

# CS540 Introduction to Artificial Intelligence

## Lecture 16

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Based on lecture slides by Jerry Zhu, Yingyu Liang, and Charles Dyer

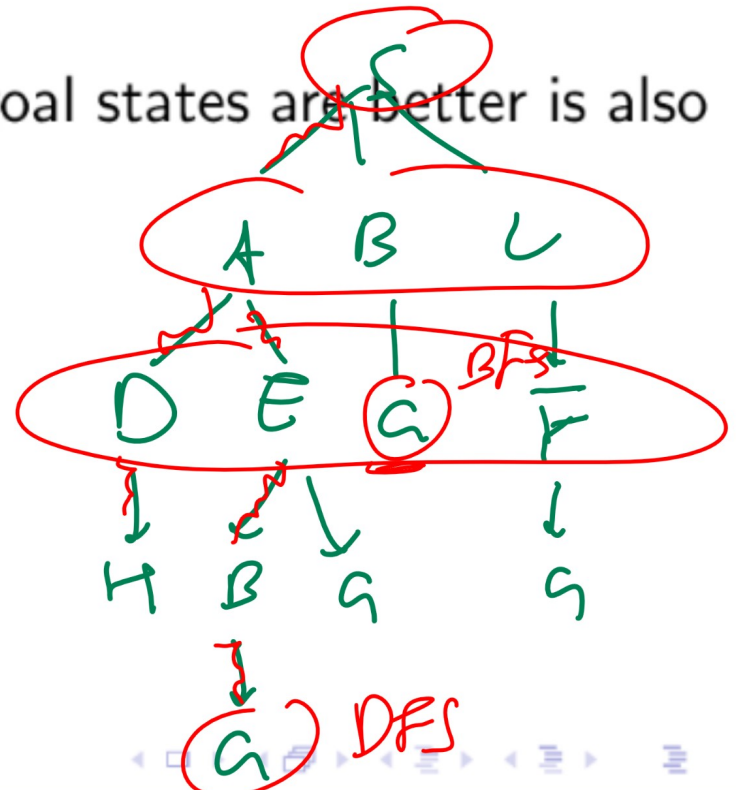
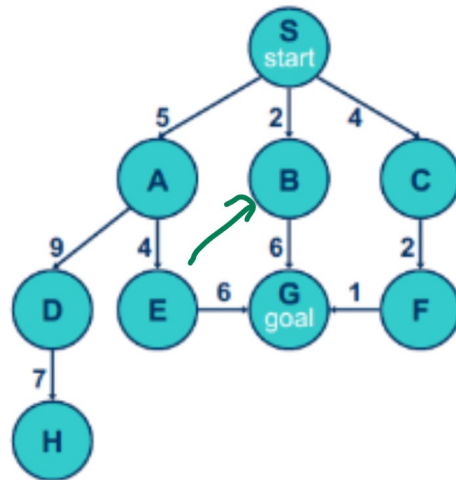
July 20, 2020

# ~~Uniformed~~ vs. Informed Search

Uninformed

Motivation

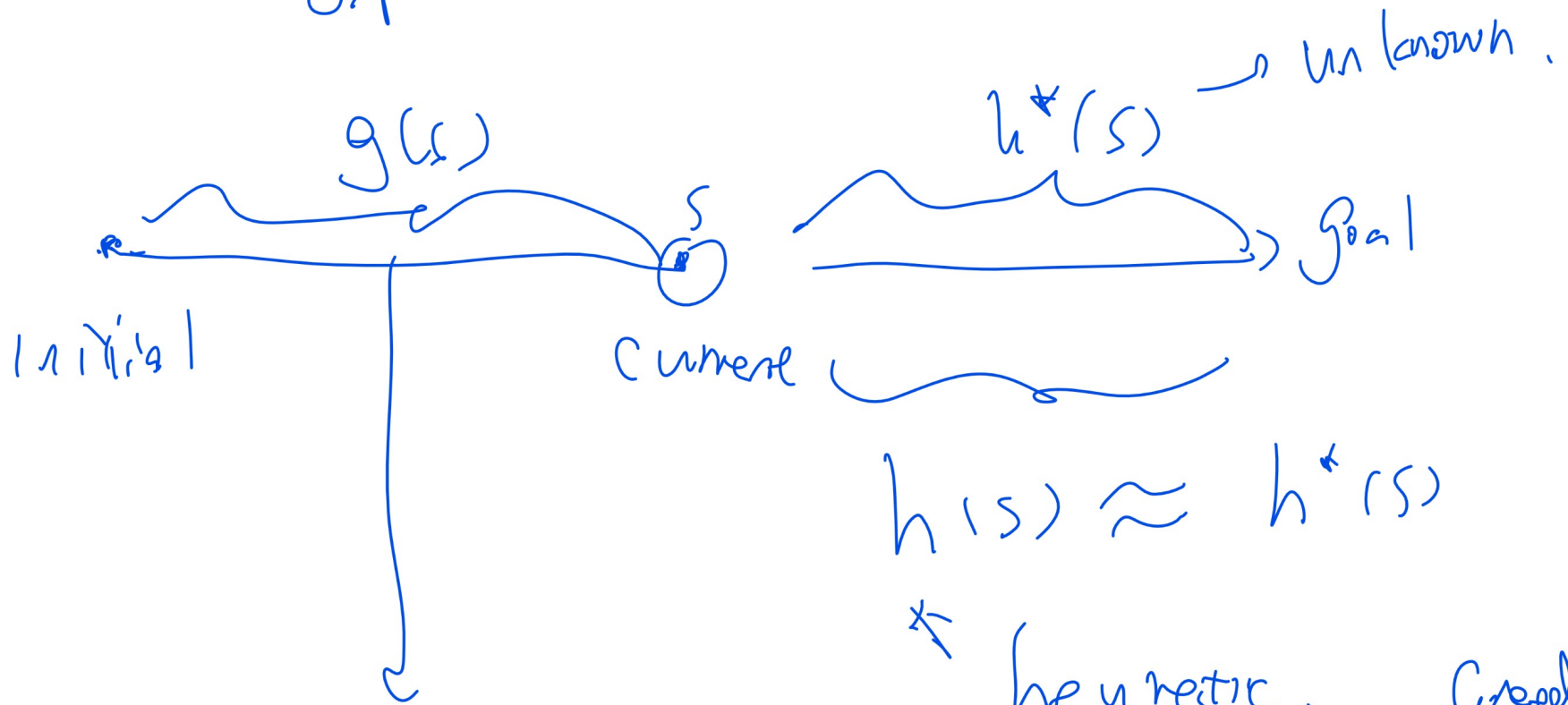
- Uninformed search means only the goal  $G$  and the successor functions  $s'$  are given.
- Informed search means which non-goal states are better is also known.



# Heuristic Diagram

Motivation

A search  $h$   $g+h$   $f$  use.  
 expand smallest  $g+h$   $f$  use.



expand smallest  $g$  first  $\rightarrow$  UCS

heuristic. Greedy  
 expand smallest  $h$  first  $\rightarrow$  BFS

# Uniform Cost Search

## Description

- Expand the vertices with the lowest current path cost  $g(s)$  first.
- It is BFS with a priority queue based on  $g(s)$ .
- It is equivalent to BFS if  $c = 1$  is constant on all edges.
- It is also called Dijkstra's Algorithm.

# UCS Example 1

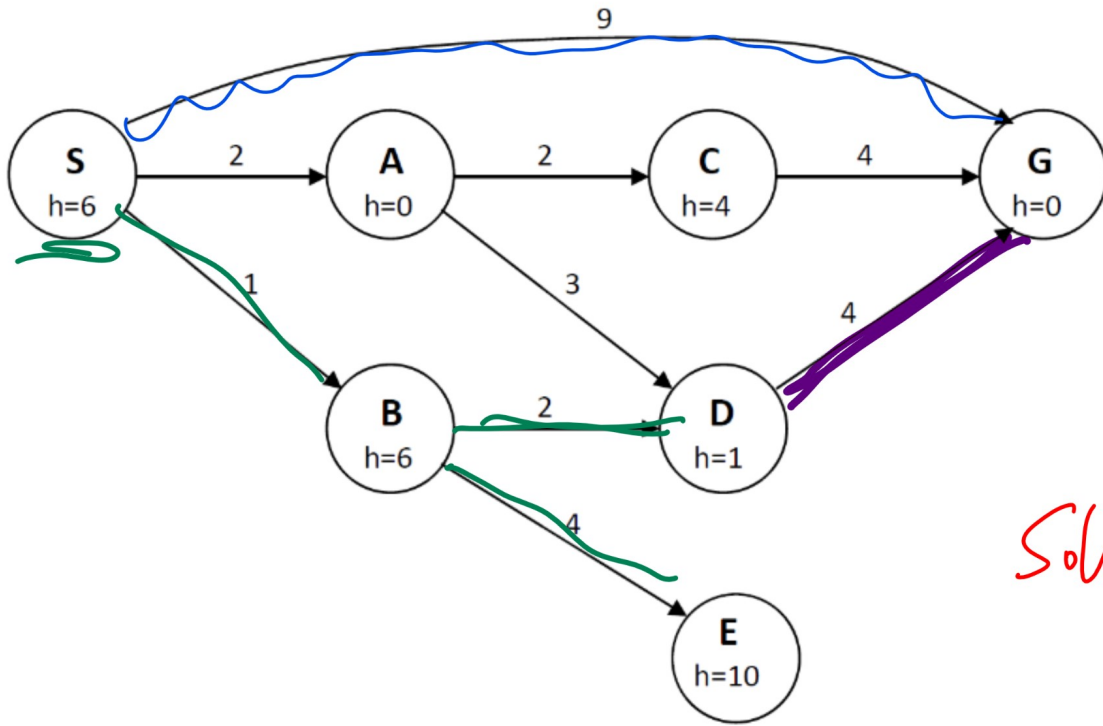
## Quiz

- Spring 2017 Midterm Q1
- Given the following adjacency matrix. Find UCS expansion path.

—	S	A	B	C	D	E	G
S	$h = 6$	2	1	—	—	—	9
A	—	$h = 0$	—	2	3	—	—
B	—	—	$h = 6$	—	2	4	—
C	—	—	—	$h = 4$	—	—	4
D	—	—	—	—	$h = 1$	—	4
E	—	—	—	—	—	$h = 10$	—
G	—	—	—	—	—	—	$h = 0$

# UCS Example 1 Diagram

Quiz



Expansion path:  
 S, B, A, D

C, D, E, G

Solution S, B, D, G

Priority Queue:  
 S, B, A  
 0, 1, 2

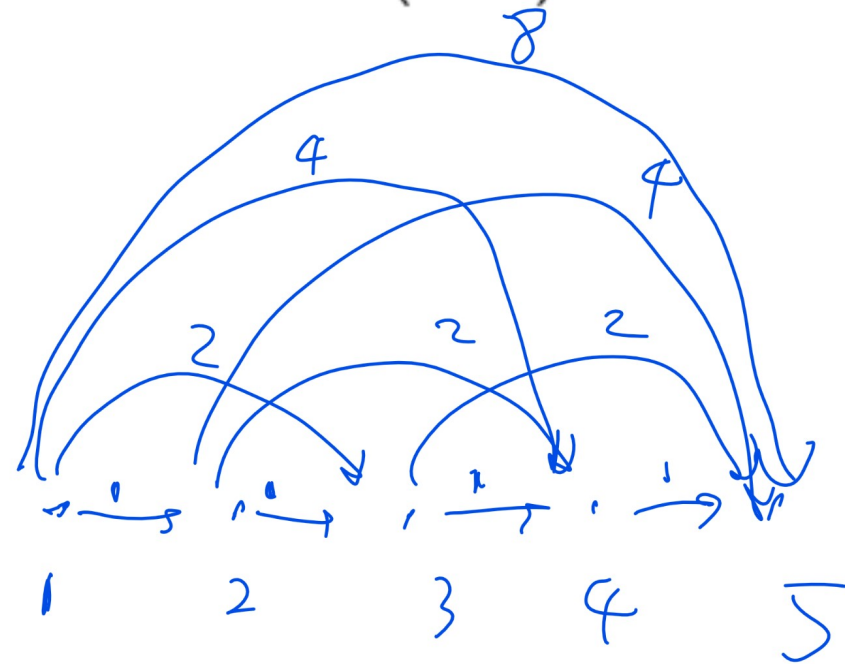
~~D~~, ~~B~~, ~~C~~, ~~D~~, ~~E~~, ~~G~~, G, G, G  
 1+2, 2+2, 2+3, 1+4, 3+4, 4+4, 4+4, 9  
 (A purple arrow points from the first 3+4 to the second 3+4)

# UCS Example 2

## Quiz

Q 11

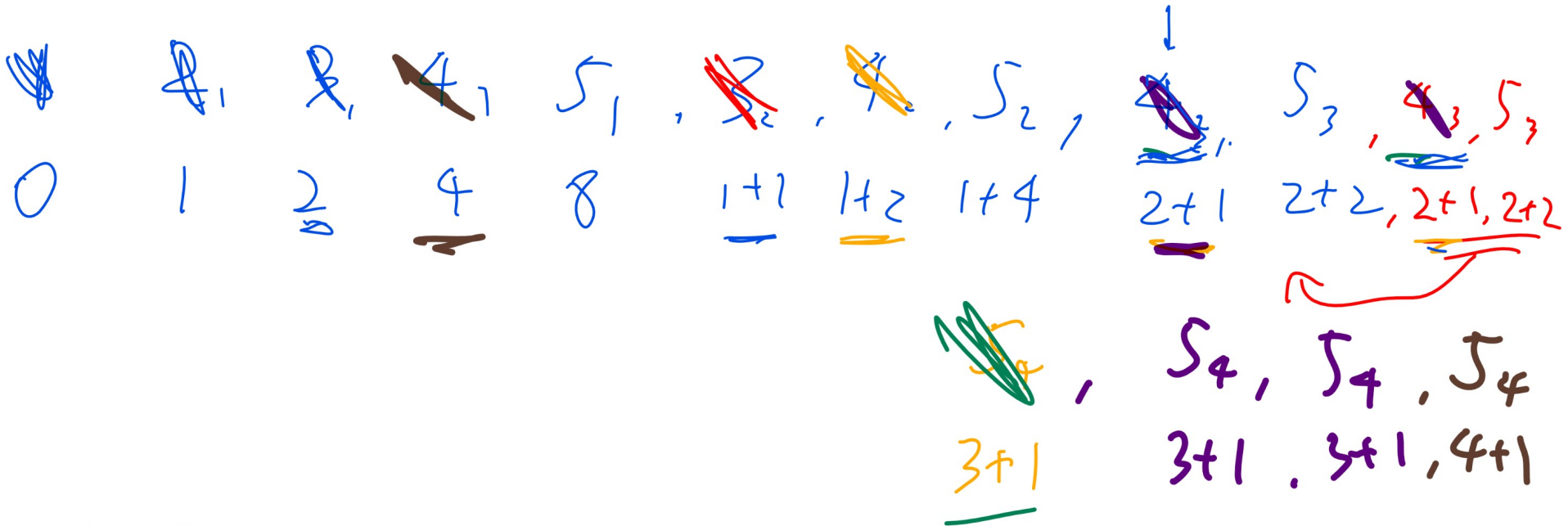
- Given that cost from state  $i$  to  $j$  is  $2^{j-i-1}$  for  $j > i$ . The initial state is 1 and goal state is 5. What is a vertex expansion sequence if Uniform Cost Search (UCS) is used?
- A: 1, 5
- B: 1, 2, 3, 4, 5
- C: 1, 2, 3, 4, 4, 5
- D: 1, 2, 3, 3, 4, 4, 5
- E: 1, 2, 3, 3, 4, 4, 4, 5



if  $c = \text{constant}$ , BFS = UCS.

# UCS Example 2 Diagram

Queue:  
selecting smallest to dequeue;  
Quiz



1, 2, 3, 4, 4, 4, 4, 5

post video later



# Uniform Cost Search

## Algorithm

- Input: a weighted digraph  $(V, E, c)$ , initial states  $I$  and goal states  $G$ .
- Output: a path from  $I$  to  $G$ .
- EnQueue initial states into a priority queue  $Q$ . Here,  $Q$  is ordered by  $g(s)$  for  $s \in Q$ .

$$Q = I$$

- While  $Q$  is not empty and goal is not deQueued, deQueue  $Q$  and enQueue its successors.

$$s = Q_{(0)} = \arg \min_{s \in Q} g(s)$$

$$Q = Q + s'(s)$$

# Uniform Cost Search Performance

## Discussion

- UCS is complete.
- UCS is optimal with any  $c$ .

# Best First Greedy Search

## Description

- Expand the vertices with the lowest heuristic cost  $h(s)$  first.
- Use a priority queue based on  $h(s)$ .

# Greedy Example 1

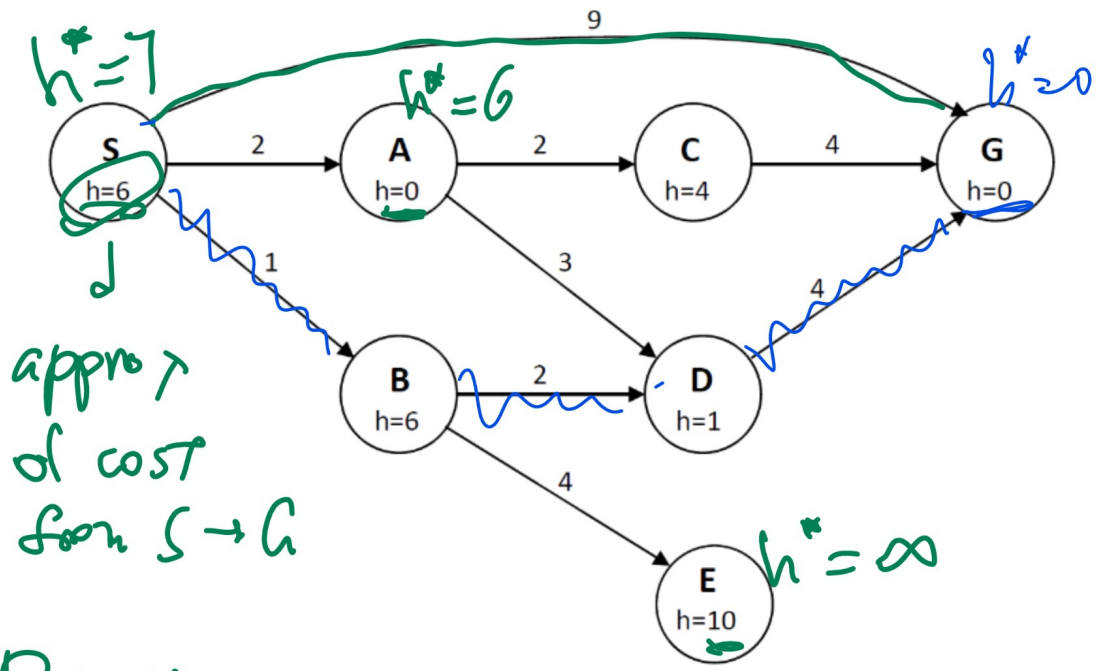
## Quiz

- Given the following adjacency matrix. Find Greedy Search expansion path.

—	S	A	B	C	D	E	G
S	$h = 6$	2	1	—	—	—	9
A	—	$h = 0$	—	2	3	—	—
B	—	—	$h = 6$	—	2	4	—
C	—	—	—	$h = 4$	—	—	4
D	—	—	—	—	$h = 1$	—	4
E	—	—	—	—	—	$h = 10$	—
G	—	—	—	—	—	—	$h = 0$

# Greedy Example 1 Diagram

Quiz



approx of cost for S → G

expansion:

S, A, G

Solution

S → G

Priority Queue:

~~A~~<sub>S</sub>    ~~G~~<sub>S</sub> ~~D~~<sub>A</sub> ~~C~~<sub>A</sub> ~~S~~<sub>S</sub>    B<sub>S</sub>

h            0            0 1 4 6            6

BFS  
 Greedy  
 is wrong  
 in general.

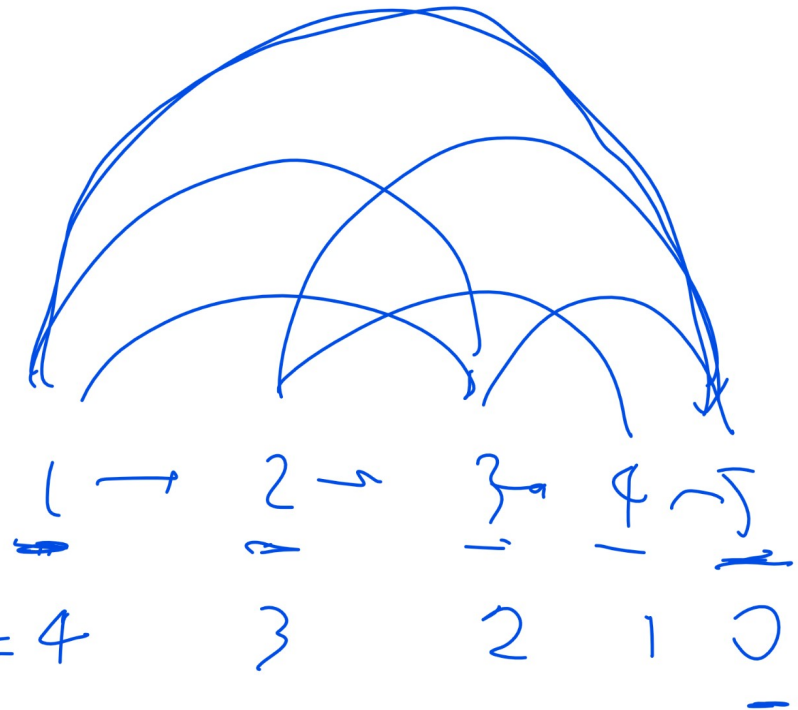
# Greedy Example 2

## Quiz

Q12

Given that cost from state  $i$  to  $j$  is  $2^{j-i-1}$  for  $j > i$ . The heuristic is  $h(i) = 5 - i$ . The initial state is 1 and goal state is 5. What is a vertex expansion sequence if Best First Greedy Search is used?

- A: 1, 5
- B: 1, 2, 3, 4, 5
- C: 1, 2, 3, 4, 4, 5
- D: 1, 2, 3, 3, 4, 4, 5
- E: 1, 2, 3, 3, 4, 4, 4, 5





# Best First Greedy Search

## Algorithm

- Input: a weighted digraph  $(V, E, c)$ , initial states  $I$  and goal states  $G$ , and the heuristic function  $h(s), s \in V$ .
- Output: a path from  $I$  to  $G$ .
- EnQueue initial states into a priority queue  $Q$ . Here,  $Q$  is ordered by  $h(s)$  for  $s \in Q$ .

$$Q = I$$

- While  $Q$  is not empty and goal is not deQueued, deQueue  $Q$  and enQueue its successors.

$$s = Q_{(0)} = \arg \min_{s \in Q} h(s)$$

$$Q = Q + s'(s)$$



# Best First Greedy Search Performance

## Discussion

- Greedy is incomplete.
- Greedy is not optimal.

# A Search

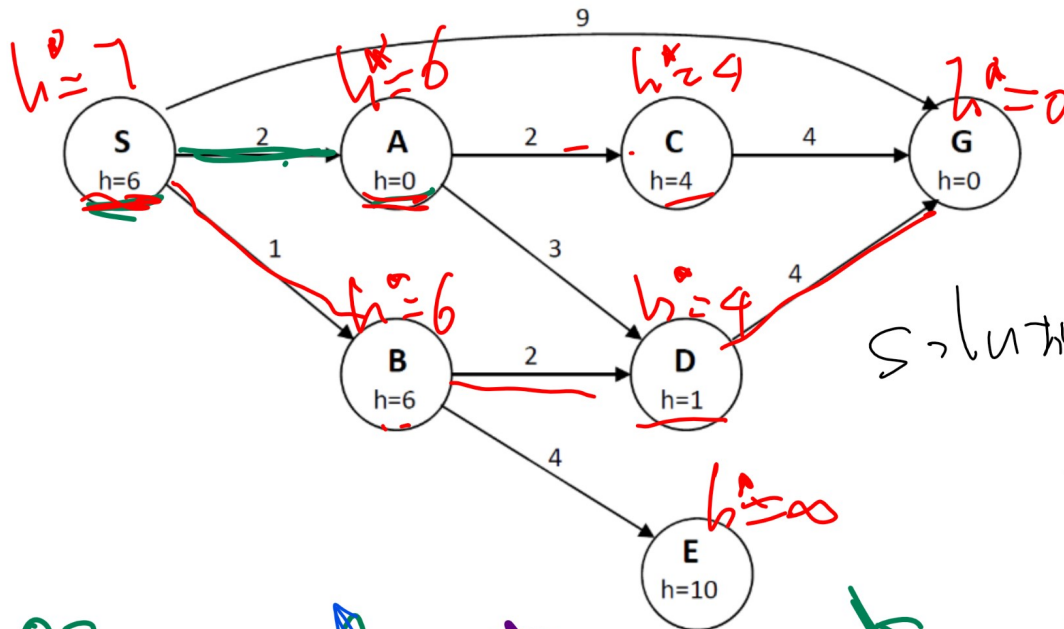
## Description

- Expand the vertices with the lowest total cost  $g(s) + h(s)$  first.
- Use a priority queue based on  $g(s) + h(s)$ .
- A stands for Always be optimistic?



# A Search Example 1 Diagram

Quiz



expansion

S, A, D, B, D, G

Solution:

$S \rightarrow B \rightarrow D \rightarrow G$

path	<del>A</del>	<del>B</del>	<del>D</del>	<del>S</del>	<del>B</del>	<del>G</del>	C <sub>A</sub>	C <sub>S</sub>	C <sub>D</sub>	F <sub>B</sub>
g	2	3	5	0	1	7	4	9	9	5
h	0	1	1	6	6	0	4	0	4	10
g+h	2	4	6	6	7	7	8	9	9	15

# A Search

## Algorithm

- Input: a weighted digraph  $(V, E, c)$ , initial states  $I$  and goal states  $G$ , and the heuristic function  $h(s), s \in V$ .
- Output: a path from  $I$  to  $G$ .
- EnQueue initial states into a priority queue  $Q$ . Here,  $Q$  is ordered by  $g(s) + h(s)$  for  $s \in Q$ .

$$Q = I$$

- While  $Q$  is not empty and goal is not deQueued, deQueue  $Q$  and enQueue its successors.

$$s = Q_{(0)} = \arg \min_{s \in Q} g(s) + h(s)$$

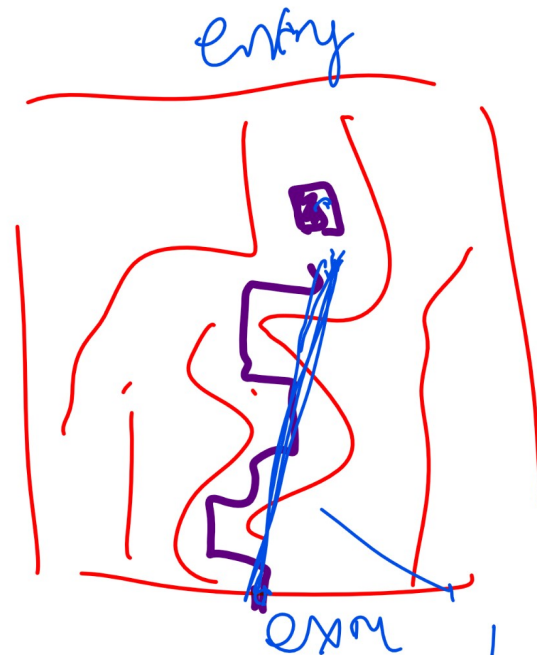
$$Q = Q + s'(s)$$

# A Search Performance

## Discussion

- A is complete.
- A is not optimal.

*h* can be wrong



Euclidean  
 $h = \text{dist}$   
 $\leq h^*$

# A Star Search

## Description

- A\* search is A search with an admissible heuristic.

$$0 \leq h \leq h^*$$

# Admissible Heuristic

## Definition

- A heuristic is admissible if it never over estimates the true cost.

$$0 \leq h(s) \leq h^*(s)$$



# Dominated Heuristic

## Definition

- One heuristic,  $h_1$ , is dominated by another,  $h_2$ , if:

$$h_1(s) \leq h_2(s) \leq h^*(s), \forall s \in S$$

- If  $h_2$  dominates  $h_1$ , then  $h_2$  is better than  $h_1$  since  $A^*$  using  $h_1$  expands at least as many states (or more) than  $A^*$  using  $h_2$ .
- If  $h_2$  dominated  $h_1$ ,  $A^*$  with  $h_2$  is better informed than  $A^*$  with  $h_1$ .

# Admissible Heuristic 8 Puzzle Example

## Quiz

$$0 \leq h \leq h^*$$

• Which ones (select multiple) of the following are admissible heuristic function for the 8 Puzzle?

- A:  $h(s)$  = number of tiles in the wrong position. ✓
- B:  $h(s) = 0$ . ✓  $0 \leq h \leq h^* \rightarrow$  UCS
- C:  $h(s) = 1$ .  $h(goal) = 1 >$   $h^*(goal) = 0$
- D:  $h(s)$  = sum of Manhattan distance between each tile and its goal location.  $\leq h^*$  ✓
- E:  $h(s)$  = sum of Euclidean distance between each tile and its goal location.  $\leq h_D \leq h^*$  ✓

# Admissible Heuristic General Example 1

## Quiz

$$0 \leq h \leq h^*$$

• Which ones (select multiple) of the following are admissible heuristic function?

• A:  $h(s) = h^*(s)$ .

• B:  $h(s) = \max\{2, h^*(s)\}$ .

✓ • C:  $h(s) = \min\{2, h^*(s)\}$ .

• D:  $h(s) = h^*(s) - 2$ .

✗ • E:  $h(s) = \sqrt{h^*(s)}$ .

$$h^*(\text{goal}) = 0$$

$$h(\text{goal}) = \max\{2, 0\} = 2 > 0$$

$$\leq h^*(s)$$

$$h^*(\text{goal}) = 0$$

$$h(\text{goal}) = 0 - 2 < 0$$

$$h^*(s) = 0.25, \quad \underline{h(s)} = \sqrt{0.25} = 0.5 > 0.25 = h^*$$

# A Star Search Example 2

## Quiz

Q13 (hse)

- Given that cost from state  $i$  to  $j$  is  $2^{j-i-1}$  for  $j > i$ . How many of the following heuristic functions are admissible? For  $i = 1, 2, 3, 4, 5$ :

$0 \leq h \leq w^*$

$h(i) = 5 - i$

②  $h(i) = \sqrt{5 - i}$

③  $h(i) = \log_2(6 - i)$

④  $h(i) = 1 - \mathbb{1}_{i=5}$

$h(i) = 0$



$h_2$	$h_4$	$i$	$w^*(i)$	$h_3$
2	1	1	4	$\log_2 5$
$\sqrt{3}$	2	2	3	2
$\sqrt{2}$	1	3	2	$\log_2 3$
1	1	4	1	1
0	0	5	0	0

$\min(1, w^*)$

$i=5$   
 $i \neq 5$

- A: 1, B: 2, C: 3, D: 4, E: 5

# A Star Search Example 2

## Quiz

- Given that cost from state  $i$  to  $j$  is  $2^{j-i-1}$  for  $j > i$ . Which one of the following heuristic functions is not dominated (among the admissible ones)? For  $i = 1, 2, 3, 4, 5$ :
  - A:  $h(i) = 5 - i$
  - B:  $h(i) = \sqrt{5 - i}$
  - C:  $h(i) = \log_2(6 - i)$
  - D:  $h(i) = 1 - \mathbb{1}_{i=5}$
  - E:  $h(i) = 0$

# A Star Search with Revisit, Part I

## Algorithm

- Input: a weighted digraph  $(V, E, c)$ , initial states  $I$  and goal states  $G$ , and the heuristic function  $h(s), s \in V$ .
- Output: a path with minimum cost from  $I$  to  $G$ .
- EnQueue initial states into a priority queue  $Q$ . Here,  $Q$  is ordered by  $g(s) + h(s)$  for  $s \in Q$ .

$$Q = I$$

$$g(I) = 0$$

$$g(s) = \infty, \text{ for } s \notin I$$

- Initialize the list of visited vertices,  $P$ .

$$P = \emptyset$$

# A Star Search with Revisit, Part II

## Algorithm

- While  $Q$  is not empty and goal is not deQueued, deQueue  $Q$ , put it on  $P$  and enQueue its successors to  $Q$ , and update the cost functions.

$$s = Q_{(0)} = \arg \min_{s \in Q} g(s) + h(s)$$

$$P = P + s$$

$$Q = Q + s'(s), \text{ update } g(s') = \min \{g(s'), g(s) + c(s, s')\}$$

# A Search Performance

## Discussion

A with  $h \leq h^*$

- A\* is complete.
- A\* is optimal.



# Iterative Deepening A Star Search

## Discussion

- A\* can use a lot of memory.
- Do path checking without expanding any vertex with  $g(s) + h(s) > 1$ .
- Do path checking without expanding any vertex with  $g(s) + h(s) > 2$ .
- ...
- Do path checking without expanding any vertex with  $g(s) + h(s) > d$ .

IDA\*

# Iterative Deepening A Star Search Performance

## Discussion

- IDA\* is complete.
- IDA\* is optimal.
- IDA\* is more costly than  $A^*$ .

# Beam Search

## Discussion

- Version 1: Keep a priority queue with fixed size  $k$ . Only keep the top  $k$  vertices and discard the rest.
- Version 2: Only keep the vertices that are at most  $\epsilon$  worse than the best vertex in the queue.  $\epsilon$  is called the beam width.

$$\underbrace{g+h}_{\downarrow} \leq \epsilon$$

# Beam Search Performance

## Discussion

- Beam is incomplete.
- Beam is not optimal.