CS540 Introduction to Artificial Intelligence Lecture 19

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Course Evaluation

Admin

- Midterm grades individual adjustment (please email or private Piazza me before the final exam):
- Incorrect ID: Wisc email ID, not 10-digit campus ID, not CS email ID.
- Typos that cause the expressions to be not evaluatable.
- Operation Personal issues: everyone is allowed to use P6 grade (out of 8, but plus 2 for really good projects) to replace one part of the midterm (and final) (out of 10).
- M9 is ready, M10 not ready, M is ready (note the difference of maze format from last year, if you submitted before yesterday, you need to resubmit).

Remind Me to Start Recording Admin

- The messages you send in chat will be recorded: you can change your Zoom name now before I start recording.
- We will play the coordination game to post exam questions again: no communication in chat after I start recording.

Coordination Game

Admin

 You are not allowed to discuss anything about this question in the public chat. There will be around 10 new questions on the final exam. I will post n of them before the exam (probably next Tuesday):



- A: n = 0.
- B: n = 1 if more than 50 percent of you choose B.
- C: n = 2 if more than 75 percent of you choose C.
- D: n = 3 if more than 98 percent of you choose D.
- E: n = 0.
- I will repeat this question a second time. If you fail to coordinate both times, I will not post any of the new questions.

Coordination Game Repeat

Admin



- You are not allowed to discuss anything about this question in the public chat. There will be around 5 new questions on the final exam. I will post n of them before the exam (probably next Tuesday):
- A: n = 0.
- B: n = 1 if more than 50 percent of you choose B.
- C: n = 2 if more than 75 percent of you choose C.
- D: n = 3 if more than \Re percent of you choose D.
 - E: n = 0.

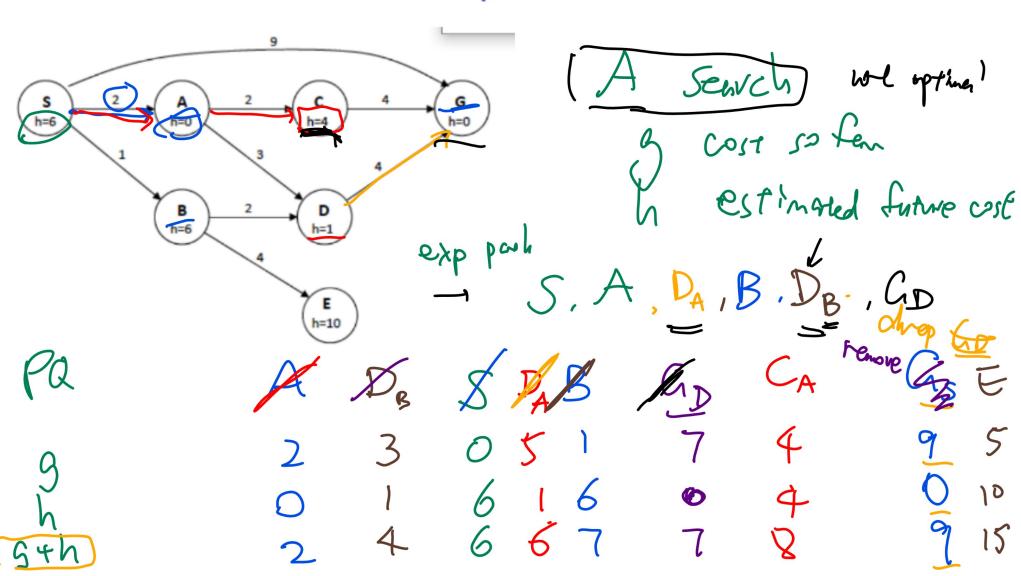
A Search Example 1

 Given the following adjacency matrix. Find A Search expansion path.

	S	А	В	С	D	E	G
S	h = 6	2	1	_	_	_	9
Α	_	h = 0	_	2	3	_	_
В	·—	1—1	h = 6	_	2	4	_
С	_	_	_	h = 4	_	_	4
D	-	_	_	_	h = 1	_	4
E	_	_	_	_	-	h = 10	_
G	_	_	_	_	_	_	h = 0

A Search Example 1 Diagram

Quiz



A Search Example 2

Quiz

Find A search expansion path

5		7	4	
_	S	A	В	G
S	h=3	1	999	_
Α	_	h = 1000	1	_
В	_	_	h = 1	1
G	_	_	_	h = 0





- B: S, B, G
- C: S, B, A, G
- D: S, B, A, B, G



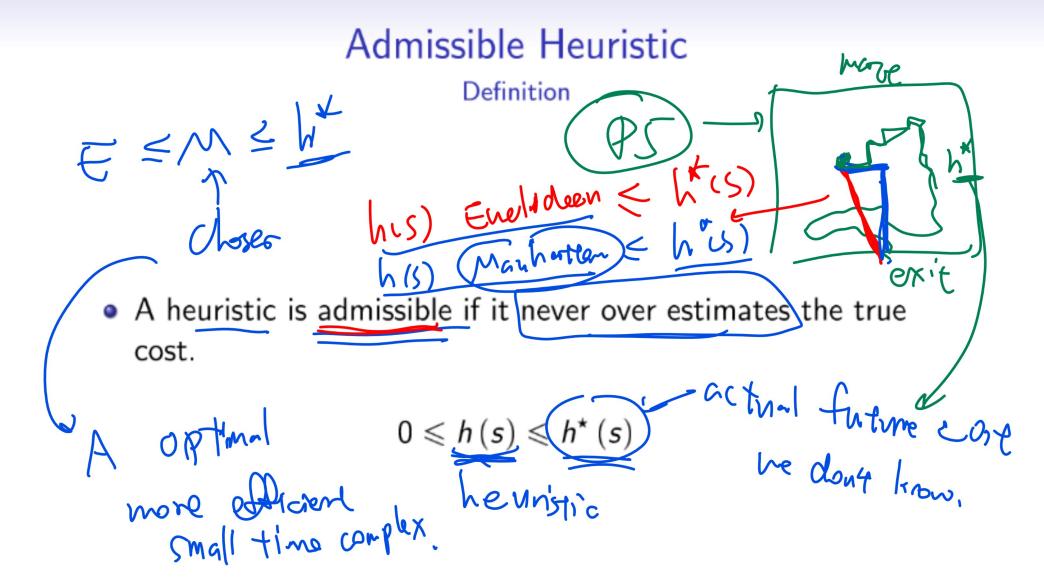




A Search Example 2 Diagram Quiz

A Search Performance

- A is complete.
- A is not optimal.



Dominated Heuristic

Definition

• One heuristic, h_1 , is dominated by another, h_2 , if:

$$h_1(s) \leqslant h_2(s) \leqslant h^*(s), \forall s \in S$$

- If h_2 dominates h_1 , then h_2 is better than h_1 since A^* using h_1 expands at least as many states (or more) than A^* using h_2 .
- If h₂ dominated h₁, A* with h₂ is better informed than A* with h₁.

Non-Optimal Heuristic

Definition

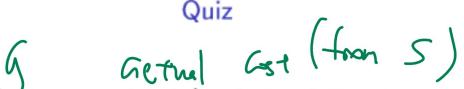
A search + admissible heuristic



- If optimality is not required and a satisfying solution is acceptable, then the heuristic should be as close as possible, either under or over, to the actual cost.
- This results in fewer states being expanded compared to using poor but admissible heuristics.

hus) = 0 admissibb

Admissible Heuristic General Example 1



- Which ones (select multiple) of the following are admissible heuristic function? a Cthal and Chron
- A: $h(s) = h^*(s) \cdot 2$.
- B: $h(s) = \sqrt{h^*(s)}$.
- C: $h(s) = h^*(s) + 1$.
- D: $h(s) = \min\{1, h^*(s)\}.$

$$E/h(s) = h^*(s) \cdot \frac{1}{2} h'(s) \text{ at } h$$

$$F: h(s) = h^*(s)^2.$$

$$6. h(s) = \max\{1, h^*(s)\}.$$

$$h^{\star}(s)-1$$

$$h'(s) = h^*(s)^2$$

$$h'(s) > 1$$

$$h'(s) > 1$$

$$h'(s) > 1$$

$$h'(s) = max\{1, h^*(s)\}. \iff af G \rightarrow h'(s) = 0, h'(s) = 1$$

Admissible Heuristic General Example 1 Again

 Which ones (select multiple) of the following are admissible heuristic function?

$$\bullet B. h(s) = \sqrt{h^*(s)}.$$

$$\&: h(s) = h^{\star}(s) + 1. > h^{\star}$$

• D;
$$h(s) = \min \{1, h^*(s)\}. \le \bigwedge^* (s)$$

$$E: h(s) = h^*(s) \cdot \frac{1}{2}.$$

•
$$f: h(s) = h^*(s)^2$$

•
$$G: h(s) = \max\{1, h^*(s)\}.$$

• H:
$$h(s) = h^*(s) - 1$$
.

$$\frac{h_{1}(\lambda,y)}{5} \leq x$$

A Search Performance

- A[⋆] is complete. [⋆]
- A[⋆] is optimal.

Iterative Deepening A Star Search



- A* can use a lot of memory.
- Do path checking without expanding any vertex with g(s) + h(s) > 1.
- Do path checking without expanding any vertex with g(s) + h(s) > 2.
- ...
- Do path checking without expanding any vertex with g(s) + h(s) > d.

Iterative Deepening A Star Search Performance

- IDA* is complete.
- IDA* is optimal.
- IDA* is more costly than A*.

Beam Search

- Version 1: Keep a priority queue with fixed size k. Only keep the top k vertices and discard the rest.
- Version 2: Only keep the vertices that are at most ε worse than the best vertex in the queue. ε is called the beam width.

Beam Search Performance

- Beam is incomplete.
- Beam is not optimal

Traveling Salesperson Example

Motivation



- Local search is about searching through a state space by iteratively improving the cost to find an optimal or near-optimal state.
- The successor states are called the neighbors (sometimes move set).
- The assumption is that similar (nearby) solutions have similar costs.

Hill Climbing (Valley Finding) Description

- Start at a random state.
- Move to the best neighbor state (one of the successors).
- Stop when all neighbors are worse than the current state.
- The idea is similar to gradient descent.

Boolean Satisfiability Example 1

- Assume all variables A, B, C, D, E are set to True. How many of the following clauses are satisfied?
- $A \vee \neg B \vee C$
- $\bullet \neg A \lor C \lor D$
- B ∨ D ∨ ¬E
- \bullet $\neg C \lor \neg D \lor \neg E$
- \bullet $\neg A \lor \neg C \lor E$

Boolean Satisfiability Example 2

- Assume all variables A, B, C, D, E are set to True. Which one
 of the variables should be changed to False to maximize the
 number of clauses satisfied?
- $A \vee \neg B \vee C$
- \bullet $\neg A \lor C \lor D$
- $B \vee D \vee \neg E$
- \bullet $\neg C \lor \neg D \lor \neg E$
- $\bullet \neg A \lor \neg C \lor E$

Boolean Satisfiability Example 3

- Assume all variables A, B, C, D, E are set to True. Which one
 of the variables should be changed to False to maximize the
 number of clauses satisfied?
- ¬A ∨ ¬B ∨ ¬E
- ¬A ∨ ¬B ∨ ¬D
- \bullet $\neg A \lor \neg C \lor \neg D$
- \bullet $\neg B \lor \neg C \lor \neg D$
- \bullet $\neg C \lor \neg D \lor \neg E$

Hill Climbing

Algorithm

- Input: state space S and cost function f.
- Output: $s^* \in S$ that minimizes f(s).
- Start at a random state s₀.
- At iteration t, find the neighbor that minimizes f.

$$s_{t+1} = \arg\min_{s \in s'(s_t)} f(s)$$

Stop when none of the neighbors have a lower cost.

stop if
$$f(s_{t+1}) \leq f(s_t)$$

Random Restarts

Discussion

 A simple modification is picking random initial states multiple times and finding the best among the local minima.

First Choice Hill Climbing

- If there are too many neighbors, randomly generate neighbors until a better neighbor is found.
- This method is called first choice hill climbing.

Simulated Annealing

Description

- Each time, a random neighbor is generated.
- If the neighbor has a lower cost, move to the neighbor.
- If the neighbor has a higher cost, move to the neighbor with a small probability.
- Stop until bored.
- It is a version of Metropolis-Hastings Algorithm.