

Lecture 5: Minima of Convex Functions; Algorithmic Setup

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1 Minima of convex functions

Consider the constrained problem

$$\min_{x \in \mathcal{X}} f(x). \quad (\text{P})$$

Recall definition of convex functions.

Theorem 1. Consider the problem (P). Suppose f is convex, and \mathcal{X} is convex, closed and non-empty. Then:

1. Any local solution to (P) is also a global solution.
2. The set of global solutions to (P) is convex.

Proof. Part 1: Suppose f.p.o.c. that x^* is a local but not a global solution. Then there exists $\bar{x} \in \mathcal{X}$ such that $f(\bar{x}) < f(x^*)$. As \mathcal{X} is convex, for all $\alpha \in (0, 1)$,

$$(1 - \alpha)x^* + \alpha\bar{x} \in \mathcal{X}.$$

As f is convex, for all $\alpha \in (0, 1)$:

$$f((1 - \alpha)x^* + \alpha\bar{x}) \leq (1 - \alpha)f(x^*) + \alpha f(\bar{x}) < f(x^*).$$

Hence every neighborhood of x^* must include a point $(1 - \alpha)x^* + \alpha\bar{x}$ for some $\alpha > 0$ that will have a strictly lower function value. So x^* cannot be a local solution, a contradiction.

Part 2: Let $x^*, \bar{x} \in \mathcal{X}$ be any two global solutions.

\mathcal{X} is convex $\implies (1 - \alpha)x^* + \alpha\bar{x} \in \mathcal{X}$.

f is convex \implies

$$f((1 - \alpha)x^* + \alpha\bar{x}) \leq (1 - \alpha)f(x^*) + \alpha f(\bar{x}) = f(x^*) = f(\bar{x})$$

$\implies f((1 - \alpha)x^* + \alpha\bar{x}) = f(x^*)$, so $(1 - \alpha)x^* + \alpha\bar{x}$ is also a global solution \implies the set of global solution is convex. \square

1.1 Differentiable convex functions

Theorem 2 (Equivalent characterization of convexity). *The following are true.*

1. Let $f : \mathbb{R}^d \rightarrow \bar{\mathbb{R}}$ be continuously differentiable. The function f is convex if and only if

$$\forall x, y : f(y) \geq f(x) + \langle \nabla f(x), y - x \rangle. \quad (1)$$

(A picture. From local to global.)

2. Let $f : \mathbb{R}^d \rightarrow \bar{\mathbb{R}}$ be twice continuously differentiable. The function f is convex if and only if

$$\forall x : \nabla^2 f(x) \succcurlyeq 0.$$

Proof. Part 1, “only if”: By convexity of f , for any $\alpha \in (0, 1)$:

$$\begin{aligned} f((1-\alpha)x + \alpha y) &\leq (1-\alpha)f(x) + \alpha f(y) \\ \xrightarrow{\text{rearranging}} f(y) - f(x) &\geq \frac{f(x + \alpha(y-x)) - f(x)}{\alpha} \stackrel{\text{Taylor's}}{\geq} \frac{\langle \nabla f(x), \alpha(y-x) \rangle + o(\alpha)}{\alpha}. \end{aligned}$$

Taking $\alpha \rightarrow 0$ gives (1)

Part 1, “if”: Take any x, y and $\alpha \in (0, 1)$. Set $z = (1-\alpha)x + \alpha y$. Apply (1) to x, z and y, z :

$$f(x) \geq f(z) + \alpha \langle \nabla f(z), x - y \rangle, \quad (2)$$

$$f(y) \geq f(z) + (1-\alpha) \langle \nabla f(z), y - x \rangle. \quad (3)$$

(2) $\times (1-\alpha)$ + (3) $\times \alpha$ gives

$$(1-\alpha)f(x) + \alpha f(y) \geq f(z),$$

which implies convexity of f .

Part 2: See Wright-Recht, Lemma 2.9. □

Theorem 3 (Sufficient condition for global optimality). Consider the problem (P), where f is continuously differentiable and convex. If $x^* \in \mathcal{X}$ and $\nabla f(x^*) = 0$, then x^* is a global minimizer of f .

Proof. Use Part 1 of Theorem 2:

$$\forall x : f(x) \geq f(x^*) + \langle \nabla f(x^*), x - x^* \rangle = f(x^*).$$

□

Remark 1. Theorem 3 holds for both unconstrained (i.e., $\mathcal{X} = \mathbb{R}^d$) and constrained problems. Using terminology from last time, x^* being a stationary point is sufficient for global optimality. For unconstrained problem, this is also necessary (Lecture 4, Theorem 1). For constrained problem, this may not be necessary (example).

2 Strongly convex functions

We use Euclidean norm $\|\cdot\|_2$ in this section.

Definition 1 (Strong convexity). Given $m > 0$, we say that $f : \mathbb{R}^d \rightarrow \bar{\mathbb{R}}$ is *strongly convex* with modulus/parameter m (or *m-strongly convex for short*), if

$$\forall x, y \in \mathbb{R}^d : f((1-\alpha)x + \alpha y) \leq (1-\alpha)f(x) + \alpha f(y) - \frac{m}{2}(1-\alpha)\alpha \|y - x\|_2^2.$$

Remark 2. Verify yourself that the above is equivalent to convexity of the function $f(x) - \frac{m}{2}\|x\|_2^2$.

Theorem 4 (Equivalent characterization of strong convexity). The following hold.

1. Suppose f is continuously differentiable. Then f is m -strong convexity if and only if

$$f(y) \geq f(x) + \langle \nabla f(x), y - x \rangle + \frac{m}{2} \|y - x\|_2^2.$$

(A picture. Compare with convexity only. Complements L -smoothness.)

2. Suppose f is twice continuously differentiable. Then f is m -strong convexity if and only if

$$\forall x : \nabla^2 f(x) \succeq mI.$$

(Compare with L -smoothness)

Proof. Apply Theorem 2 to the function $f(x) - \frac{m}{2} \|x\|_2^2$. □

Theorem 5. Suppose that $f : \mathbb{R}^d \rightarrow \bar{\mathbb{R}}$ is continuously differentiable and m -strongly convex for some $m > 0$. If $x^* \in \mathcal{X}$ satisfies $\nabla f(x^*) = 0$, then x^* is the unique global minimizer of f .

Proof. By Part 1 of Theorem 4:

$$f(x) \geq f(x^*) + \langle \nabla f(x^*), x - x^* \rangle + \underbrace{\frac{m}{2} \|x - x^*\|_2^2}_{>0 \text{ unless } x=x^*}.$$

□

3 Algorithmic setup

1. First-order oracle:

$$x \longrightarrow \text{oracle} \longrightarrow f(x), \nabla f(x)$$

2. Second-order oracle:

$$x \longrightarrow \text{oracle} \longrightarrow f(x), \nabla f(x), \nabla^2 f(x)$$

All algorithms we consider in this course are iterative:

- start with some x_0
- at iteration $k = 0, 1, 2, \dots$
 - get oracle answers for x_k , choose x_{k+1}

4 Basic descent methods

Take the form

$$x_{k+1} = x_k + \alpha_k p_k, \quad k = 0, 1, \dots$$

Definition 2. $p \in \mathbb{R}^d$ is a descent direction for f at x if

$$f(x + tp) < f(x)$$

for all sufficiently small $t > 0$.

Proposition 1. *If f is continuously differentiable (in a neighborhood of x), then any p such that $\langle -\nabla f(x), p \rangle > 0$ is a descent direction.*

Proof. By Taylor's theorem:

$$f(x + tp) = f(x) + t \langle \nabla f(x + \gamma tp), p \rangle$$

for some $\gamma \in (0, 1)$. We know that $\langle \nabla f(x), p \rangle < 0$. As ∇f is continuous, for all sufficiently small $t > 0$,

$$\langle \nabla f(x + \gamma tp), p \rangle < 0,$$

hence $f(x + tp) < f(x)$. □

5 Gradient descent

What would be a good descent direction?

Could try to move in the direction of $-\nabla f(x)$. Justification: Look at all p with $\|p\|_2 = 1$. Then

$$\inf_{\|p\|_2=1} \langle \nabla f(x), p \rangle = -\|\nabla f(x)\|_2$$

attained for $p = -\frac{\nabla f(x)}{\|\nabla f(x)\|_2}$.

“Simplest” descent algorithm:

$$x_{k+1} = x_k - \alpha_k \nabla f(x_k),$$

where α_k is the step size. Ideally, choose α_k small enough so that

$$f(x_{k+1}) < f(x_k)$$

when $\nabla f(x_k) \neq 0$.

Known as “gradient method”, “gradient descent”, “steepest descent” (w.r.t. the ℓ_2 norm).