Lecture 11: Acceleration via Regularization and Restarting; Lower Bounds

Yudong Chen

Last week we discussed two variants of Nesterov's accelerated gradient descent (AGD).

Algorithm 1 Nesterov's AGD, smooth and strongly convex input: initial x_0 , strong convexity and smoothness parameters m, L, number of iterations Kinitialize: $x_{-1} = x_0$, $\beta = \frac{\sqrt{L/m}-1}{\sqrt{L/m}+1}$. for k = 0, 1, ... K $y_k = x_k + \beta (x_k - x_{k-1})$ $x_{k+1} = y_k - \frac{1}{L} \nabla f(y_k)$ return x_K

Theorem 1. For Nesterov's AGD Algorithm 1 applied to m-strongly convex L-smooth f, we have

$$f(x_k) - f^* \le \left(1 - \sqrt{\frac{m}{L}}\right)^k \cdot \frac{(L+m) \|x_0 - x^*\|_2^2}{2}.$$

Equivalently, we have $f(x_k) - f^* \le \epsilon$ after at most $k = O\left(\sqrt{\frac{L}{m}}\log\frac{L\|x_0 - x^*\|_2^2}{\epsilon}\right)$ iterations.

Algorithm 2 Nesterov's AGD, smooth convex input: initial x_0 , smoothness parameter L, number of iterations Kinitialize: $x_{-1} = x_0$, $\lambda_0 = 0$, $\beta_0 = 0$. for k = 0, 1, ..., K $y_k = x_k + \beta_k (x_k - x_{k-1})$ $x_{k+1} = y_k - \frac{1}{L} \nabla f(y_k)$ $\lambda_{k+1} = \frac{1 + \sqrt{1 + 4\lambda_k^2}}{2}$, $\beta_{k+1} = \frac{\lambda_k - 1}{\lambda_{k+1}}$ return x_K

Theorem 2. For Nesterov's AGD Algorithm 2 applied to L-smooth convex f, we have

$$f(x_k) - f(x^*) \le \frac{2L \|x_0 - x^*\|_2^2}{k^2}$$

In this lecture, we will show that the two types of acceleration above are closely related: we can use one to derive the other. We then show that in a certain precise (but narrow) sense, the convergence rates of AGD are optimal among first-order methods. For this reason, AGD is also known as Nesterov's *optimal* method.

1 Acceleration via regularization

Suppose we only know the AGD method for *strongly* convex functions (Algorithm 1) and its $(1 - \sqrt{\frac{m}{L}})^k$ guarantee (Theorem 1). Can we use it as a subroutine to develop an accelerated algorithm for (non-strongly) convex functions with a $\frac{1}{k^2}$ convergence rate?

The answer is yes (up to logarithmic factors). One approach is to add a *regularizer* $\epsilon ||x||_2^2$ to f(x) and apply Algorithm 1 to the function $f(x) + \epsilon ||x||_2^2$, which is strongly convex. See HW 3.

2 Acceleration via restarting

In the opposite direction, suppose we only know the AGD method for (non-strongly) convex functions (Algorithm 2) and its $\frac{1}{k^2}$ guarantee (Theorem 2). Can we use it as a subroutine to develop an accelerated algorithm for *strongly* convex functions with a $(1 - \sqrt{\frac{m}{L}})^k$ convergence rate (equivalently, a $\sqrt{\frac{L}{m}} \log \frac{1}{\epsilon}$ iteration complexity)?

This is possible using a classical and powerful idea in optimization: *restarting*. See Algorithm 3. In each round, we run Algorithm 2 for $\sqrt{\frac{8L}{m}}$ iterations to obtain \overline{x}_{t+1} . In the next round, we restart Algorithm 2 using \overline{x}_{t+1} as the initial solution and run for another $\sqrt{\frac{8L}{m}}$ iterations. This is repeated for *T* rounds.

Algorithm 3 Restarting AGD

input: initial \overline{x}_0 , strong convexity and smoothness parameters *m*, *L*, number of rounds *T* **for** t = 0, 1, ..., T

Run Algorithm 2 with \overline{x}_t (initial solution), *L* (smoothness parameter), $\sqrt{\frac{8L}{m}}$ (number of iterations) as the input. Let \overline{x}_{t+1} be the output.

return \overline{x}_T

Exercise 1. How is Algorithm 3 different from running Algorithm 2 without restarting for $T \times \sqrt{\frac{8L}{m}}$ iterations?

2.1 Analysis

Suppose *f* is *m*-strongly convex and *L*-smooth. By Theorem 2, we know that

$$f(\overline{x}_{t+1}) - f(x^*) \le \frac{2L \|\overline{x}_t - x^*\|_2^2}{8L/m} = \frac{m \|\overline{x}_t - x^*\|_2^2}{4}.$$

By strong convexity, we have

$$f(\overline{x}_t) \geq f(x^*) + \underbrace{\langle \nabla f(x^*), \overline{x}_t - x^* \rangle}_{=0} + \frac{m}{2} \|\overline{x}_t - x^*\|_2^2,$$

hence $\|\overline{x}_t - x^*\|_2^2 \leq \frac{2}{m} (f(\overline{x}_t) - f(x^*))$. Combining, we get

$$f(\overline{x}_{t+1}) - f(x^*) \leq \frac{f(\overline{x}_t) - f(x^*)}{2}.$$

That is, each round of Algorithm 3 halves the optimality gap. It follows that

$$f(\overline{x}_T) - f(x^*) \le \left(\frac{1}{2}\right)^T \left(f(\overline{x}_0) - f(x^*)\right).$$

Therefore, $f(\overline{x}_T) - f(x^*) \le \epsilon$ can be achieved after at most

$$T = O\left(\log \frac{f(\overline{x}_0) - f(x^*)}{\epsilon}\right)$$
 rounds,

which corresponds to a total of

$$T \times \sqrt{\frac{8L}{m}} = O\left(\sqrt{\frac{L}{m}}\log\frac{f(\overline{x}_0) - f(x^*)}{\epsilon}\right)$$
 AGD iterations.

This iteration complexity is the same as Theorem 1 up to a logarithmic factor.

Remark 1. Note how strong convexity is needed in the above argument.

Remark 2. Optional reading: This overview article discusses restarting as a general/meta algorithmic technique.

3 Lower bounds

In this section, we consider a class of first-order iterative algorithms that satisfy $x_0 = 0$, and

$$x_{k+1} \in \operatorname{Lin}\left\{\nabla f(x_0), \nabla f(x_1), \dots, \nabla f(x_k)\right\}, \qquad \forall k \ge 0,$$
(1)

where the RHS denotes the linear subspace spanned by $\nabla f(x_0)$, $\nabla f(x_1)$,..., $\nabla f(x_k)$; in other words, x_{k+1} is an (arbitrary) linear combination of the gradients at the previous (k + 1) iterates.

3.1 Smooth and convex *f*

Theorem 3. There exists an L-smooth convex function f such that any first-order method in the sense of (1) must satisfy

$$f(x_k) - f(x^*) \ge \frac{3L \|x_0 - x^*\|_2^2}{32(k+1)^2}$$

Comparing with this lower bound, we see that the $\frac{L}{k^2}$ rate for AGD in Theorem 2 is optimal/unimprovable (up to constants).

Proof of Theorem **3***.* Let $A \in \mathbb{R}^{d \times d}$ be the matrix given by

$$A_{ij} = \begin{cases} 2, & i = j \\ -1, & j \in \{i - 1, i + 1\} \\ 0, & \text{otherwise.} \end{cases}$$
(2)

Explicitly,

$$A = \begin{bmatrix} 2 & -1 & 0 & 0 & \cdots & \cdots & 0 \\ -1 & 2 & -1 & 0 & \cdots & \cdots & 0 \\ 0 & -1 & 2 & -1 & 0 & \cdots & 0 \\ & & \ddots & \ddots & \ddots & & \\ 0 & \cdots & & & -1 & 2 & -1 \\ 0 & \cdots & & & & -1 & 2 \end{bmatrix}$$

Let $e_i \in \mathbb{R}^d$ denote the *i*-th standard basis vector. Consider the quadratic function

$$f(x) = \frac{L}{8}x^{\top}Ax - \frac{L}{4}x^{\top}e_1,$$

which is convex and *L*-smooth since $0 \leq A \leq 4I$. Note that $\nabla f(x) = \frac{L}{4}(Ax - e_1)$. By induction, we can show that for $k \ge 1$,

$$x_k \in \operatorname{Lin} \{e_1, Ax_1, \ldots, Ax_{k-1}\} \subseteq \operatorname{Lin} \{e_1, \ldots, e_k\}$$

Therefore, if we let $A_k \in \mathbb{R}^{d \times d}$ denote the matrix obtained by zeroing out the entries of A outside the top-left $k \times k$ block, then

$$f(x_k) = \frac{L}{8} x_k^{\top} A_k x_k - \frac{L}{4} x_k^{\top} e_1 \ge f_k^* := \min_x \left\{ \frac{L}{8} x^{\top} A_k x - \frac{L}{4} x^{\top} e_1 \right\}.$$

By setting gradient to zero, we find that the minimum above is attained by

$$x_k^* := \left(1 - \frac{1}{k+1}, 1 - \frac{2}{k+1}, \dots, 1 - \frac{k}{k+1}, 0, \dots, 0\right)^\top \in \mathbb{R}^d,$$

with $f_k^* = -\frac{L}{8} \left(1 - \frac{1}{k+1}\right)$. It follows that the global minimizer $x^* = x_d^*$ of f satisfies $f(x^*) = f_d^* =$ $-\frac{L}{8}(1-\frac{1}{d+1})$ and (since $x_0 = 0$)

$$\|x_d^* - x_0\|_2^2 = \|x_d^*\|_2^2 = \sum_{i=1}^d \left(1 - \frac{i}{d+1}\right)^2 \le \frac{d+1}{3}.$$

Combining pieces and taking d = 2k + 1, we have

$$f(x_k) - f(x^*) \ge f_k^* - f_d^* = \frac{L}{8} \left(\frac{1}{k+1} - \frac{1}{2k+2} \right)$$
$$= \frac{L}{16} \frac{k+1}{(k+1)^2}$$
$$= \frac{L}{32} \frac{d+1}{(k+1)^2}$$
$$\ge \frac{3L}{32} \frac{\|x^* - x_0\|_2^2}{(k+1)^2}.$$

3.2 Smooth and strongly convex *f*

For strongly convex functions, we have the following lower bound, which shows that the $\left(1 - \frac{1}{\sqrt{L/m}}\right)^k$ rate of AGD in Theorem 1 cannot be significantly improved.

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Theorem 4. There exists an m-strongly convex and L-smooth function such that any first-order method in the sense of (1) must satisfy

$$f(x_k) - f(x^*) \ge \frac{m}{2} \left(1 - \frac{4}{\sqrt{L/m}}\right)^{k+1} \|x_0 - x^*\|_2^2.$$

Proof. Let $A \in \mathbb{R}^{d \times d}$ be defined in (2) above and consider the function

$$f(x) = \frac{L-m}{8} \left(x^{\top} A x - 2x^{\top} e_1 \right) + \frac{m}{2} \|x\|_2^2,$$

which is *L*-smooth and *m*-strongly convex. Strong convexity implies that

$$f(x_k) - f(x^*) \ge \frac{m}{2} \|x_k - x^*\|_2^2.$$
(3)

A similar argument as above shows that $x_k \in \text{Lin} \{e_1, \ldots, e_k\}$, hence

$$\|x_k - x^*\|_2^2 \ge \sum_{i=k+1}^d x^*(i)^2,$$
(4)

where $x^*(i)$ denotes the *i*th entry of x^* . For simplicity we take $d \to \infty$ (we omit the formal limiting argument).¹ The minimizer x^* can be computed by setting the gradient of *f* to zero, which gives an infinite set of equations

$$1 - 2\frac{L/m + 1}{L/m - 1}x^*(1) + x^*(2) = 0,$$

$$x^*(k - 1) - 2\frac{L/m + 1}{L/m - 1}x^*(k) + x^*(k + 1) = 0, \qquad k = 2, 3, \dots$$

Solving these equations gives

$$x^*(i) = \left(\frac{\sqrt{L/m} - 1}{\sqrt{L/m} + 1}\right)^i, \quad i = 1, 2, \dots$$
 (5)

Combining pieces, we obtain

$$f(x_{k}) - f(x^{*}) \geq \frac{m}{2} \sum_{i=k+1}^{\infty} x^{*}(i)^{2} \qquad \text{by (3) and (4)}$$

$$\geq \frac{m}{2} \left(\frac{\sqrt{L/m} - 1}{\sqrt{L/m} + 1}\right)^{2(k+1)} \|x_{0} - x^{*}\|_{2}^{2} \qquad \text{by (5) and } x_{0} = 0$$

$$= \frac{m}{2} \left(1 - \frac{4}{\sqrt{L/m} + 1} + \frac{4}{(\sqrt{L/m} + 1)^{2}}\right)^{k+1} \|x_{0} - x^{*}\|_{2}^{2}$$

$$\geq \frac{m}{2} \left(1 - \frac{4}{\sqrt{L/m}}\right)^{k+1} \|x_{0} - x^{*}\|_{2}^{2}.$$

Remark 3. The lower bounds in Theorems 3 and 4 are in the worst-case/minimax sense: one cannot find a first-order method that achieves a better convergence rate on *all* smooth convex functions than AGD. This, however, does not prevent better rates to be achieved for a sub class of such functions. It is also possible to achieve better rates by using higher-order information (e.g., the Hessian).

¹The convergence rates for AGD in Theorems 1 and 2 do not explicitly depend on the dimension d, hence these results can be generalized to infinite dimensions.