# ECO2020 Tutorial 3

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# **1** Question 1

(Comprehensive Exam June 2005 Q3)

## 1.1 Part a

Competitive Equilibrium requires:

1. Consumers maximize utility:

$$q_{1}(\alpha) \geq \theta q_{2}(\alpha) \text{ for } \theta \in \Theta_{1}$$
$$q_{1}(\alpha) \leq \theta q_{2}(\alpha) \text{ for } \theta \in \Theta_{2}$$

2. Markets clear:

$$\int_{\Theta_1} f(\theta) \, d\theta = \alpha$$

 $\begin{array}{l} \text{Consider } \hat{\theta} = \frac{q_1(\alpha)}{q_2} \, (\alpha), \\ \text{If } \theta < \hat{\theta}, \end{array} \end{array}$ 

$$q_{1}(\alpha) = q_{2}(\alpha) \hat{\theta} > q_{2}(\alpha) \theta$$
$$\Rightarrow \theta \in \Theta_{1}$$

If  $\theta > \hat{\theta}$ ,

$$q_{1}(\alpha) = q_{2}(\alpha) \hat{\theta} < q_{2}(\alpha) \theta$$
$$\Rightarrow \theta \in \Theta_{2}$$

Therefore,

$$\Theta_1 = \begin{bmatrix} 0, \hat{\theta} \end{bmatrix}$$
$$\Theta_2 = \begin{bmatrix} \hat{\theta}, 1 \end{bmatrix}$$

#### 1.2 Part b

$$\hat{\theta} = \frac{q_1(\alpha)}{q_2}(\alpha)$$
$$= \frac{2a + \alpha}{2a + \alpha + 1}$$

Then market clearing condition becomes:

$$\begin{split} &\int_{\Theta_1} f\left(\theta\right) d\theta = \alpha \\ &\Rightarrow \int_0^{\hat{\theta}} f\left(\theta\right) d\theta = \alpha \\ &\Rightarrow F\left(\hat{\theta}\right) = \alpha \\ &\Rightarrow F\left(\frac{2a+\alpha}{2a+\alpha+1}\right) = \alpha \end{split}$$

Let 
$$G(\alpha) = F\left(\frac{2a+\alpha}{2a+\alpha+1}\right)$$
:

G is continuous since is composite of continuous functions on  $\alpha \in [0, 1]$ . The support is [0, 1], compact convex,

Apply Brower fixed point theorem:

$$\exists \alpha^* \text{ such that } G(\alpha^*) = \alpha^*$$

# **2** Question 2

(Comprehensive Exam June 2007 Q3)

#### 2.1 Part a

Let allocations be functions  $x:[0,1] \to \mathbb{R}_+$  and  $y:[0,1] \to \mathbb{R}_+$ 

Pareto Efficient allocation requires:

1. Individual feasibility:

$$x\left( heta 
ight) \geq 0$$
 and  $y\left( heta 
ight) \geq 0$  for  $heta \in \left[ 0,1 
ight]$ 

2. Aggregate feasibility:

$$\int_{0}^{1} x(\theta) = 1$$
$$\int_{0}^{1} y(\theta) = z$$

3. Optimality:

There does not exist  $(\tilde{x}, \tilde{y})$  such that:

$$(\tilde{x}(\theta), \tilde{y}(\theta)) \succeq_{\theta} (x(\theta), y(\theta) \text{ for } \theta \in [0, 1]$$

with strict inequality on some non-empty open interval  $(a, b) \in [0, 1]$ .

## 2.2 Part b

Walrasian Equilibrium requires:

1. Consumers maximize utility:

$$\max_{x(\theta),y(\theta)} x(\theta)^{\theta} y(\theta)^{1-\theta} \text{ such that } p_x x(\theta) + p_y y(\theta) = p_x + p_y z \text{ for } \theta \in [0,1]$$

2. Markets clear:

$$\int_{0}^{1} x(\theta) = 1$$
$$\int_{0}^{1} y(\theta) = z$$

### 2.3 Part c

Normalize  $p_y = 1$ , then

$$x(\theta) = \theta \frac{p_x + z}{p_x}$$
$$y(\theta) = (1 - \theta)(p_x + z)$$

Check X market clearing:

$$\int_{0}^{1} \theta \frac{p_{x} + z}{p_{x}} = 1$$
$$\Rightarrow \frac{1}{2} \frac{p_{x} + z}{p_{x}} = 1$$
$$\Rightarrow p_{x} = z$$

Or check Y market clearing:

$$\int_0^1 (1-\theta) (p_x + z) = z$$
$$\Rightarrow \frac{1}{2} (p_x + z) = z$$
$$\Rightarrow p_x = z$$

Therefore, WE has allocations:

$$x(\theta) = \theta \frac{z+z}{z} = 2\theta$$
$$y(\theta) = (1-\theta)(z+z) = (1-\theta)2z$$

# **3** Question 3

(Comprehensive Exam June 2007 Q4)

### 3.1 Part a

Walrasian Equilibrium requires:

1. Consumers maximize utility:

$$\max_{x_A, y_A} \min \{2x_A, y_A\} \text{ such that } p_x x_A + p_y y_A = p_x + p_y$$
$$\max_{x_B, y_B} \min \{x_B, 2y_B\} \text{ such that } p_x x_B + p_y y_B = p_x + p_y$$
$$\max_{x_C, y_C} \min \{2x_C, y_C\} \text{ such that } p_x x_C + p_y y_C = p_x + p_y$$

2. Markets clear:

$$x_A + x_B + x_C = 3$$
$$y_A + y_B + y_C = 3$$

Maybe not normalize  $p_y = 1$  this time:

set  $2x_A = y_A$  and  $x_B = 2y_B$  and  $2x_C = y_C$ 

Substitute them into the budget constraints:

$$x_A = \frac{p_x + p_y}{p_x + 2p_y}$$
$$y_A = \frac{2(p_x + p_y)}{p_x + 2p_y}$$
$$x_B = \frac{2(p_x + p_y)}{2p_x + p_y}$$
$$y_B = \frac{p_x + p_y}{2p_x + p_y}$$
$$x_C = \frac{p_x + p_y}{p_x + 2p_y}$$
$$y_C = \frac{2(p_x + p_y)}{p_x + 2p_y}$$

Check X market clearing:

$$\frac{p_x + p_y}{p_x + 2p_y} + \frac{2(p_x + p_y)}{2p_x + p_y} + \frac{p_x + p_y}{p_x + 2p_y} = 3$$
$$\Rightarrow \frac{2p_y}{p_x + 2p_y} = \frac{p_y}{2p_x + p_y}$$
$$\Rightarrow p_y = 0 \text{ or } 4p_x + 2p_y = p_x + 2p_y$$
$$\Rightarrow p_y = 0 \text{ or } p_x = 0$$

AND check Y market clearing:

$$\frac{2(p_x + p_y)}{p_x + 2p_y} + \frac{p_x + p_y}{2p_x + p_y} + \frac{2(p_x + p_y)}{p_x + 2p_y} = 3$$
$$\Rightarrow \frac{2p_x}{p_x + 2p_y} = \frac{p_x}{p_x + 2p_y}$$
$$\Rightarrow p_x = 0 \text{ or } 2p_x + 4p_y = p_x + 2p_y$$
$$\Rightarrow p_x = 0$$

Note that  $p_y = 0$  is NOT a solution here.

Therefore, WE has allocations:

$$x_A = \frac{1}{2}, x_B = 2, x_C = \frac{1}{2}$$
  
 $y_A = 1, y_B = 1, y_C = 1$ 

### 3.2 Part b

Core allocation requires:

$$\min\left\{2 \cdot \frac{1+\varepsilon}{2}, 1+\varepsilon\right\} \ge \min\left\{2x_A, y_A\right\}$$
$$\min\left\{2-\varepsilon, 2 \cdot (1-2\varepsilon)\right) \ge \min\left\{x_B, 2y_B\right\}$$
$$\min\left\{2 \cdot \frac{1+\varepsilon}{2}, 1+\varepsilon\right\} \ge \min\left\{2x_C, y_C\right\}$$

1. Not blocked by  $\left\{ A\right\} ,\left\{ B\right\} ,\left\{ C\right\} :$ 

$$x_A = y_A = 1$$
$$x_B = y_B = 1$$
$$x_C = y_C = 1$$

2. Not blocked by  $\left\{ A,B\right\} ,\left\{ A,C\right\} ,\left\{ B,C\right\} :$ 

$$x_A + x_B = y_A + y_B = 2$$
$$x_A + x_C = y_A + y_C = 2$$
$$x_B + x_C = y_B + y_C = 2$$

3. Not blocked by  $\{A, B, C\}$ :

$$x_A + x_B + x_C = y_A + y_B + y_C = 3$$

First simplify the conditions:

$$1 + \varepsilon \ge \min \{2x_A, y_A\}$$
  
$$2 - 4\varepsilon \ge \min \{x_B, 2y_B\} \text{ this assumes } \varepsilon \ge 0$$
  
$$1 + \varepsilon \ge \min \{2x_C, y_C\}$$

One person coalition:

$$1 + \varepsilon \ge 1$$
 and  $2 - 4\varepsilon \ge 1$   
 $\Rightarrow 0 \le \varepsilon \le \frac{1}{4}$ 

Two people coalition:

$$2 - 4\varepsilon \ge \min\left\{2 - \frac{1+\varepsilon}{2}, 2 \cdot (2 - (1+\varepsilon))\right\}$$
$$\Rightarrow 2 - 4\varepsilon \ge \min\left\{\frac{3}{2} - \frac{\varepsilon}{2}, 2 - 2\varepsilon\right\}$$
$$\Rightarrow 2 - 4\varepsilon \ge \frac{3}{2} - \frac{\varepsilon}{2}\left(\text{ since } \varepsilon \le \frac{1}{4}\right)$$
$$\Rightarrow \frac{1}{2} \ge \frac{7}{2}\varepsilon$$
$$\Rightarrow \varepsilon \le \frac{1}{7}$$

Three people coalition:

It is Pareto optimal.

Therefore, any  $\varepsilon \in \left[0, \frac{1}{7}\right]$  implies the allocation is in the Core.

# 4 Question 4

(Comprehensive Exam August 2008 Q4)

### 4.1 Part a

Let  $c_a^b$  represent the *a*-th period consumption of consumer *b*.

Core allocation requires:

1. Not blocked by  $\{1\}, \{2\}, ..., \{N\}$ :

$$\log (c_n^n) + \log (c_{n+1}^n) \ge \log (\theta) + \log (\theta') \text{ for } n > 0$$
$$\log (c_1^0) \ge \log (\theta')$$

$$\begin{split} c_{N+1}^{N} &\leq \theta' \Rightarrow c_{N}^{N} < \theta \text{ are blocked by } \{N\} \\ \Rightarrow c_{N}^{N} &\geq \theta \Rightarrow c_{N}^{N-1} \leq \theta' \Rightarrow c_{N-1}^{N-1} < \theta \text{ are blocked by } \{N-1\} \\ \Rightarrow c_{N-1}^{N-1} &\geq \theta \Rightarrow c_{N-1}^{N-2} \leq \theta' \Rightarrow c_{N-2}^{N-2} < \theta \text{ are blocked by } \{N-2\} \\ \Rightarrow \dots \\ \Rightarrow c_{2}^{2} &\geq \theta \Rightarrow c_{2}^{1} \leq \theta' \Rightarrow c_{1}^{1} < \theta \text{ are blocked by } \{1\} \end{split}$$

Now,

$$c_1^0 < \theta' \text{ are blocked by } \{0\}$$
  

$$\Rightarrow c_1^1 \le \theta \Rightarrow c_1^1 = \theta \Rightarrow c_2^1 < \theta' \text{ are blocked by } \{1\} \Rightarrow c_2^1 = \theta'$$
  

$$\Rightarrow c_2^2 \le \theta \Rightarrow c_2^2 = \theta \Rightarrow c_3^2 < \theta' \text{ are blocked by } \{2\} \Rightarrow c_3^2 = \theta'$$
  

$$\Rightarrow \dots$$
  

$$\Rightarrow c_N^N = \theta$$

Therefore, only consuming endowment is in the Core.

#### 4.2 Part b

Walrasian equilibrium requires:

1. Consumers maximize utility:

$$\max_{c_n^n, c_{n+1}^n} \log (c_n^n) + \log (c_{n+1}^n) \text{ such that } .p_n c_n^n + p_{n+1} c_{n+1}^n = p_n \theta + p_{n+1} \theta'$$
$$\max_{c_1^0} \log (c_1^0) \text{ such that } p_0 c_1^0 = p_0 \theta'$$

2. Markets clear:

$$c_n^{n-1} + c_n^n = \theta + \theta'$$
$$c_{N+1}^N = \theta'$$

Consider any sequence of prices  $\{p_n\}_{n=0}^N$ :

Start from consumer 0, her budget constraint implies:

 $c_1^0=\theta'$ 

Then market clearing condition of  $c_1$  implies:

$$c_1^1 = \theta + \theta' - c_1^0 = \theta + \theta' - \theta' = \theta$$

Then consumer 1's budget constraint implies:

$$c_2^1 = \theta'$$

Then market clearing condition of  $c_2$  implies:

$$c_2^2 = \theta + \theta' - c_2^1 = \theta + \theta' - \theta' = \theta$$
  
...  
$$c_N^N = \theta$$
  
$$V_{N+1} = \theta'$$

#### 4.3 Part c

The argument in Part b is forward induction. Therefore it applies to  $N \to \infty.$ 

 $c_N^N$ 

#### 4.4 Part d

The argument in Part a is backward induction. Therefore it does not apply.

Consider a coalition of  $\{0, 1, 2, ...\}$ :

The allocation with:

$$c_n^n = c_n^{n-1} = \frac{\theta + \theta'}{2}$$
 for  $n \ge 0$ 

is preferred by all consumers since:

$$\log\left(\frac{\theta+\theta'}{2}\right) > \log\left(\theta'\right) \text{ for consumer } 0$$
$$2\log\left(\frac{\theta+\theta'}{2}\right) > \log\left(\theta\right) + \log\left(\theta'\right) \text{ for consumers } n > 0$$

by concavity of log.

## **5** Question 5

(Comprehensive Exam June 2004 Q4)

#### 5.1 Part a

Z is blocked by coalition  $\{A_1,A_2,B_1\},$  need to check:

$$u_{A1}(x_{A1}, y_{A1}) \ge u_{A1}(Z_A^x, Z_A^y)$$
$$u_{A2}(x_{A2}, y_{A2}) \ge u_{A2}(Z_A^x, Z_A^y)$$
$$u_{B1}(x_{B1}, y_{B1}) \ge u_{B1}(Z_B^x, Z_B^y)$$

with

$$x_{A1} + x_{A2} + x_{B1} = 2\omega_A^x + \omega_B^x$$
$$y_{A1} + y_{A2} + y_{B1} = 2\omega_A^y + \omega_B^y$$

Consider:

$$x_{A1} = x_{A2} = \frac{1}{2} \left( \omega_A^x + Z_A^x \right)$$
$$y_{A1} = y_{A2} = \frac{1}{2} \left( \omega_A^y + Z_A^y \right)$$
$$x_{B1} = Z_B^x$$
$$y_{B1} = Z_B^y$$

where

 $(x_{A1}, y_{A1})$  and  $(x_{A2}, y_{A2})$  are located at the midpoint between  $\omega$  and Z (between Z and B), strictly preferred for A1 and A2.

 $(x_{B1}, y_{B1})$  is located at Z, indifferent for B1.

and

$$x_{A1} + x_{A2} + x_{B1}$$
$$= \omega_A^x + Z_A^x + Z_B^x$$
$$= \omega_A^x + \omega_A^x + \omega_B^x$$
$$y_{A1} + y_{A2} + y_{B1}$$
$$= \omega_A^y + Z_A^y + Z_B^y$$

#### 5.2 Part b

Suppose there are  ${\cal N}$  players of each type, then:

Z is blocked by coalition  $\{A_1, A_2, ..., A_N, B_1, B_2, ..., B_{N-1}\}$ , need to check:

$$u_{Ai}(x_{Ai}, y_{Ai}) \ge u_{Ai}(Z_A^x, Z_A^y) \text{ for } i \in \{1, 2, ..., N\}$$
$$u_{Bj}(x_{Bj}, y_{Bj}) \ge u_{Bj}(Z_B^x, Z_B^y) \text{ for } j \in \{1, 2, ..., N-1\}$$

with

$$\sum_{i=1}^{N} x_{Ai} + \sum_{j=1}^{N-1} x_{Bj} = N\omega_A^x + (N-1)\omega_B^x$$
$$\sum_{i=1}^{N} y_{Ai} + \sum_{j=1}^{N-1} y_{Bj} = N\omega_A^y + (N-1)\omega_B^y$$

Consider:

$$\begin{aligned} x_{Ai} &= \frac{1}{N}\omega_A^x + \frac{N-1}{N}Z_A^x \text{ and} \\ y_{Ai} &= \frac{1}{N}\omega_A^y + \frac{N-1}{N}Z_A^y \text{ for } i \in \{1, 2, ..., N\} \\ x_{Bj} &= Z_B^x \text{ and} \\ y_{Bj} &= Z_B^y \text{ for } j \in \{1, 2, ..., N-1\} \end{aligned}$$

where

 $(x_{Ai}, y_{Aj})$ s are located at some point between  $\omega$  and Z for large N (between Z and B), strictly preferred for  $A_i$ .

 $(x_{Bj}, y_{Bj})$  is located at Z, indifferent for Bj. and

$$\sum_{i=1}^{N} x_{Ai} + \sum_{j=1}^{N-1} x_{Bj}$$
  
=  $\omega_A^x + (N-1) (Z_A^x + Z_B^x)$   
=  $\omega_A^x + (N-1) (\omega_A^x + \omega_B^x)$   
=  $N\omega_A^x + (N-1) \omega_B^x$   
 $\sum_{i=1}^{N} y_{Ai} + \sum_{j=1}^{N-1} y_{Bj}$   
=  $\omega_A^y + (N-1) (Z_A^y + Z_B^y)$   
=  $\omega_A^y + (N-1) (\omega_A^y + \omega_B^y)$   
=  $N\omega_A^y + (N-1) \omega_B^y$