

# CS540 Introduction to Artificial Intelligence

## Lecture 19

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Based on lecture slides by Jerry Zhu, Yingyu Liang, and Charles Dyer

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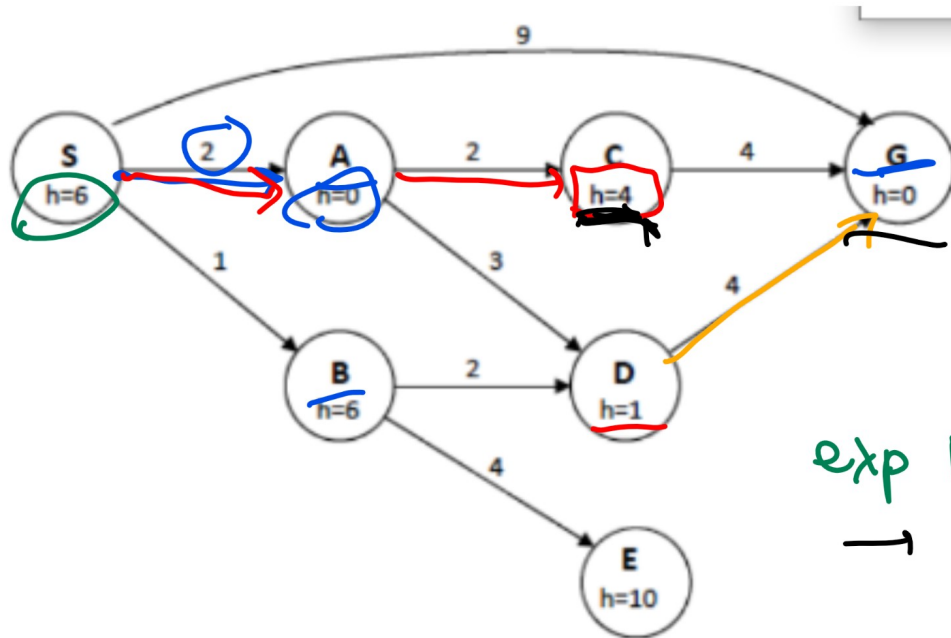






# A Search Example 1 Diagram

Quiz



A search we optimize!  
 cost so far  
 estimated future cost

exp path

→ S, A, D<sub>A</sub>, B, D<sub>B</sub>, G, D

PQ  
 g  
 h  
g+h

<del>A</del>	<del>D<sub>B</sub></del>	<del>S</del>	<del>D<sub>A</sub></del>	<del>B</del>	<del>G</del>	<del>D</del>	C <sub>A</sub>	<del>C<sub>B</sub></del>	<del>E</del>
2	3	0	5	1	7	4	4	9	5
0	1	6	1	6	0	4	4	0	10
2	4	6	6	7	7	8	9	9	15

# A Search Example 2

## Quiz

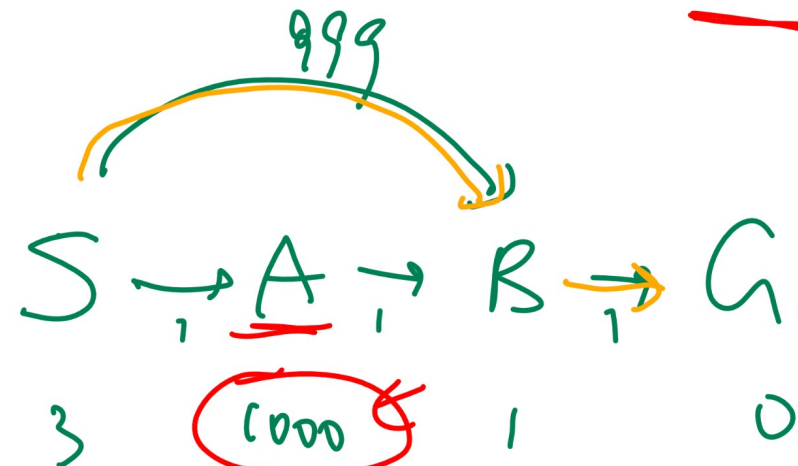
- Find A search expansion path

-	<u>S</u> <sup>1</sup>	A <sup>2</sup>	B <sup>3</sup>	G <sup>4</sup>
S	$h = 3$	1	999	-
A	-	$h = 1000$	1	-
B	-	-	$h = 1$	1
G	-	-	-	$h = 0$

Q4

- ~~A: S, A, B, G~~
- B: S, B, G
- C: S, B, A, G
- D: S, B, A, B, G

~~A~~ ~~B~~ ~~G~~  
 3 1001 1000-  
 incorrect



overestimating





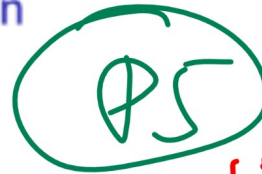
# A Search Performance

## Discussion

- A is complete.
- A is not optimal.

# Admissible Heuristic

Definition



$$C \leq M \leq \underline{h^*}$$

↑  
closer

$h(s)$  Euler's  $\leq h^*(s)$   
 $h(s)$  Manhattan  $\leq h^*(s)$

- A heuristic is admissible if it never over estimates the true cost.

A optimal  
more efficient  
small time complex.

$$0 \leq \underline{h(s)} \leq \underline{h^*(s)}$$

heuristic

actual future cost  
we don't know.

# Dominated Heuristic

## Definition

- One heuristic,  $h_1$ , is dominated by another,  $h_2$ , if:

$$h_1(s) \leq h_2(s) \leq h^*(s), \forall s \in S$$

- If  $h_2$  dominates  $h_1$ , then  $h_2$  is better than  $h_1$  since  $A^*$  using  $h_1$  expands at least as many states (or more) than  $A^*$  using  $h_2$ .
- If  $h_2$  dominated  $h_1$ ,  $A^*$  with  $h_2$  is better informed than  $A^*$  with  $h_1$ .

# Non-Optimal Heuristic

## Definition

A search + admissible heuristic



- If optimality is not required and a satisfying solution is acceptable, then the heuristic should be as close as possible, either under or over, to the actual cost.
- This results in fewer states being expanded compared to using poor but admissible heuristics.

$h(s) = 0$     admissible  
 $\Rightarrow$  ~~YES~~

# Admissible Heuristic General Example 1

## Quiz

Q5

Which ones (select multiple) of the following are admissible heuristic function?

- A:  $h(s) = h^*(s) \cdot 2$ .
- B:  $h(s) = \sqrt{h^*(s)}$ .
- C:  $h(s) = h^*(s) + 1$ .
- D:  $h(s) = \min\{1, h^*(s)\}$ .

• E:  $h(s) = h^*(s) \cdot \frac{1}{2}$

• F:  $h(s) = h^*(s)^2$

• G:  $h(s) = \max\{1, h^*(s)\}$

• H:  $h(s) = h^*(s) - 1$

Actual cost (from S)

Actual cost (from G)

$$0 \leq h(s) \leq h^*(s)$$

$\forall s$

$h^*(s)$  at G  $h(G) = 0$  ✓

$h^*(s) > 1$   $h = h^* < h^*$

at G  $\rightarrow h^*(s) = 0, h(s) = 1$

at G  $\rightarrow h^*(s) = 0$   
 $\rightarrow -1$  at G

## Admissible Heuristic General Example 1 Again

## Quiz

- Q5
- Which ones (select multiple) of the following are admissible heuristic function?

~~A~~:  $h(s) = h^*(s) \cdot 2.$   $> h^*$

~~B~~:  $h(s) = \sqrt{h^*(s)}.$

~~C~~:  $h(s) = h^*(s) + 1.$   $> h^*$

D:  $h(s) = \min\{1, h^*(s)\}.$   $\leq h^*(s)$

E:  $h(s) = h^*(s) \cdot \frac{1}{2}.$

~~F~~:  $h(s) = h^*(s)^2.$

~~G~~:  $h(s) = \max\{1, h^*(s)\}.$

H:  $h(s) = h^*(s) - 1.$

$$\min(x, y) \leq x$$

$$\min(x, y) \leq y$$

$$\sqrt{0.25} = 0.5$$

$$0.25 < 0.5$$

check	0
check	0-1
check	1+

$$0 \leq h(s) \leq h^*(s) \forall s$$

# A Search Performance

## Discussion

- $A^*$  is complete. ✓
- $A^*$  is optimal. ✓





# Iterative Deepening A Star Search Performance

## Discussion

- IDA\* is complete. ✓
- IDA\* is optimal. ✓
- IDA\* is more costly than  $A^*$ .

# Beam Search

## Discussion

- Version 1: Keep a priority queue with fixed size  $k$ . Only keep the top  $k$  vertices and discard the rest.
- Version 2: Only keep the vertices that are at most  $\epsilon$  worse than the best vertex in the queue.  $\epsilon$  is called the beam width.

# Beam Search Performance

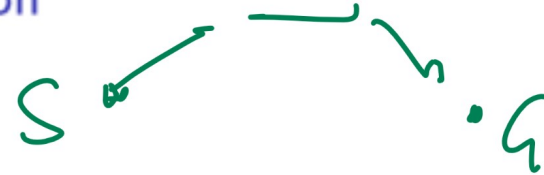
## Discussion

- Beam is incomplete. ✓
- Beam is not optimal. ✓



## Local Search

Motivation



- Local search is about searching through a state space by iteratively improving the cost to find an optimal or near-optimal state.
- The successor states are called the neighbors (sometimes move set).
- The assumption is that similar (nearby) solutions have similar costs.

→ [gradient descent]

$w$  →  $w$  →  $w^*$   
 $C(w)$                        $C$                        $C$

# Hill Climbing (Valley Finding)

## Description

- Start at a random state.
- Move to the best neighbor state (one of the successors).
- Stop when all neighbors are worse than the current state.
- The idea is similar to gradient descent.

# Boolean Satisfiability Example 1

## Quiz

- Assume all variables  $A$ ,  $B$ ,  $C$ ,  $D$ ,  $E$  are set to True. How many of the following clauses are satisfied?
- $A \vee \neg B \vee C$
- $\neg A \vee C \vee D$
- $B \vee D \vee \neg E$
- $\neg C \vee \neg D \vee \neg E$
- $\neg A \vee \neg C \vee E$



## Boolean Satisfiability Example 2

### Quiz

- Assume all variables  $A$ ,  $B$ ,  $C$ ,  $D$ ,  $E$  are set to True. Which one of the variables should be changed to False to maximize the number of clauses satisfied?
- $A \vee \neg B \vee C$
- $\neg A \vee C \vee D$
- $B \vee D \vee \neg E$
- $\neg C \vee \neg D \vee \neg E$
- $\neg A \vee \neg C \vee E$

# Boolean Satisfiability Example 3

## Quiz

- Assume all variables  $A$ ,  $B$ ,  $C$ ,  $D$ ,  $E$  are set to True. Which one of the variables should be changed to False to maximize the number of clauses satisfied?
- $\neg A \vee \neg B \vee \neg E$
- $\neg A \vee \neg B \vee \neg D$
- $\neg A \vee \neg C \vee \neg D$
- $\neg B \vee \neg C \vee \neg D$
- $\neg C \vee \neg D \vee \neg E$

# Hill Climbing

## Algorithm

- Input: state space  $S$  and cost function  $f$ .
- Output:  $s^* \in S$  that minimizes  $f(s)$ .
- Start at a random state  $s_0$ .
- At iteration  $t$ , find the neighbor that minimizes  $f$ .

$$s_{t+1} = \arg \min_{s \in S'(s_t)} f(s)$$

- Stop when none of the neighbors have a lower cost.

$$\text{stop if } f(s_{t+1}) \leq f(s_t)$$

# Random Restarts

## Discussion

- A simple modification is picking random initial states multiple times and finding the best among the local minima.

# First Choice Hill Climbing

## Discussion

- If there are too many neighbors, randomly generate neighbors until a better neighbor is found.
- This method is called first choice hill climbing.

