◆□▶ ◆□▶ ◆三▶ ◆三▶ 三三 のへぐ

CS540 Introduction to Artificial Intelligence Lecture 17

Young Wu

Based on lecture slides by Jerry Zhu, Yingyu Liang, and Charles Dyer

August 3, 2022

DFS 00000000000

▲□▶ ▲□▶ ▲□▶ ▲□▶ ▲□ ● ● ●

Learning vs Search

- In reinforcement learning, the reward and state transition need to be learned by taking actions.
- In search problems, the reward and state transitions are given.
- The problem is to find a sequence of actions that lead to the goal with minimum cost.

DFS 00000000000

▲□▶ ▲□▶ ▲□▶ ▲□▶ ▲□ ● ● ●

Search Problem

- State space S is the set of all valid configurations.
- Initial states *I* and goal states *G* are subsets of *S*.
- Successor function s' (s) given the current state s is the set of states reachable in one step from s.
- There is a cost (or negative reward) associated with moving from s to s' (s).
- The search problem is the problem of finding a solution path from a state in *I* to a state in *G*, usually with minimum total cost.

BFS

DFS 0000000000

▲□▶ ▲□▶ ▲ 三▶ ▲ 三▶ 三 のへぐ

State Space Motivation

- The states need to represent all necessary information about the game.
- The actions are discrete and deterministic and are determined by the successor function.
- Each possible action at state *s* is associated with a state in the set *s'*(*s*).

BFS

DFS 0000000000

◆□▶ ◆□▶ ◆三▶ ◆三▶ 三三 - のへで

Sizes of State Space

- Tic Tac Toe: 10^3
- Checkers: 10^{20}
- Chess: 10⁵⁰
- Go: 10¹⁷⁰

BFS

DFS 00000000000

▲ロ ▶ ▲周 ▶ ▲ 国 ▶ ▲ 国 ▶ ● の Q @

State Space Graph Definition

- A state space can be represented by a weighted directed graph (V, E, c).
- V is the set of vertices (also called nodes).
- *E* is the set of edges (also called arcs). Each edge is directed from one vertex to another vertex and represents an action.
- *c* is the cost (also called weights) associated with each edge. The costs are positive.

Search Problem on Graph Definition

- Search starts at an initial state and finishes if one of the goal states is reached.
- The solution is a path in the graph from an initial state to a goal state.
- The cost of a solution is the sum of edge costs on the solution path.
- The optimal solution is the solution with the lowest cost.

▲□▶ ▲□▶ ▲□▶ ▲□▶ ▲□ ● ● ●

Expansion Definition

- Vertices that are explored so far are stored in a tree called the state space search tree.
- Expanding a vertex means to generate all successor vertices and add them (and the associated edges) to the state space search tree.
- The leaves of the search tree are unexpanded and are called the frontier (sometimes called the fringe).
- The search strategies differ in the order in which the vertices are expanded.

BFS

DFS 0000000000

▲ロ ▶ ▲周 ▶ ▲ 国 ▶ ▲ 国 ▶ ● の Q @

Performance Definition

- A search strategy is complete if it finds at least one solution.
- A search strategy is optimal if it finds the optimal solution.
- For uninformed search, the costs are assumed to be 1 for all edges c = 1.

▲□▶ ▲□▶ ▲□▶ ▲□▶ ▲□ ● ● ●

Complexity Definition

- The time complexity of a search strategy is the worst case maximum number of vertices expanded.
- The space complexity of a search strategy is the worst case maximum number of states stored in the frontier at a single time.
- Notation: the goals are *d* edges away from the initial state. This means assuming a constant cost of 1, the optimal solution has cost *d*. The maximum depth of the graph is *D*.
- Notation: the branching factor is *b*, the maximum number of actions associated with a state.

$$b = \max_{s \in V} \left| s'\left(s\right) \right|$$

▲□▶ ▲□▶ ▲ 三▶ ▲ 三▶ 三三 - のへぐ

Breadth First Search

- Use Queue (FIFO) for the frontier.
- Remove from the front, add to the back.

▲□▶ ▲□▶ ▲□▶ ▲□▶ □ のQで

Breadth First Search

- Input: a weighted digraph (V, E, c), initial states I and goal states G.
- Output: a path from I to G.
- EnQueue initial states.

$$Q = I$$

• While Q is not empty and goal is not deQueued, deQueue Q and enQueue its successors.

$$s = Q_0$$

 $Q = Q + s' (s)$

Breadth First Search Performance

- BFS is complete.
- BFS is optimal with c = 1.

Breadth First Search Complexity

• Time complexity: the worst case occurs when the goal is the last vertex at depth *d*.

$$T = b + b^2 + \ldots + b^d$$

• Space complexity: the worst case is storing all vertices at depth *d* is in the frontier.

$$S = b^d$$

BiDirectional Search

- BFS from the initial states and goal states at the same time.
- The search stops when the two frontiers meet (have non-empty intersection) in the middle.
- The time and space complexity is the same as BFS with depth $\frac{d}{d}$
 - 2

BFS

DFS •000000000

▲□▶ ▲□▶ ▲ 三▶ ▲ 三▶ 三三 - のへぐ

Depth First Search

Description

- Use Stack (LIFO) for the frontier.
- Remove from the front, add to the front.

Depth First Search

- Input: a weighted digraph (V, E, c), initial states I and goal states G.
- Output: a path from I to G.
- Push initial states.

$$S = I$$

• While S is not empty and goal is not popped, pop S and push its successors.

$$s = S_0$$
$$S = s'(s) + S$$

BFS

DFS 0000000000

Depth First Search Performance

• DFS is incomplete if $D = \infty$.

• DFS is not optimal.

▲ロ ▶ ▲周 ▶ ▲ 国 ▶ ▲ 国 ▶ ● の Q @

Depth First Search Complexity

Time complexity: the worst case occurs when the goal is the root of the last subtree expanded in the whole graph.

$$T = b^{D-d+1}...+b^{D-1}+b^{D}$$

• Space complexity: the worst case is storing all vertices sharing the parents with vertices in the current path.

$$S = (b-1)D + 1$$

BFS 00000 DFS 0000000000

▲□▶ ▲□▶ ▲□▶ ▲□▶ ▲□ ● ● ●

Iterative Deepening Search

- DFS but stop if path length > 1
- repeat DFS but stop if path length > 2
- ...
- repeat DFS but stop if path length > d

▲□▶ ▲□▶ ▲□▶ ▲□▶ □ のQで

Iterative Deepening Search Algorithm

- Input: a weighted digraph (V, E, c), initial states I and goal states G.
- Output: a path from I to G.
- Perform DFS on the digraph restricted to vertices with depth $\leqslant 1$ from the initial state.
- Perform DFS on the digraph restricted to vertices with depth $\leqslant 2$ from the initial state.
- Repeat until the goal is deQueued.

Iterative Deepening Search Performance

- IDS is complete.
- IDS is optimal with c = 1.

▲□▶ ▲□▶ ▲□▶ ▲□▶ □ のQで

Iterative Deepening Search Complexity Discussion

• Time complexity: the worst case occurs when the goal is the last vertex at depth *d*.

$$T = db + (d - 1)b^{2} + \dots + 3b^{d-2} + 2b^{d-1} + 1b^{d}$$

• Space complexity: it has the same space complexity as DFS. $S = (b-1) \ d$

BFS

DFS 00000000000

▲□▶ ▲□▶ ▲□▶ ▲□▶ ▲□ ● ● ●

Non-Tree Search

- If the state space is not a tree, search strategies need to remember the states that are already expanded.
- A vertex should be removed from the frontier if it is already expanded.

Uniformed vs Informed Search

- Uninformed search means only the goal G and the successor function s' are given.
- Informed search means which non-goal states are better is also known.
- Usually, iterative deepening is used for uninformed search.