maximize student enrollment

Courses	Students
A B C	2
A B D	7
A D E	3
B C D	4
B D E	10
C D E	5

# Example: Scheduling Courses

Let's run through an example:

- State: course assignment to time slot

M	M	F	T	M
A	B	C	D	E

= MMFTM

- Here:
  - Courses A, B, E scheduled Mon/Wed
  - Course D scheduled Tue/Thu
  - Course C scheduled Fri/Sat

Courses	Students
A B C	2
A B D	7
A D E	3
B C D	4
B D E	10
C D E	5

# Example: Scheduling Courses

Value of a state? Say MMFTM

Courses	Students	Can enroll?
A B C	2	No
A B D	7	No
A D E	3	No
B C D	4	Yes
B D E	10	No
C D E	5	Yes

- Here  $4+5=9$  students can enroll in desired courses

# Example: Scheduling Courses

## First step:

- Randomly initialize and evaluate states

MMFTM = 9

MMFTM = 26%

TTFMM = 4

TTFMM = 11%

FMTTF = 19

FMTTF = 54%

MTTTF = 3

MTTTF = 9%

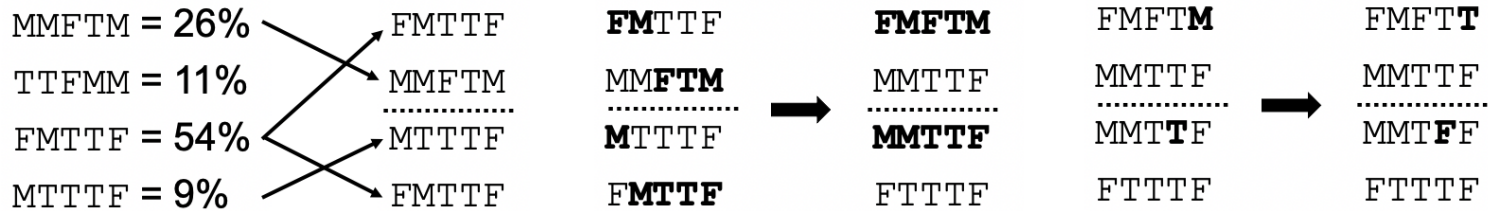
- Calculate reproduction probabilities

Courses	Students
A B C	2
A B D	7
A D E	3
B C D	4
B D E	10
C D E	5

# Example: Scheduling Courses

Next steps:

- Select parents using reproduction probabilities
- Perform crossover
- Randomly mutate new children



# Example: Scheduling Courses

Continue:

- Now, get our function values for updated population
- Calculate reproduction probabilities

$$FMFTT = 11$$

$$FMFTT = 39\%$$

$$MMTTF = 13$$

$$MMTTF = 46\%$$

$$MMTFF = 4$$

$$MMTFF = 14\%$$

$$FTTTF = 0$$

$$FTTTF = 0\%$$

Courses	Students
A B C	2
A B D	7
A D E	3
B C D	4
B D E	10
C D E	5

# Variations & Concerns

Many **possibilities**:

- Parents survive to next generation
- Ranking instead of exact value of  $f(s)$  for reproduction probabilities

Some **challenges**

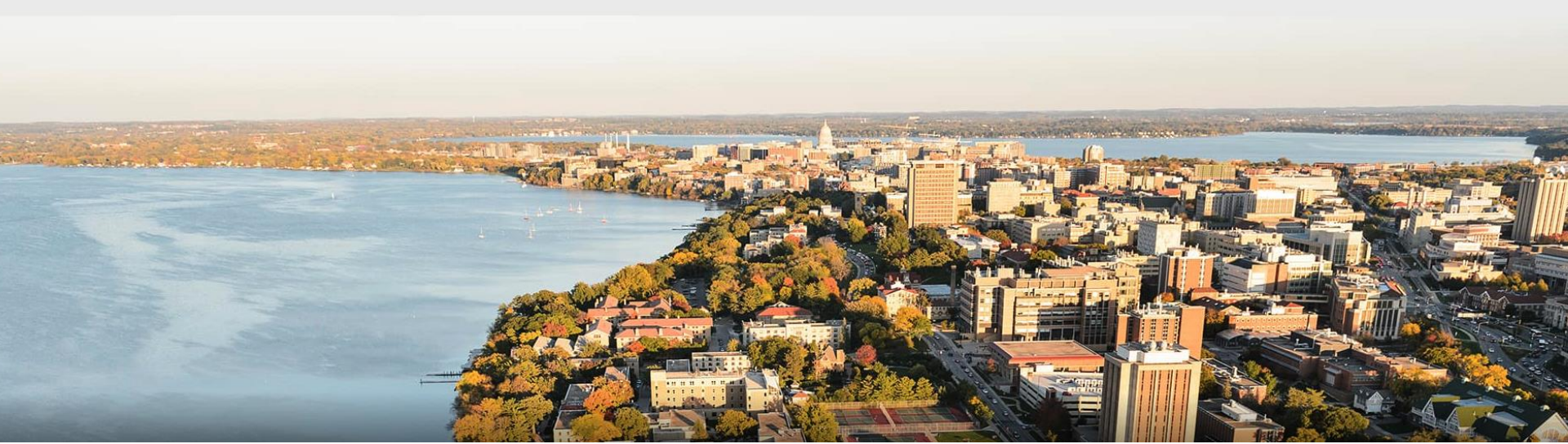
- State encoding
- Lack of diversity: converge too soon
- Must pick a lot of parameters





# Summary

- Challenging optimization problems
  - First, try hill climbing. Simplest solution
- Simulated annealing
  - More sophisticated approach; helps with local optima
- Genetic algorithms
  - Biology-inspired optimization routine



**Acknowledgements:** Adapted from materials by Fred Sala, Jerry Zhu + Tony Gitter (University of Wisconsin), Andrew Moore