

# CS540 Introduction to Artificial Intelligence

## Lecture 15

Young Wu

Based on lecture slides by Jerry Zhu and Yingyu Liang

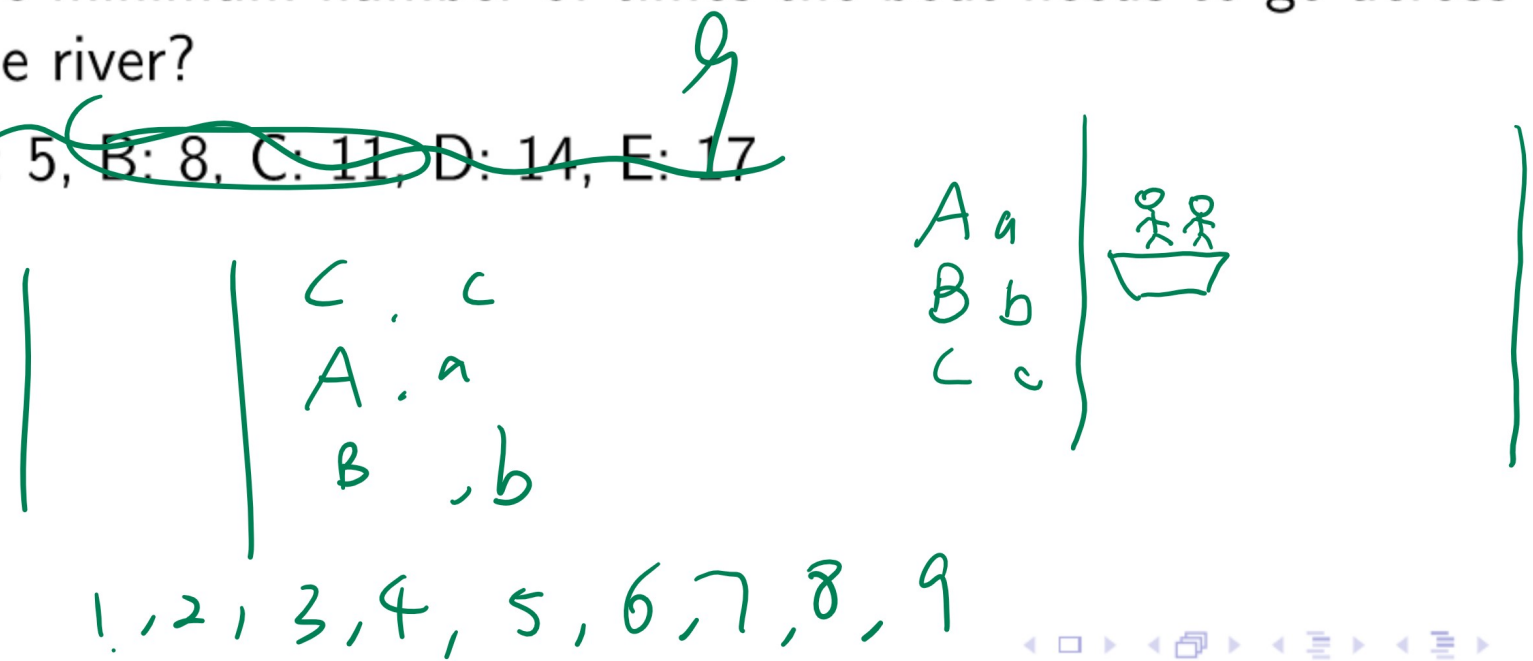
July 15, 2019

# River Corssing Problem

## Quiz (Participation)

- Three married couples need to cross the river. The boat holds no more than two people. No woman can be in the presence of another man unless her husband is also present. What is the minimum number of times the boat needs to go across the river?

- ~~A: 5, B: 8, C: 11, D: 14, E: 17~~





# Search Problem Applications

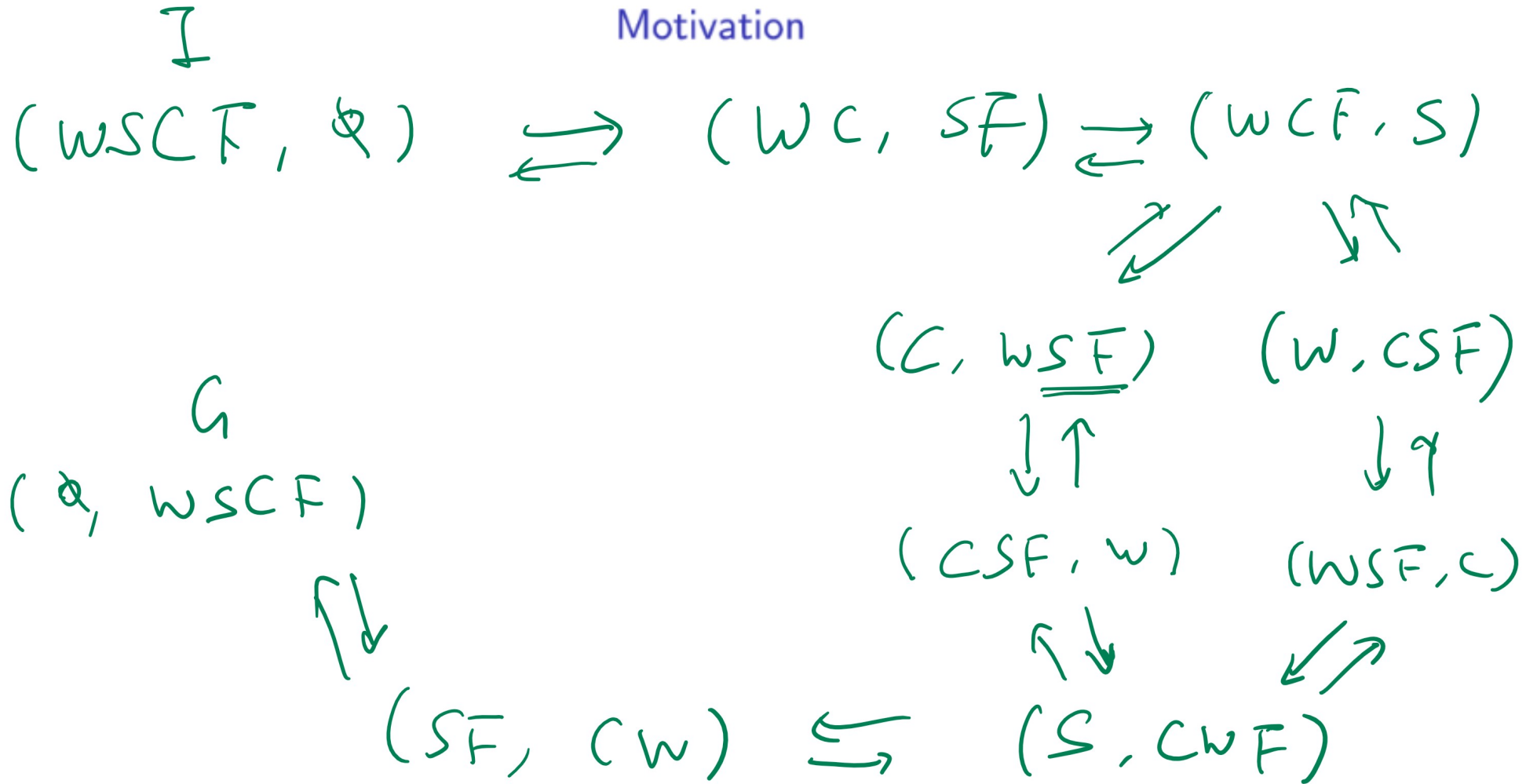
## Motivation

- Puzzles and games.
- Navigation: route finding.
- Motion planning.
- Scheduling.



# Wolf, Sheep, Cabbage Example

Motivation



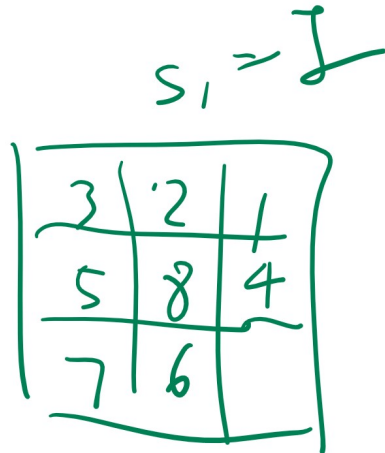


# State Space

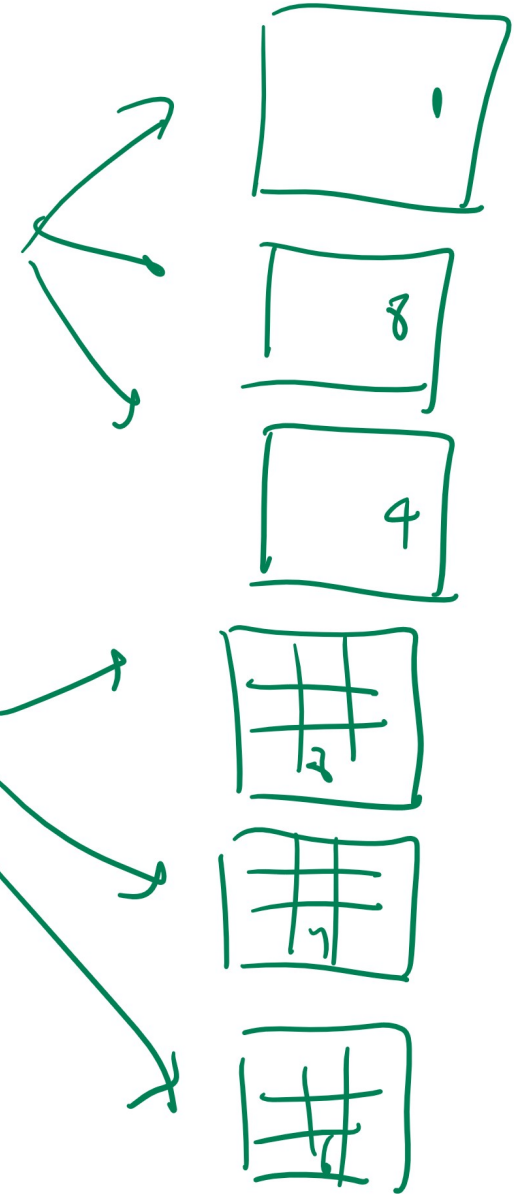
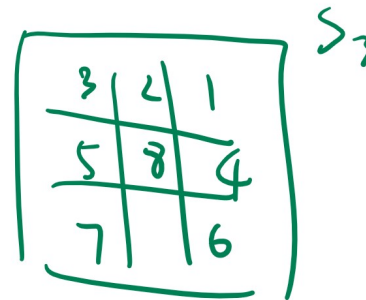
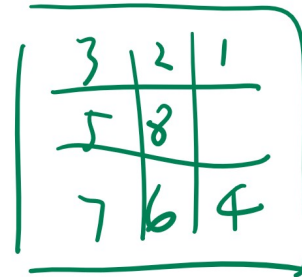
## Motivation

- The states need to represent all necessary information about the game.
- The actions are discrete and deterministic and are determined by the successor function.
- Each possible action at state  $s$  is associated with a state in the set  $s'(s)$ .

# 8 Puzzle Example



Motivation



# Sizes of State Space

## Motivation

- Tic Tac Toe:  $10^3$
- Checkers:  $10^{20}$
- Chess:  $10^{50}$
- Go:  $10^{170}$

# State Space Graph

## Definition

- A state space can be represented by a weighted directed graph  $(V, E, c)$ .
- $V$  is the set of vertices (also called nodes).
- $E$  is the set of edges (also called arcs). Each edge is directed from one vertex to another vertex and represents an action.
- $c$  is the cost (also called weights) associated with each edge. The costs are positive.

# Search Problem on Graph

## Definition

- Search starts at an initial state and finishes if one of the goal states is reached.
- The solution is a path in the graph from an initial state to a goal state.
- The cost of a solution is the sum of edge costs on the solution path.
- The optimal solution is the solution with the lowest cost.

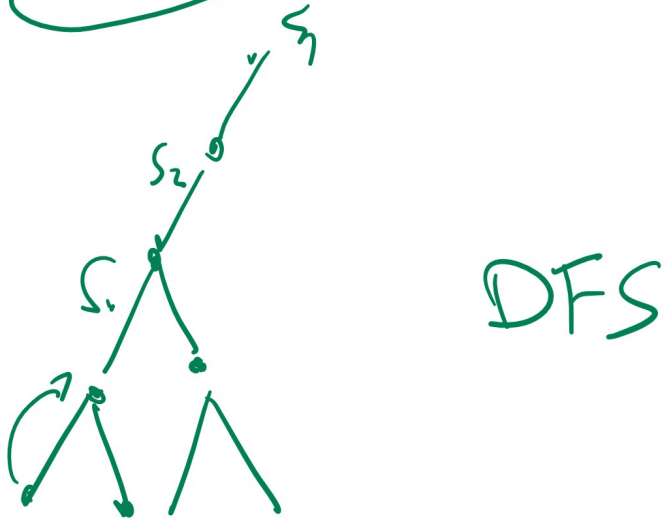
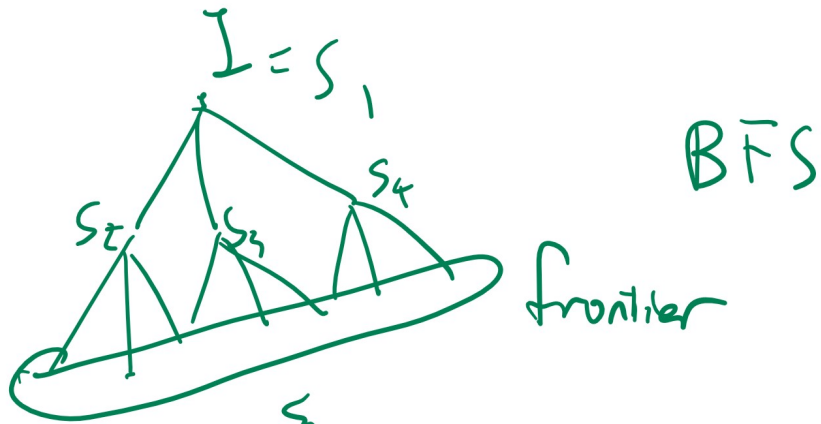






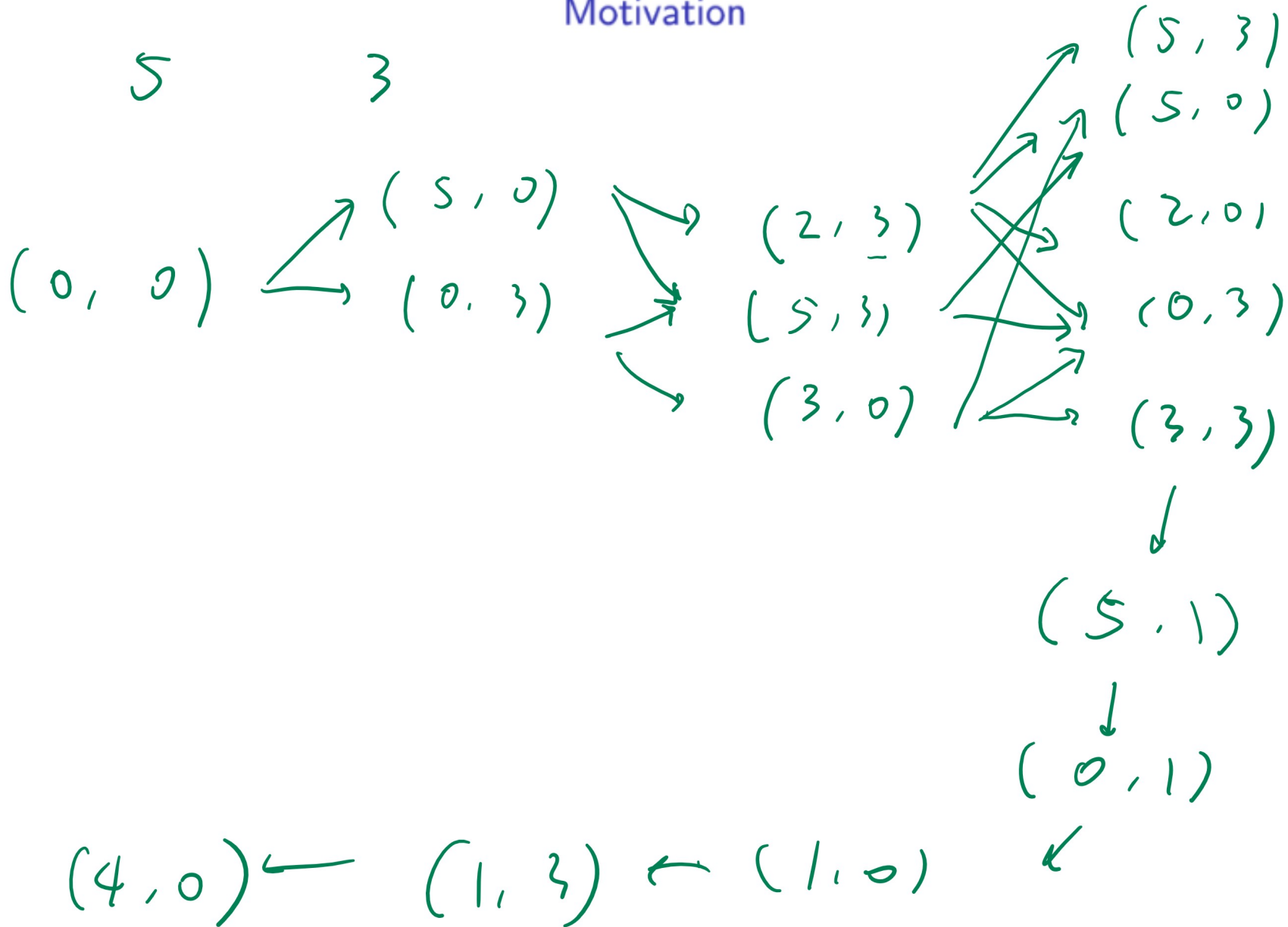
# Water Jugs Example, Part I

Motivation



# Water Jugs Example, Part II

Motivation



# Performance

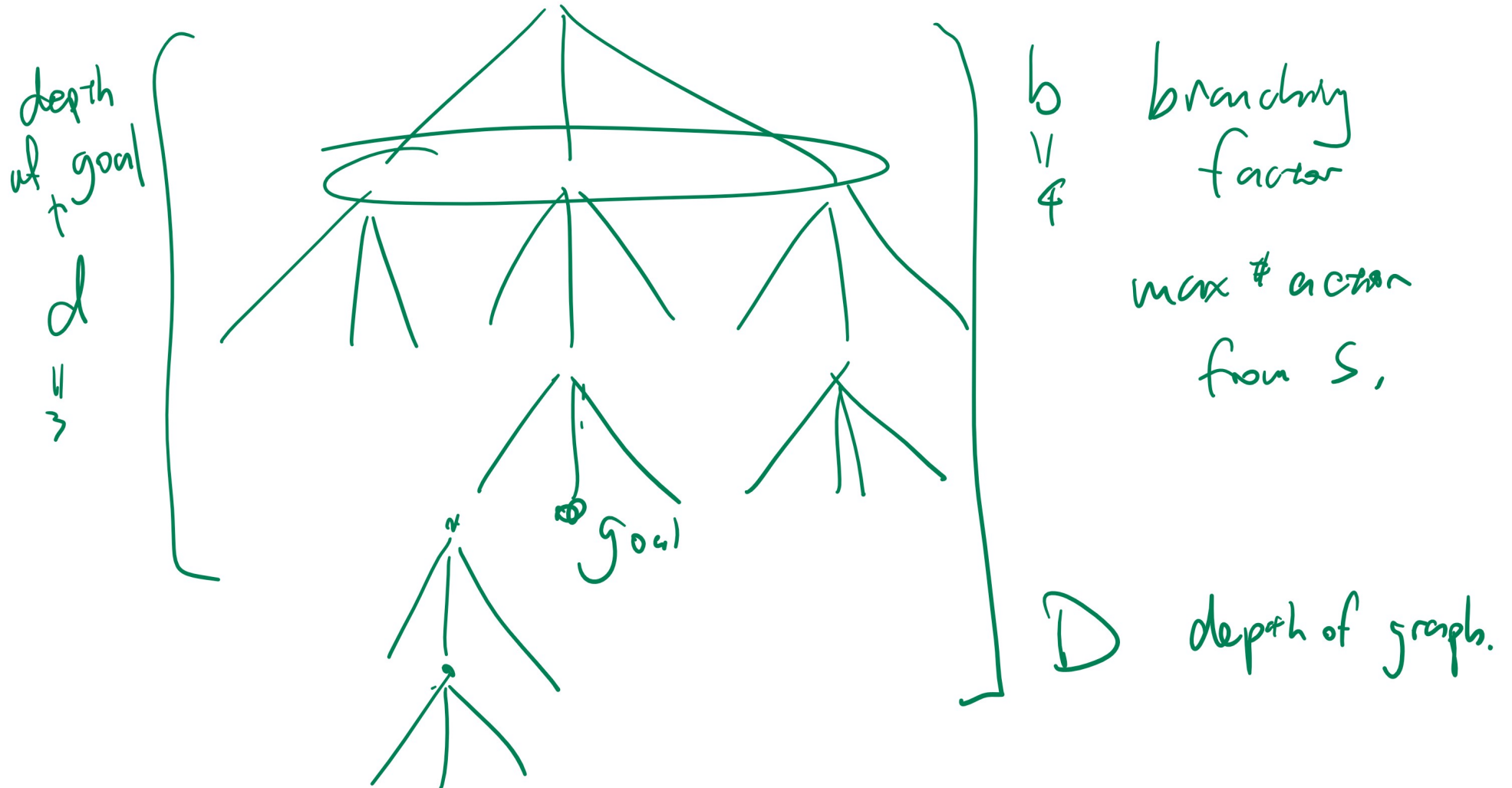
## Definition

- A search strategy is complete if it finds at least one solution.
- A search strategy is optimal if it finds the optimal solution.
- For uninformed search, the costs are assumed to be 1 for all edges  $c = 1$ .



# Search Tree Diagram

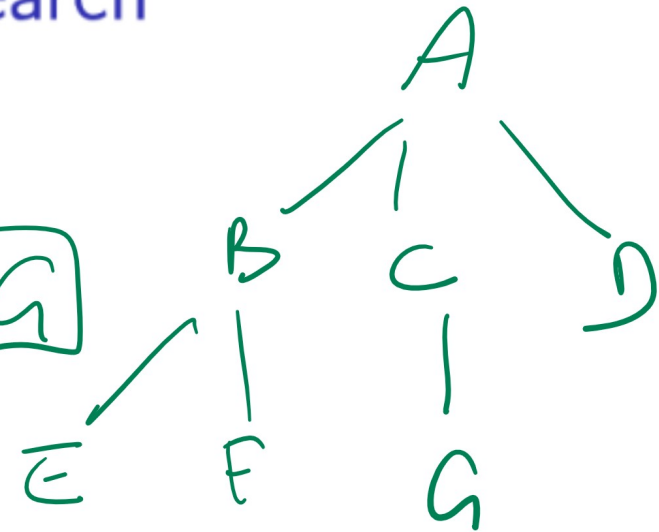
## Definition



# Breadth First Search

Description

Q = A B ~~C~~ D E F G



- Use Queue (FIFO) for the frontier.
- Remove from the front, add to the back.

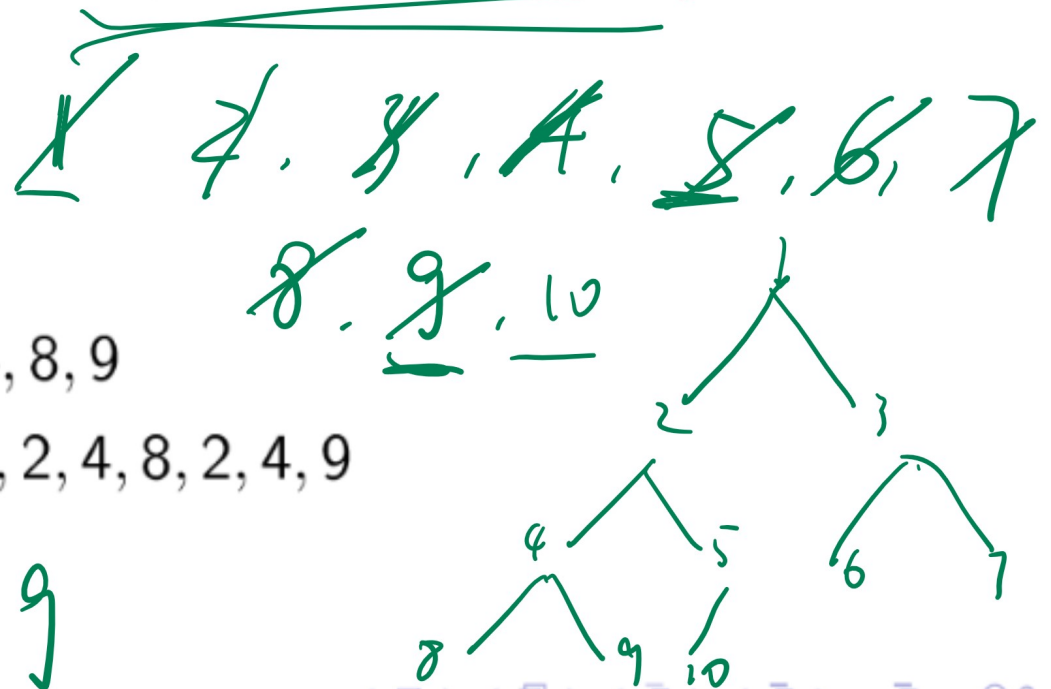
# BFS Example

## Quiz (Graded)



- Fall 2018 Midterm Q2, Fall 2017 Midterm Q13, Fall 2010 Final Q2
- Suppose the states are positive integers between 1 and 10, initial state is 1, goal state is 9, successors of  $i$  is  $2i$  and  $2i + 1$  (if exist). What a BFS expansion sequence?

- A: 1, 2, 3, 4, 5, 6, 7, 8, 9
- B: 1, 2, 4, 8, 3, 5, 7, 9
- C: 1, 2, 4, 8, 9
- D: 1, 2, 3, 2, 4, 5, 3, 6, 7, 2, 4, 8, 9
- E: 1, 2, 3, 2, 4, 2, 5, 3, 6, 3, 7, 2, 4, 8, 2, 4, 9



1, 2, 3, 4, 5, 6, 7, 8, 9









# Breadth First Search Performance

## Discussion

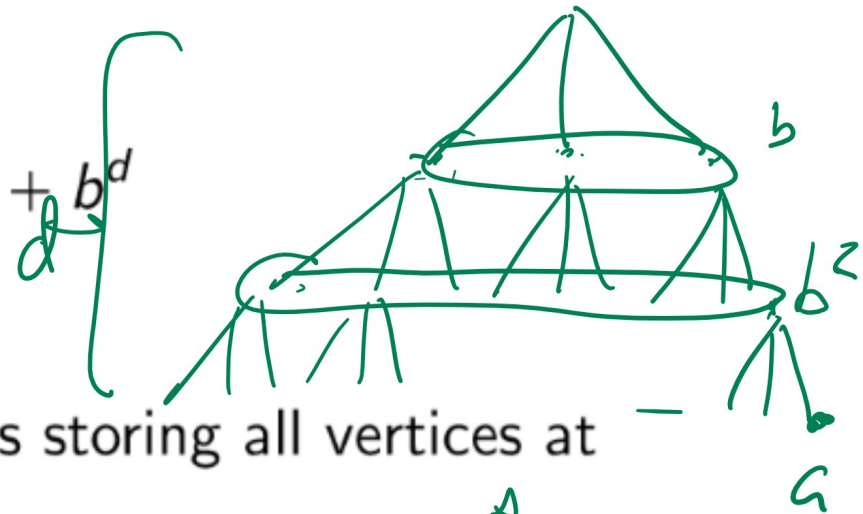
- BFS is complete.
- BFS is optimal with  $c = 1$ .

# Breadth First Search Complexity

## Discussion

- Time complexity: the worst case occurs when the goal is the last vertex at depth  $d$ .

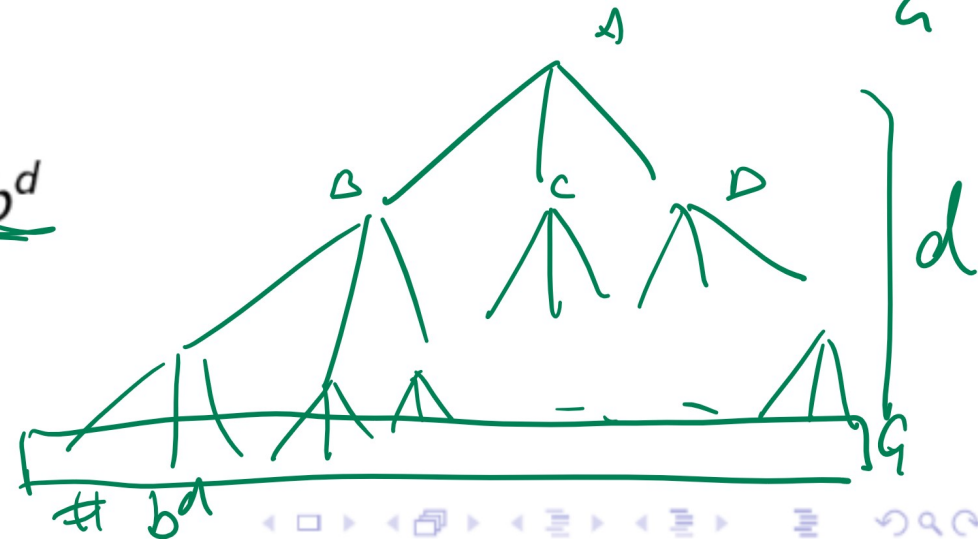
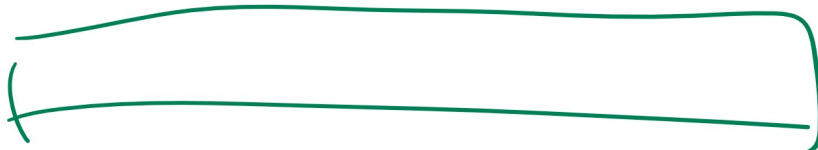
$$T = \underbrace{b} + \underbrace{b^2} + \dots + \underbrace{b^d}_d$$



- Space complexity: the worst case is storing all vertices at depth  $d$  is in the frontier.

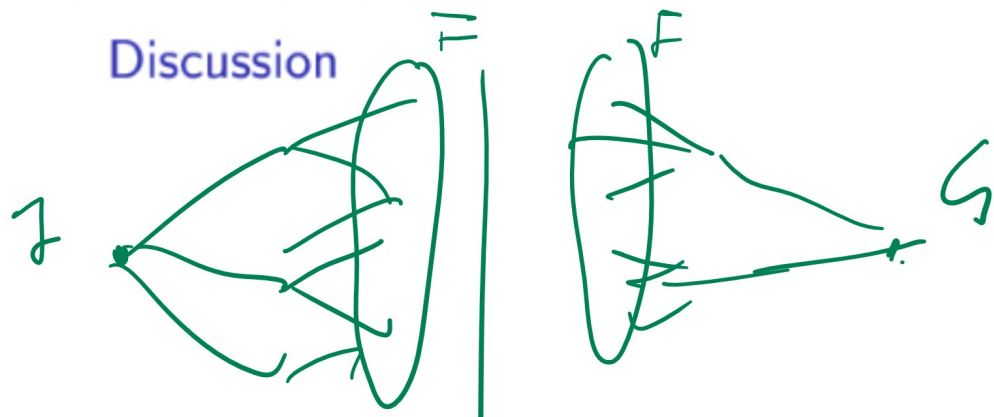
~~A~~ B C D

$$S = \underline{b^d}$$



# BiDirectional Search

Discussion



- BFS from the initial states and goal states at the same time.
- The search stops when the two frontiers meet (have non-empty intersection) in the middle.
- The time and space complexity is the same as BFS with depth  $\frac{d}{2}$ .

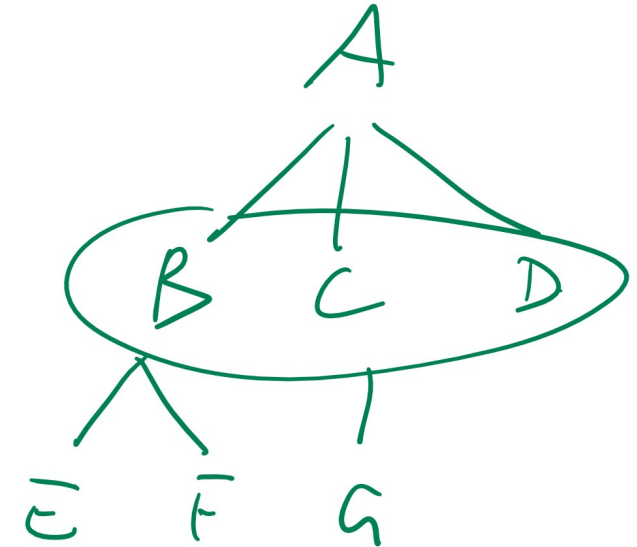
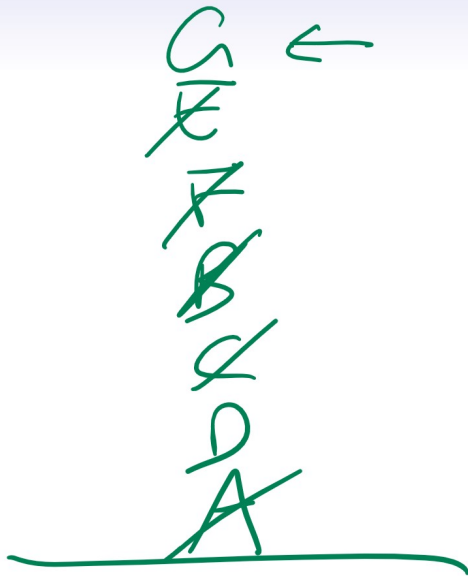
$$T = 2 \left( b + b^2 + \dots + b^{\frac{d}{2}} \right)$$

$$S = 2 b^{\frac{d}{2}}$$

$Q(I) \cap Q(G) \neq \emptyset \rightarrow \text{meet.}$

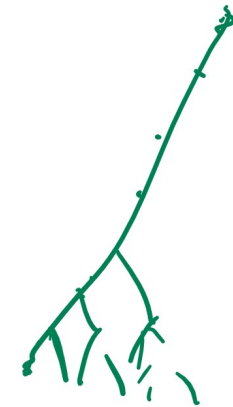
# Depth First Search

## Description



- Use Stack (LIFO) for the frontier.
- Remove from the front, add to the front.

A B E F C G  
list of expanded vertices.



# DFS Example

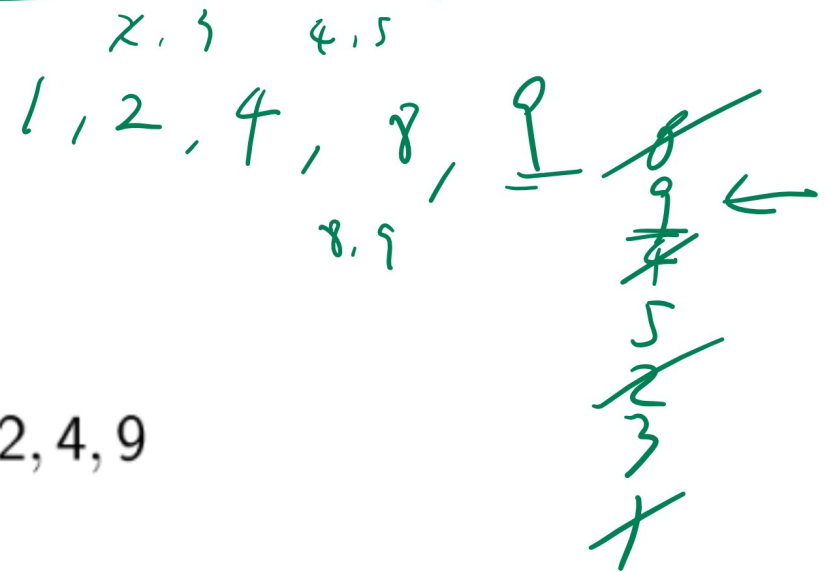
## Quiz (Graded)

- Fall 2018 Midterm Q2, Fall 2017 Midterm Q13, Fall 2010 Final Q2

0.5

- Suppose the states are positive integers between 1 and 10, initial state is 1, goal state is 9, successors of  $i$  is  $2i$  and  $2i + 1$  (if exist). What a DFS expansion sequence?

- A: 1, 2, 3, 4, 5, 6, 7, 8, 9
- B: 1, 2, 4, 8, 3, 5, 7, 9
- C: 1, 2, 4, 8, 9
- D: 1, 2, 3, 2, 4, 5, 3, 6, 7, 2, 4, 8, 9
- E: 1, 2, 3, 2, 4, 2, 5, 3, 6, 3, 7, 2, 4, 8, 2, 4, 9







# Depth First Search

## Algorithm

- Input: a weighted digraph  $(V, E, c)$ , initial states  $I$  and goal states  $G$ .
- Output: a path from  $I$  to  $G$ .
- Push initial states.

$$S = I$$

- While  $S$  is not empty and goal is not ~~found~~, pop  $S$  and push its successors.

*popped*

$$s = S_0$$

$$S = s'(s) + S$$



# Depth First Search Performance

## Discussion

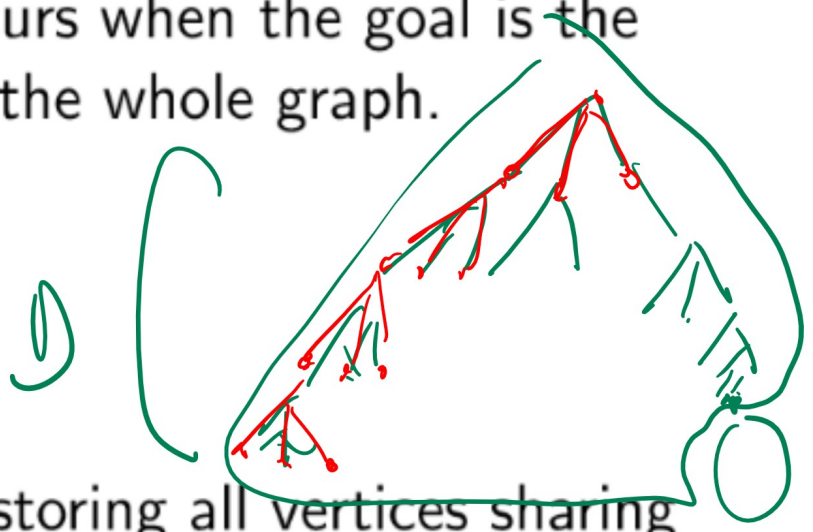
- DFS is incomplete if  $D = \infty$ .
- DFS is not optimal.

# Depth First Search Complexity

## Discussion

- Time complexity: the worst case occurs when the goal is the root of the last subtree expanded in the whole graph.

$$T \approx b^D$$



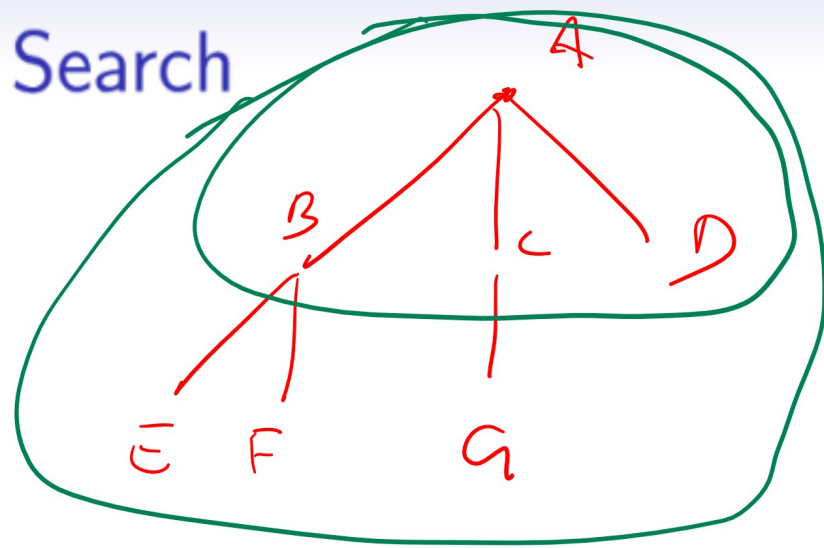
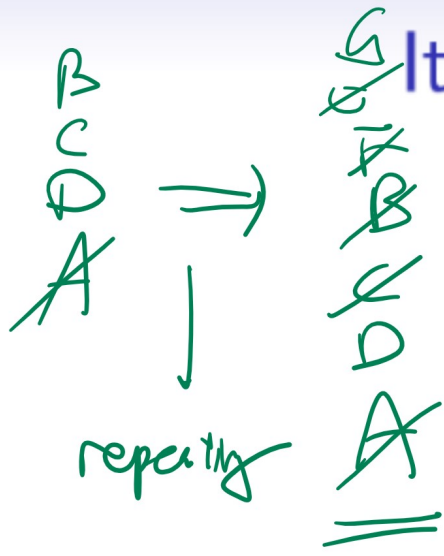
- Space complexity: the worst case is storing all vertices sharing the parents with vertices in the current path.

$$S = (b - 1) D$$

	BFS	DFS
Time	small	large
Space	large	small

# Iterative Deepening Search

Description

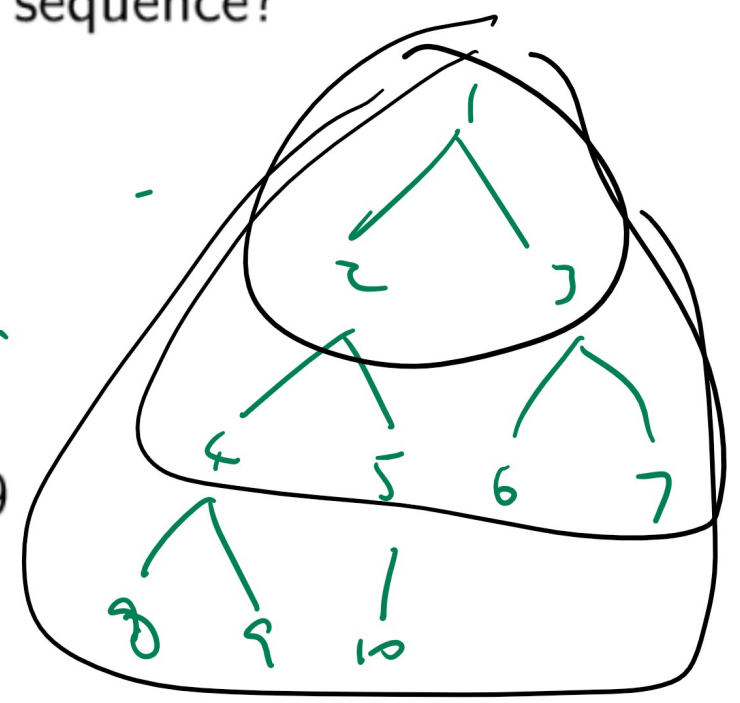


- DFS but stop if path length > 1
- repeat DFS but stop if path length > 2
- ...
- repeat DFS but stop if path length >  $d$

# IDS Example

## Quiz (Graded)

- Fall 2018 Midterm Q2, Fall 2017 Midterm Q13, Fall 2010 Final Q2
- Suppose the states are positive integers between 1 and 10, initial state is 1, goal state is 9, successors of  $i$  is  $2i$  and  $2i + 1$  (if exist). What a IDS expansion sequence?
- A: 1, 2, 3, 4, 5, 6, 7, 8, 9
- B: 1, 2, 4, 8, 3, 5, 7, 9
- C: 1, 2, 4, 8, 9
- D: 1, 2, 3, 2, 4, 5, 3, 6, 7, 2, 4, 8, 9
- E: 1, 2, 3, 2, 4, 2, 5, 3, 6, 3, 7, 2, 4, 8, 2, 4, 9



1, 2, 3, 1 2 4 5 3 6 7  
1, 2, 4, 8, 9





# Iterative Deepening Search Performance

## Discussion

- IDS is complete.
- IDS is optimal with  $c = 1$ .

# Iterative Deepening Search Complexity

## Discussion

- Time complexity: the worst case occurs when the goal is the last vertex at depth  $d$ .

$$T = db + (d - 1)b^2 + \dots + 3b^{d-1} + 2b^{d-1} + 1b^d \approx O(b^d)$$

same as  
DFS

- Space complexity: it has the same space complexity as DFS.

$$S = \underline{\underline{(b - 1) D}}$$





# Configuration Space

## Discussion



