

# CS540 Introduction to Artificial Intelligence

## Lecture 16

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Based on lecture slides by Jerry Zhu and Yingyu Liang

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# Bridge and Torch Game, Part I

## Quiz (Participation)

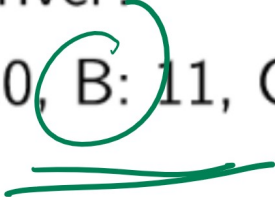
- Four people with one flashlight (torch) want to go across a river. The bridge can hold two people at a time, and they must cross with the flashlight. The time it takes for each person to cross the river:

A	B	C	D
<u>1</u>	<u>2</u>	3	4

$$2 + 1 + 4 + 2 + 2$$

$$2 + 1 + 3 + 1 + 4$$

- What is the minimum total time required for everyone to cross the river?
- A: 10, B: 11, C: 12, D: 13, E: 14



# Bridge and Torch Game, Part II

## Quiz (Participation)

- Four people with one flashlight (torch) want to go across a river. The bridge can hold two people at a time, and they must cross with the flashlight. The time it takes for each person to cross the river:

A	B	C	D
1	2	4	5

$$\begin{array}{r} 2 + 1 + 5 + 2 + 2 \\ \hline 12 \end{array}$$

- What is the minimum total time required for everyone to cross the river?

- A: 10, B: 11, C: 12, D: 13, E: 14

CD AB

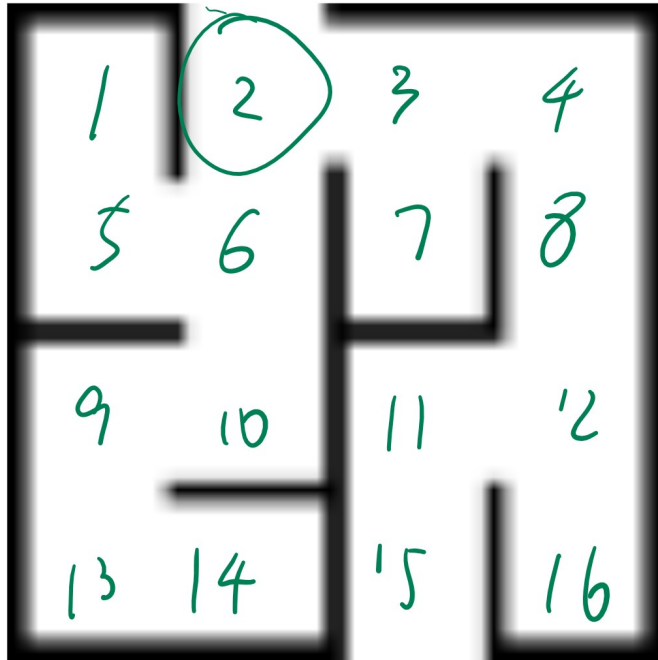
~~$$2 + 1 + 4 + 1 + 5 = 13$$~~





# Maze Example, Part I

## Motivation



BFS

~~2~~ ~~3~~ ~~6~~ ~~4~~ 7 5 10  
2  
 8

DFS

2 3 6 | 2 3 6 4 7  
 5 10  
 depth 1 tree      depth 2 tree



# Uniformed vs. Informed Search

## Motivation

- Uninformed search means only the goal  $G$  and the successor functions  $s'$  are given.
- Informed search means which non-goal states are better is also known.

# Heuristic

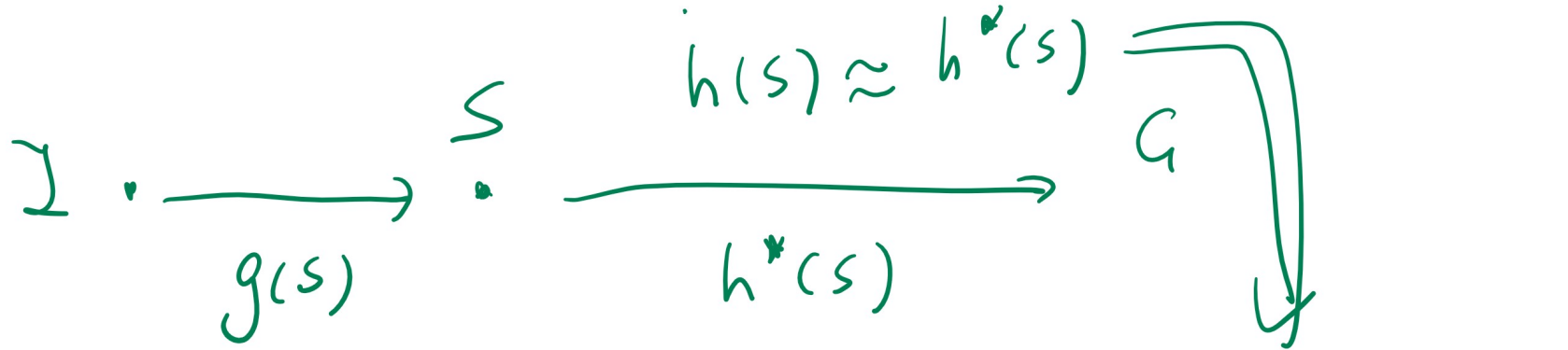
## Motivation

- The additional information is usually given as a heuristic cost from a state  $s$  to the goal.
- The cost of the path from the start to a vertex  $s$  in the frontier is  $g(s)$ .
- The cost from  $s$  to the goal,  $h^*(s)$ , is estimated by  $h(s)$ . This estimate may not be accurate.

$$h(s) \approx h^*(s)$$

# Heuristic Diagram

## Motivation



① expand  $s$  with smallest  $g(s)$  first  
 ↳ UCS (Dijkstra)  
 uniform cost search

② —————  $h(s)$  first  
 ↳ best first greedy search

③ —————  $g(s) + h(s)$  first  $\Rightarrow$  A search.



# UCS Example, Part I

Quiz (Graded)

- Spring 2017 Midterm Q1
- Given the following adjacency matrix. Find UCS expansion path.

—	S	A	B	C	D	E	G
S	$h = 6$	2	1	—	—	—	9
A	—	$h = 0$	—	2	3	—	—
B	—	—	$h = 6$	—	2	4	—
C	—	—	—	$h = 4$	—	—	4
D	—	—	—	—	$h = 1$	—	4
E	—	—	—	—	—	$h = 10$	—
G	—	—	—	—	—	—	$h = 0$



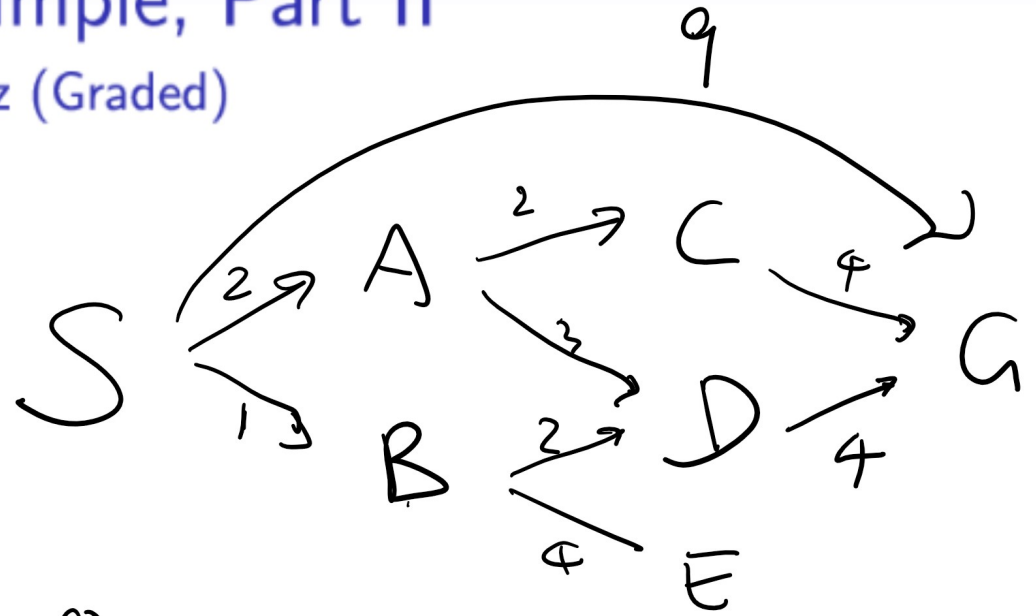
# UCS Example, Part II

Quiz (Graded)

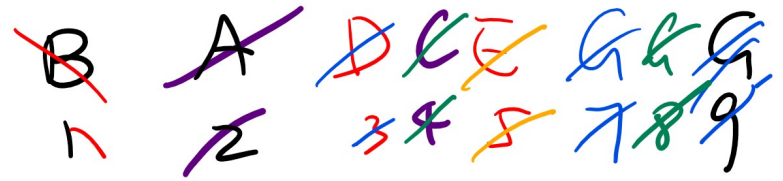
Q4

expansion order

- A: S B A D C E G
- B: S B D G
- C: S A G
- D: S G
- E: S A D B D G



off Q



expansion



# Uniform Cost Search

## Algorithm

- Input: a weighted digraph  $(V, E, c)$ , initial states  $I$  and goal states  $G$ .
- Output: a path from  $I$  to  $G$ .
- EnQueue initial states into a priority queue  $Q$ . Here,  $Q$  is ordered by  $g(s)$  for  $s \in Q$ .

$$Q = I$$

- While  $Q$  is not empty and goal is not deQueued, deQueue  $Q$  and enQueue its successors.

$$s = Q_{(0)} = \arg \min_{s \in Q} g(s)$$

$$Q = Q + s'(s)$$

# Uniform Cost Search Performance

## Discussion

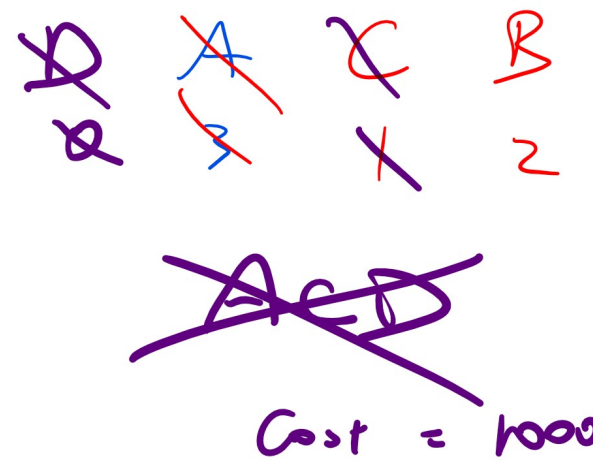
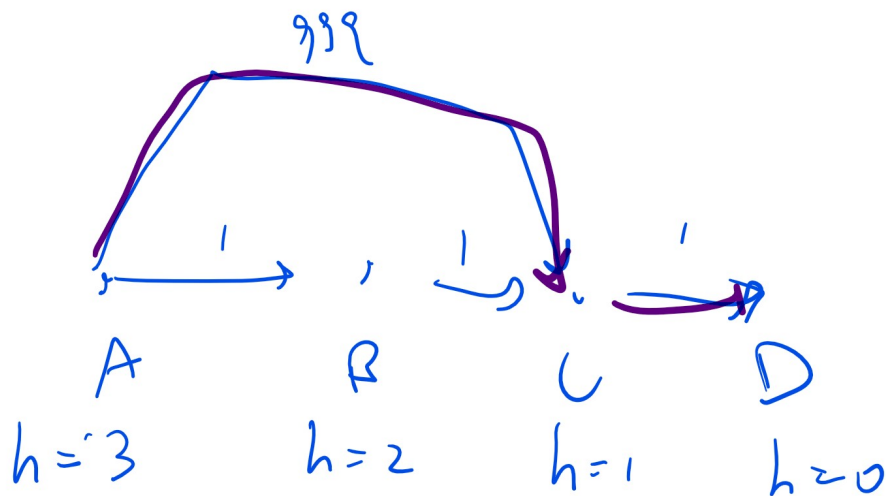
- UCS is complete.
- UCS is optimal with any  $c. \geq 0$

no heuristic

# Best First Greedy Search

## Description

- Expand the vertices with the lowest heuristic cost  $h(s)$  first.
- Use a priority queue based on  $h(s)$ .



# Greedy Example, Part I

## Quiz (Graded)

- Given the following adjacency matrix. Find Greedy Search expansion path.

—	S	A	B	C	D	E	G
S	$h = 6$	2	1	—	—	—	9
A	—	$h = 0$	—	2	3	—	—
B	—	—	$h = 6$	—	2	4	—
C	—	—	—	$h = 4$	—	—	4
D	—	—	—	—	$h = 1$	—	4
E	—	—	—	—	—	$h = 10$	—
G	—	—	—	—	—	—	$h = 0$

# Greedy Example, Part II

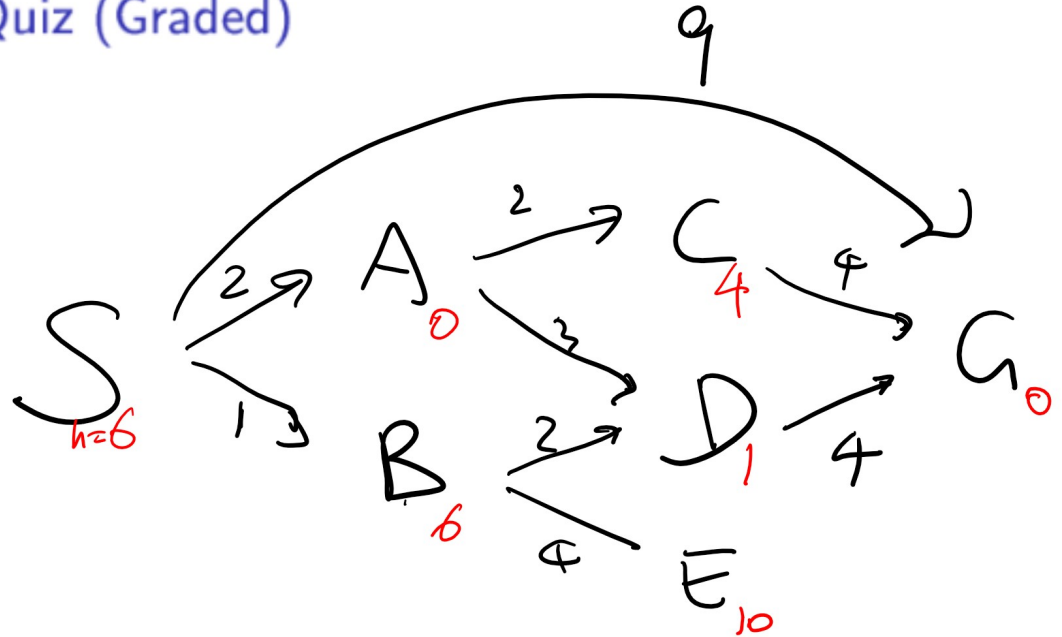
Quiz (Graded)

Q6

expensive path

put on final

- A: S B A D C E G
- B: S B D G
- C: S A G
- ~~D: S G~~
- E: S A D B D G



$h^*(D) = 4$

$h(D) = 1$

$h(s) =$

<del>6</del>	<del>0</del>	<del>6</del>	1	4	6
S	A	B	D	C	E

# Best First Greedy Search

## Algorithm

- Input: a weighted digraph  $(V, E, c)$ , initial states  $I$  and goal states  $G$ , and the heuristic function  $h(s), s \in V$ .
- Output: a path from  $I$  to  $G$ .
- EnQueue initial states into a priority queue  $Q$ . Here,  $Q$  is ordered by  $h(s)$  for  $s \in Q$ .

$$Q = I$$

- While  $Q$  is not empty and goal is not deQueued, deQueue  $Q$  and enQueue its successors.

$$s = Q_{(0)} = \arg \min_{s \in Q} h(s)$$

$$Q = Q + s'(s)$$









# A Search Example, Part II

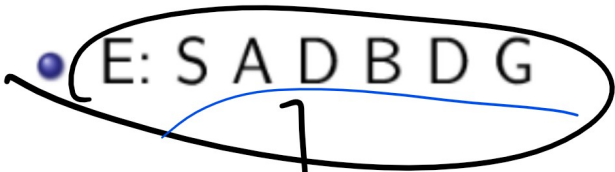
## Quiz (Graded)

expansion path

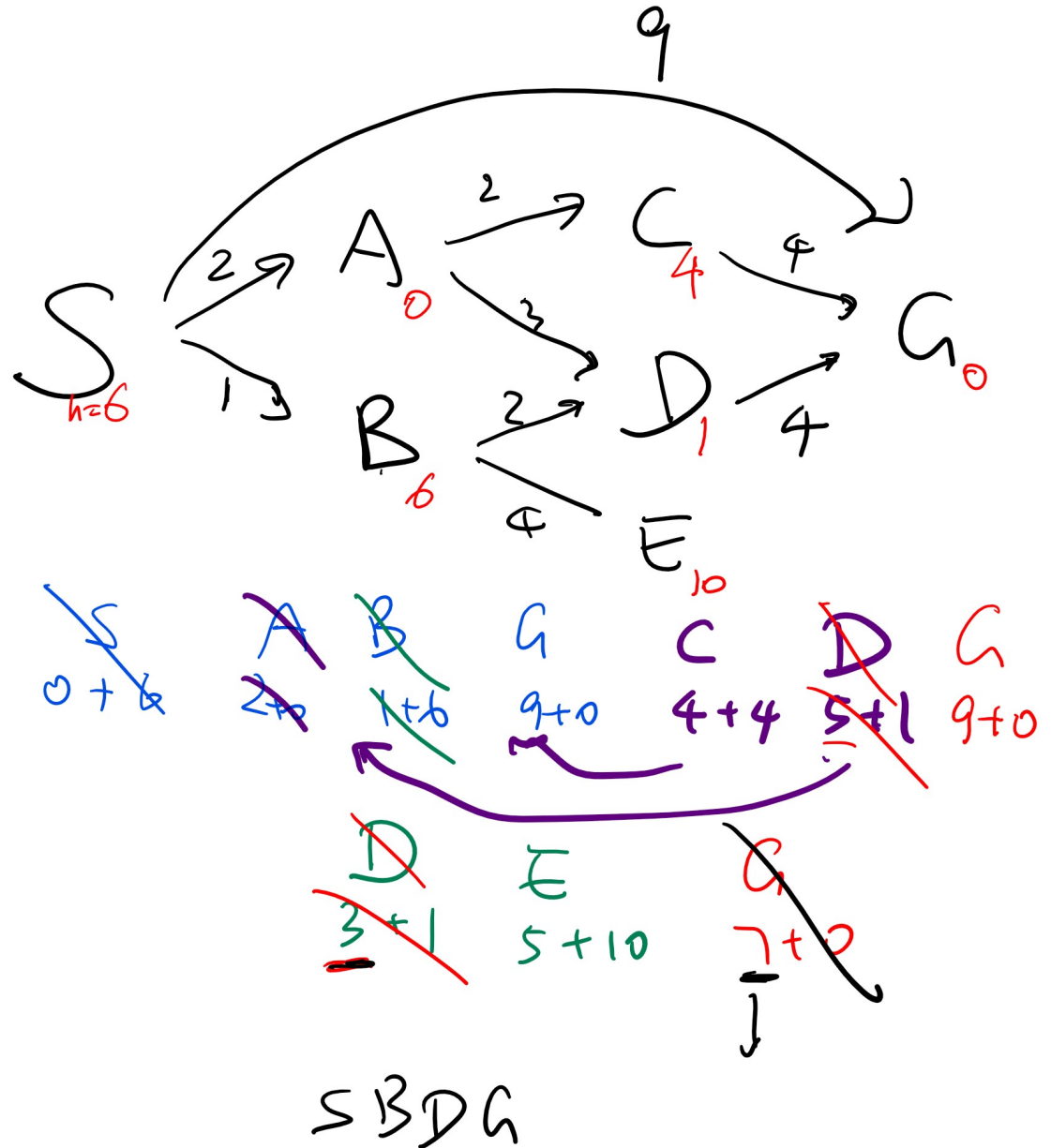
(Q7)

A search.

- A: S B A D C E G
- B: S B D G
- C: S A G
- D: S G
- E: S A D B D G



S A D B D G



# A Search

## Algorithm

- Input: a weighted digraph  $(V, E, c)$ , initial states  $I$  and goal states  $G$ , and the heuristic function  $h(s), s \in V$ .
- Output: a path from  $I$  to  $G$ .
- EnQueue initial states into a priority queue  $Q$ . Here,  $Q$  is ordered by  $g(s) + h(s)$  for  $s \in Q$ .

$$Q = I$$

- While  $Q$  is not empty and goal is not deQueued, deQueue  $Q$  and enQueue its successors.

$$s = Q_{(0)} = \arg \min_{s \in Q} g(s) + h(s)$$

$$Q = Q + s'(s)$$

# A Search Performance

## Discussion

- A is complete.
- A is not optimal.

# A Star Search

## Description

- $A^*$  search is A search with an admissible heuristic.

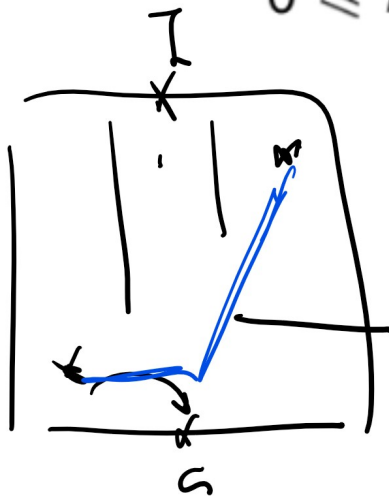


# Admissible Heuristic

## Definition

- A heuristic is admissible if it never over estimates the true cost.

$$0 \leq h(s) \leq h^*(s)$$



$h(s) = \text{Euclidean distance} \leq$   
actual cost

# Admissible Heuristic 8 Puzzle Example

## Quiz (Graded)

Q10

Which ones (select multiple) of the following are admissible heuristic function for the 8 Puzzle?

A:  $h(s)$  = number of tiles in the wrong position.

B:  $h(s) = 0$ .

C:  $h(s) = 1$ .  $\Rightarrow h(G) = 0$

D:  $h(s)$  = sum of Manhattan distance between each tile and its goal location.

E:  $h(s)$  = sum of Euclidean distance between each tile and its goal location.

$\leq h^*(s)$

1	2	3
4	5	6
7	8	

C=1

$0 \leq h(s) \leq h^*(s)$

$\leq h^*(s)$

$\leq h_0(s)$

usually requires more steps

# Admissible Heuristic General Example

## Quiz (Graded)

- Which ones (select multiple) of the following are admissible heuristic function?

- A:  $h(s) = h^*(s)$ . ✓
- B:  ~~$h(s) = \max\{2, h^*(s)\}$~~
- C:  $h(s) = \min\{2, h^*(s)\}$ . ✓
- D:  ~~$h(s) = h^*(s) - 2$~~   $\leq 0$
- E:  $h(s) = \sqrt{h^*(s)}$ . ✓  
0

# Dominated Heuristic

## Definition

- One heuristic,  $h_1$ , is dominated by another,  $h_2$ , if:

$$h_1(s) \leq h_2(s) \leq h^*(s), \forall s \in S$$

- If  $h_2$  dominates  $h_1$ , then  $h_2$  is better than  $h_1$  since  $A^*$  using  $h_1$  expands at least as many states (or more) than  $A^*$  using  $h_2$ .
- If  $h_2$  dominated  $h_1$ ,  $A^*$  with  $h_2$  is better informed than  $A^*$  with  $h_1$ .

# Non-Optimal Heuristic

## Definition

- If optimality is not required and a satisfying solution is acceptable, then the heuristic should be as close as possible, either under or over, to the actual cost.
- This results in fewer states being expanded compared to using poor but admissible heuristics.

Informed Search

○○○○○○○○

UCS

○○○○○

Greedy

○○○○○

A

○○○○○○○○○○○○●○○○○○○○○

# A Star Search Maze Example

Quiz (Graded)

# A Star Search with Revisit Example, Part I

## Quiz (Graded)

- Given the following adjacency matrix. Find A\* Search expansion path.

—	S	A	B	C	D	E	G
S	$h = 6$	2	1	—	—	—	9
A	—	$h = 0$	—	2	3	—	—
B	—	—	$h = 6$	—	2	4	—
C	—	—	—	$h = 4$	—	—	4
D	—	—	—	—	$h = 1$	—	4
E	—	—	—	—	—	$h = 10$	—
G	—	—	—	—	—	—	$h = 0$



# A Star Search with Revisit Example, Part II

## Quiz (Graded)

- A: S B A D C E G
- B: S B D G
- C: S A G
- D: S G
- E: S A D B D G

# A Star Search with Revisit, Part I

## Algorithm

- Input: a weighted digraph  $(V, E, c)$ , initial states  $I$  and goal states  $G$ , and the heuristic function  $h(s), s \in V$ .
- Output: a path with minimum cost from  $I$  to  $G$ .
- EnQueue initial states into a priority queue  $Q$ . Here,  $Q$  is ordered by  $g(s) + h(s)$  for  $s \in Q$ .

$$Q = I$$

$$g(I) = 0$$

$$g(s) = \infty, \text{ for } s \notin I$$

- Initialize the list of visited vertices,  $P$ .

$$P = \emptyset$$

# A Star Search with Revisit, Part II

## Algorithm

- While  $Q$  is not empty and goal is not deQueued, deQueue  $Q$ , put it on  $P$  and enQueue its successors to  $Q$ , and update the cost functions.

$$s = Q_{(0)} = \arg \min_{s \in Q} g(s) + h(s)$$

$$P = P + s$$

$$Q = Q + s'(s), \text{ update } g(s') = \min \{g(s'), g(s) + c(s, s')\}$$

# A Search Performance

## Discussion

- $A^*$  is complete.
- $A^*$  is optimal.

# Iterative Deepening A Star Search

## Discussion

- $A^*$  can use a lot of memory.
- Do path checking without expanding any vertex with  $g(s) + h(s) > 1$ .
- Do path checking without expanding any vertex with  $g(s) + h(s) > 2$ .
- ...
- Do path checking without expanding any vertex with  $g(s) + h(s) > d$ .

# Iterative Deepening A Star Search Performance

## Discussion

- IDA\* is complete.
- IDA\* is optimal.
- IDA\* is more costly than A\*.

# Beam Search

## Discussion

- Version 1: Keep a priority queue with fixed size  $k$ . Only keep the top  $k$  vertices and discard the rest.
- Version 2: Only keep the vertices that are at most  $\epsilon$  worse than the best vertex in the queue.  $\epsilon$  is called the beam width.

# Beam Search Performance

## Discussion

- Beam is incomplete.
- Beam is not optimal.