

# CS540 Introduction to Artificial Intelligence

## Lecture 16

Young Wu

Based on lecture slides by Jerry Zhu, Yingyu Liang, and Charles Dyer

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# Uniformed vs. Informed Search

## Motivation

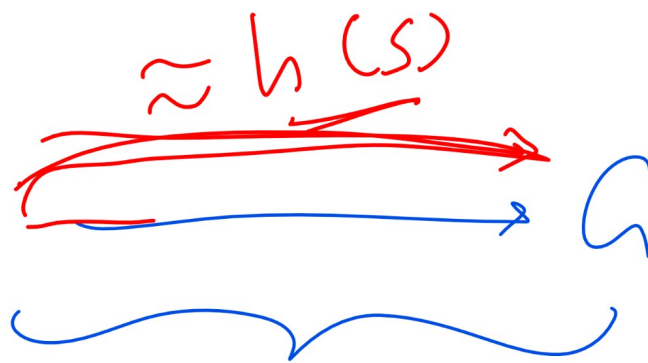
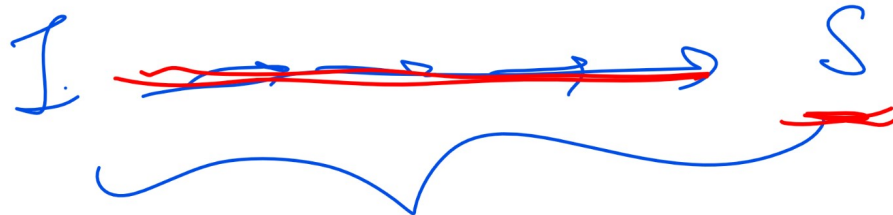
- Uninformed search means only the goal  $G$  and the successor functions  $s'$  are given.
- Informed search means which non-goal states are better is also known.

# Heuristic Diagram

## Motivation

~~$g + h$~~   $\approx$  total cost

heuristic approximation



cost so far =  $g(s)$

?  $h^*(s)$  -> remaining cost.





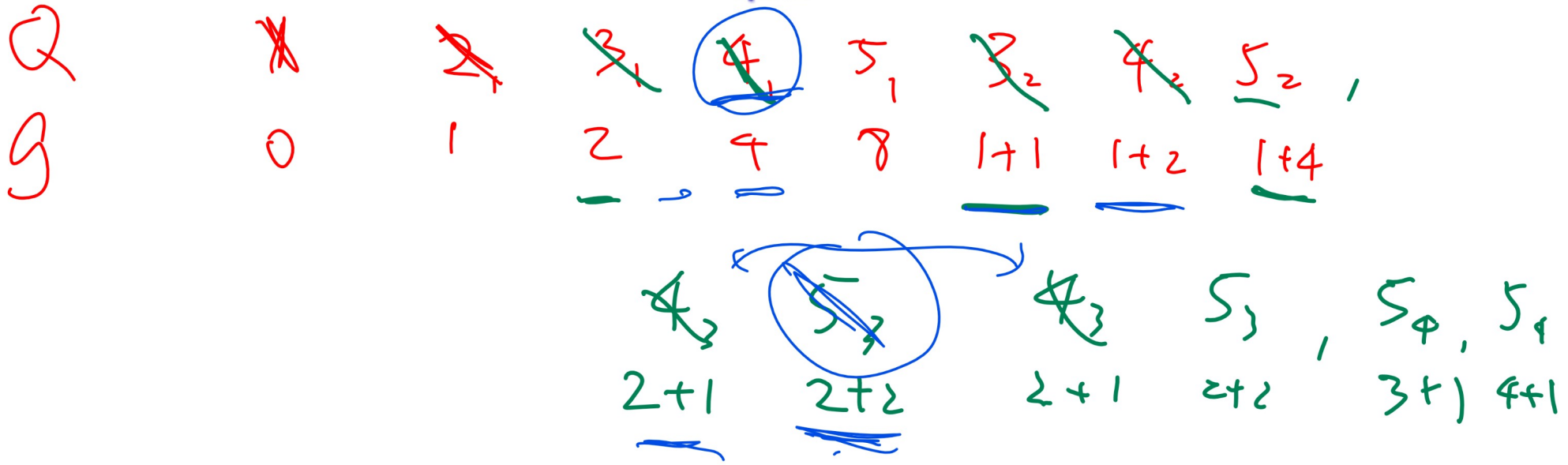






# UCS Example 2 Diagram

Quiz



record a video.



# Uniform Cost Search Performance

## Discussion

- UCS is complete.
- UCS is optimal with any  $c$ .



mg, mp

# Best First Greedy Search

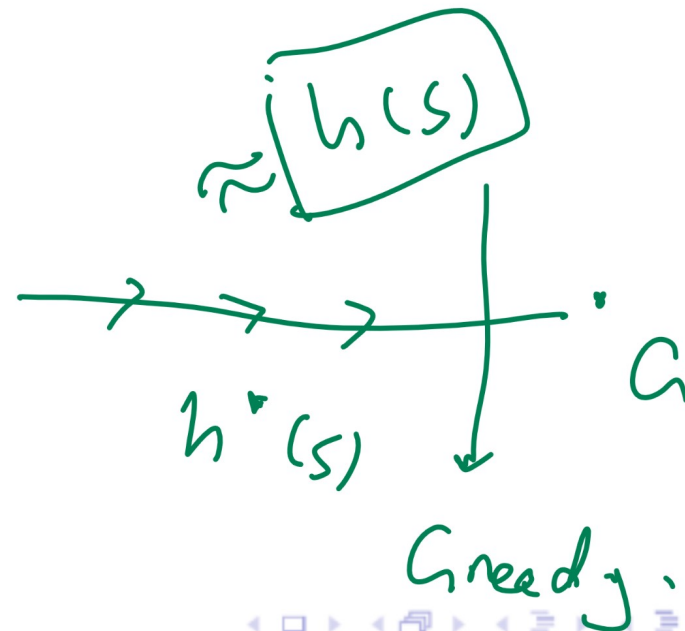
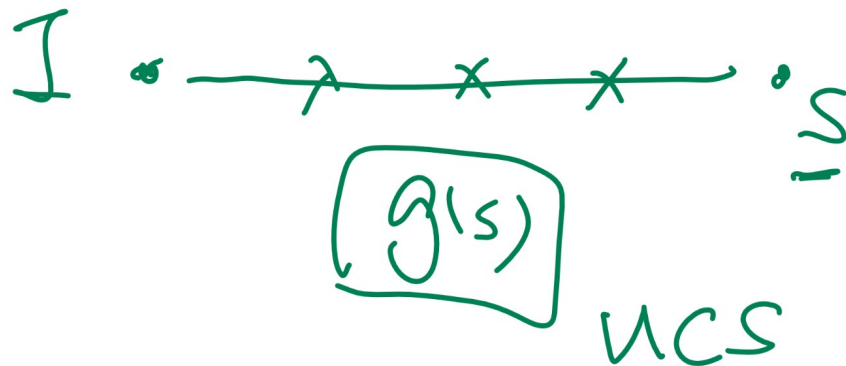
## Description

PS → bugs.

↳ want for SPT.

- Expand the vertices with the lowest heuristic cost  $h(s)$  first.
- Use a priority queue based on  $h(s)$ .

A  $g(s) + h(s)$



# Greedy Example 1

## Quiz

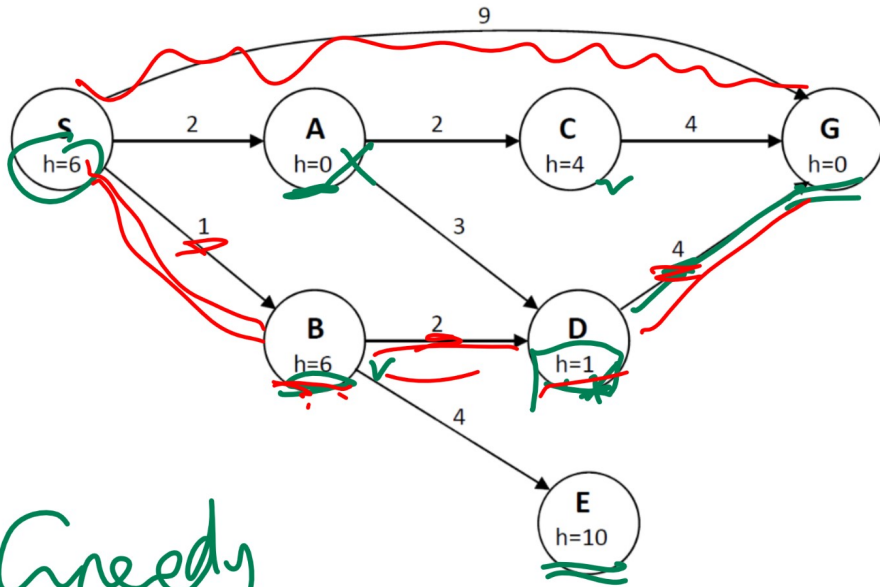
- Given the following adjacency matrix. Find Greedy Search expansion path.

—	S	A	B	C	D	E	G
S	$h = 6$	2	1	—	—	—	9
A	—	$h = 0$	—	2	3	—	—
B	—	—	$h = 6$	—	2	4	—
C	—	—	—	$h = 4$	—	—	4
D	—	—	—	—	$h = 1$	—	4
E	—	—	—	—	—	$h = 10$	—
G	—	—	—	—	—	—	$h = 0$



# Greedy Example 1 Diagram

Quiz



UCS

SBDG lowest cost

7

Greedy

Q:

~~A~~ S ~~C~~ D C A S

B<sub>s</sub>

h:

0 0 1 4 6

6

expansion.

S A G

solutions

~~SG~~

→ not a solution

# Greedy Example 2

## Quiz

Q1

- Given that cost from state  $i$  to  $j$  is  $2^{j-i-1}$  for  $j > i$ . The heuristic is  $h(i) = 5 - i$ , The initial state is 1 and goal state is 5. What is a vertex expansion sequence if Best First Greedy Search is used?

- A: 1, 5
- B: 1, 2, 3, 4, 5
- C: 1, 2, 3, 4, 4, 5
- D: 1, 2, 3, 3, 4, 4, 5
- ~~E: 1, 2, 3, 3, 4, 4, 4, 5~~

UCS

Q:

h:

↑

⊗

4

5

$h=4$   
 $= 5-1$

4

3

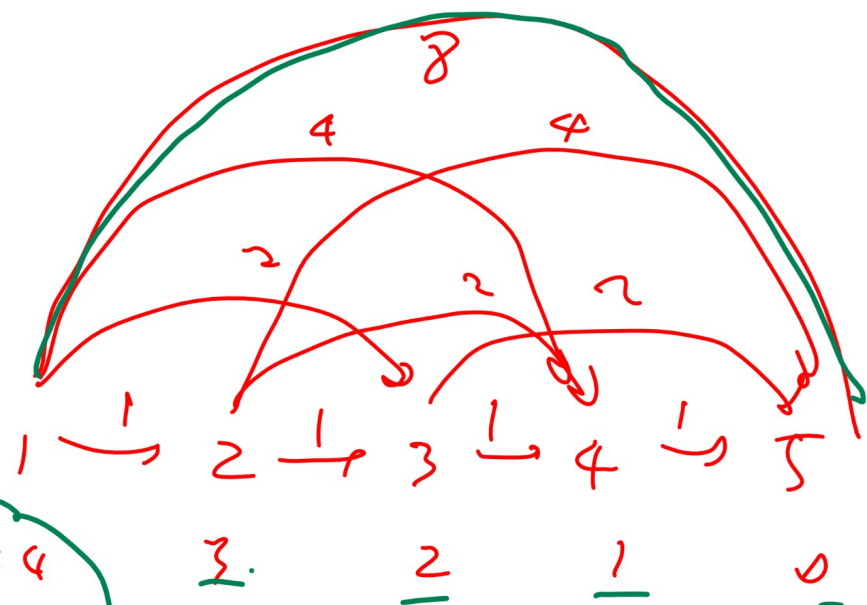
2

0

1

2

3



1, 5





# Best First Greedy Search Performance

## Discussion

- Greedy is incomplete.
- Greedy is not optimal.

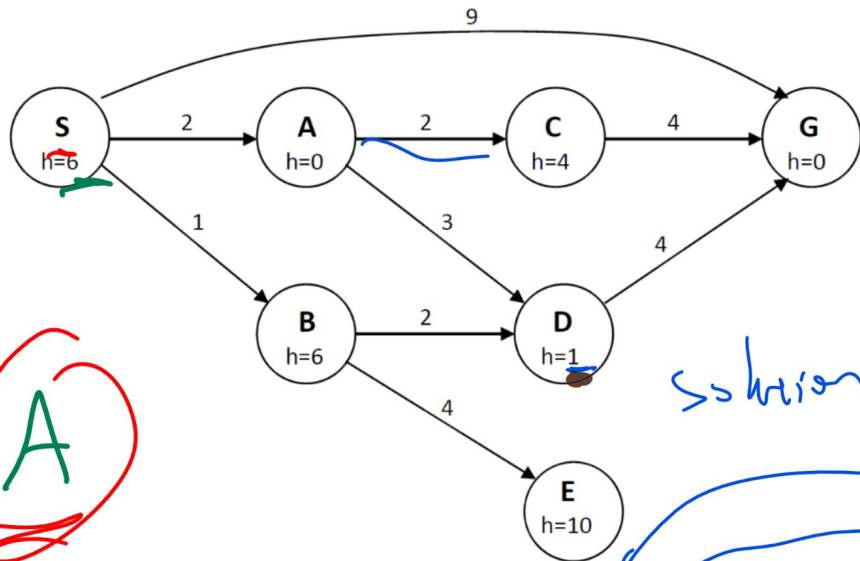






# A Search Example 1 Diagram

Quiz



A

expansion path: S, A, D, B  
D, G

Solution

SBDG

Q	<del>A<sub>S</sub></del>	<del>D<sub>B</sub></del>	<del>D<sub>A</sub></del>	<del>E<sub>C</sub></del>	<del>B<sub>D</sub></del>	<u>G<sub>D</sub></u>	C <sub>A</sub>	G <sub>S</sub>	G <sub>D</sub>	E <sub>B</sub>
g	<u>2</u>	1+2	<u>2</u> +3	0	<u>1</u>	3+4	<u>2</u> +2	9	5+4	1+9
h	0	1	1	6	6	0	4	0	0	10
g+h	2	4	6	6	7	7	8	9	9	15

back tracker

[S] A [B] C [D] G  
Initial S [S] A ~~[B]~~ ~~[D]~~



# A Search

## Algorithm

- Input: a weighted digraph  $(V, E, c)$ , initial states  $I$  and goal states  $G$ , and the heuristic function  $h(s), s \in V$ .
- Output: a path from  $I$  to  $G$ .
- EnQueue initial states into a priority queue  $Q$ . Here,  $Q$  is ordered by  $g(s) + h(s)$  for  $s \in Q$ .

$$Q = I$$

- While  $Q$  is not empty and goal is not deQueued, deQueue  $Q$  and enQueue its successors.

$$s = Q_{(0)} = \arg \min_{s \in Q} g(s) + h(s)$$

$$Q = Q + s'(s)$$

# A Search Performance

## Discussion

- A is complete. ✓
- A is not optimal. ✗

# A Star Search

## Description

- $A^*$  search is A search with an <sup>h</sup> ~~admissible~~ heuristic.

# Admissible Heuristic

## Definition

- A heuristic is admissible if it never over estimates the true cost.

$$0 \leq h(s) \leq h^*(s)$$

*Handwritten notes: The equation is enclosed in a red box. A red arrow points from the word "true" to  $h^*(s)$ . The terms  $h(s)$  and  $h^*(s)$  are underlined with red lines.*

*Handwritten notes: "A search  $h$ " is underlined with a red line. A red double-headed arrow points to "A\*", which is also underlined with a red line. The word "Optimal" is written to the right and underlined with a red line.*

# Dominated Heuristic

## Definition

- One heuristic,  $h_1$ , is dominated by another,  $h_2$ , if:

$$h_1(s) \leq h_2(s) \leq h^*(s), \forall s \in S$$

*in better  $h_1$*

- If  $h_2$  dominates  $h_1$ , then  $h_2$  is better than  $h_1$  since  $A^*$  using  $h_1$  expands at least as many states (or more) than  $A^*$  using  $h_2$ .
- If  $h_2$  dominated  $h_1$ ,  $A^*$  with  $h_2$  is better informed than  $A^*$  with  $h_1$ .

# Admissible Heuristic 8 Puzzle Example

## Quiz

$C = 1$

• Which ones (select multiple) of the following are admissible heuristic function for the 8 Puzzle?

$h(\text{goal}) = 0$

• A:  $h(s) =$  number of tiles in the wrong position.

admissible  
no good

• B:  $h(s) = 0.$  ✓

UCS

• ~~C:  $h(s) = 1.$~~

$h(\text{goal}) = 0 \neq 1$

worst possible admissible  $h.$

• D:  $h(s) =$  sum of Manhattan distance between each tile and its goal location.

$h^*$

✓ best

• E:  $h(s) =$  sum of Euclidean distance between each tile and its goal location. ✓

# Admissible Heuristic General Example 1

Quiz

$s \rightarrow$  any config of the game puzzle.

- Which ones (select multiple) of the following are admissible heuristic function?

A:  $h(s) = h^*(s) \leq h^*$

B:  $h(s) = \max\{2, h^*(s)\}$

C:  $h(s) = \min\{2, h^*(s)\} \leq h^*$

D:  $h(s) = h^*(s) - 2 \leq -2$

E:  $h(s) = \sqrt{h^*(s)}$

$h(\text{goal}) = 2 \neq 0$



$h^* > 1$

$h^* = 0.25$

$\sqrt{h^*} < h^*$

$\sqrt{h^*} = h = 0.5 > h^*$

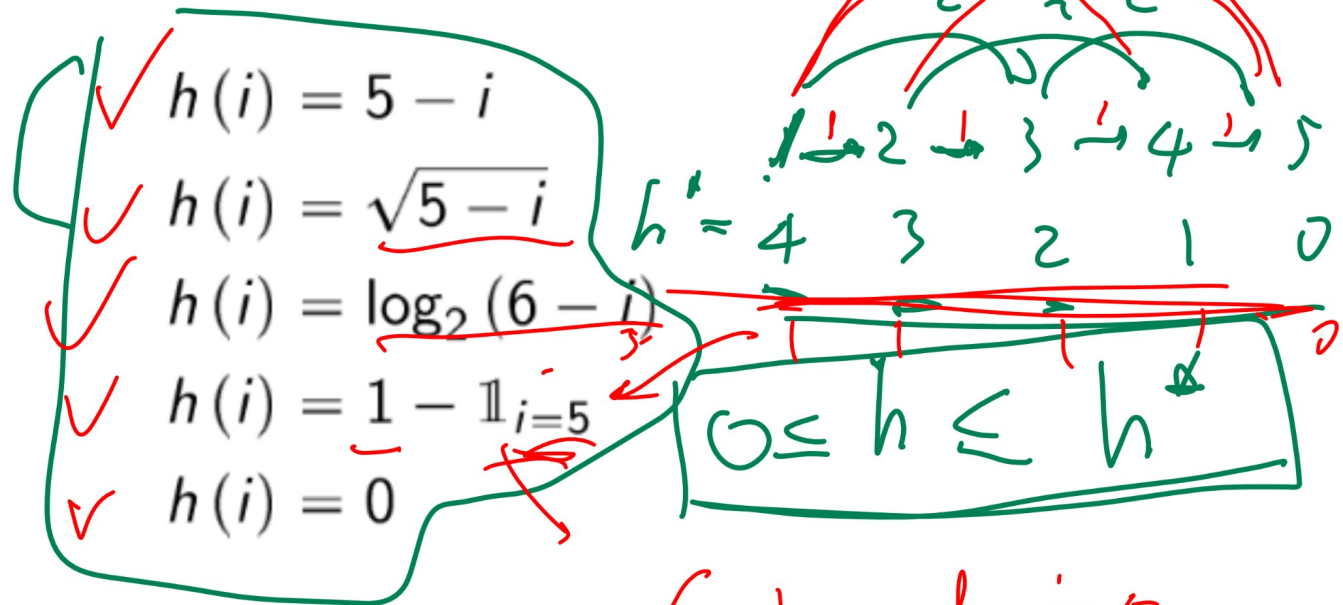


# A Star Search Example 2

## Quiz

- Given that cost from state  $i$  to  $j$  is  $2^{j-i-1}$  for  $j > i$ . How many of the following heuristic functions are admissible? For  $i = 1, 2, 3, 4, 5$ :

Q2



- A: 1, B: 2, C: 3, D: 4, E: 5



# A Star Search Example 2

## Quiz

Q3 (1.5c), A

- Given that cost from state  $i$  to  $j$  is  $2^{j-i-1}$  for  $j > i$ . Which one of the following heuristic functions is not dominated (among the admissible ones)? For  $i = 1, 2, 3, 4, 5$ :

- A:  $h(i) = 5 - i = h^*$
- B:  $h(i) = \sqrt{5 - i}$
- C:  $h(i) = \log_2(6 - i)$
- D:  $h(i) = 1 - \mathbb{1}_{i=5}$
- E:  $h(i) = 0$

E is dom by D.

$0 \leq h_1 \leq h_2 \leq h^*$   
 $h_1$  dominated by  $h_2$   
 $h_2$  dominates  $h_1$

# A Star Search with Revisit, Part I

## Algorithm

- Input: a weighted digraph  $(V, E, c)$ , initial states  $I$  and goal states  $G$ , and the heuristic function  $h(s), s \in V$ .
- Output: a path with minimum cost from  $I$  to  $G$ .
- EnQueue initial states into a priority queue  $Q$ . Here,  $Q$  is ordered by  $g(s) + h(s)$  for  $s \in Q$ .

$$Q = I$$

$$g(I) = 0$$

$$g(s) = \infty, \text{ for } s \notin I$$

- Initialize the list of visited vertices,  $P$ .

$$P = \emptyset$$

# A Star Search with Revisit, Part II

## Algorithm

- While  $Q$  is not empty and goal is not deQueued, deQueue  $Q$ , put it on  $P$  and enQueue its successors to  $Q$ , and update the cost functions.

$$s = Q_{(0)} = \arg \min_{s \in Q} g(s) + h(s)$$

$$P = P + s$$

$$Q = Q + s'(s), \text{ update } g(s') = \min \{g(s'), g(s) + c(s, s')\}$$



# Iterative Deepening A Star Search

## Discussion

- $A^*$  can use a lot of memory.
- Do path checking without expanding any vertex with  $g(s) + h(s) > 1$ .
- Do path checking without expanding any vertex with  $g(s) + h(s) > 2$ .
- ...
- Do path checking without expanding any vertex with  $g(s) + h(s) > d$ .



# Beam Search

## Discussion



- Version 1: Keep a priority queue with fixed size  $k$ . Only keep the top  $k$  vertices and discard the rest.
- Version 2: Only keep the vertices that are at most  $\epsilon$  worse than the best vertex in the queue.  $\epsilon$  is called the beam width.

