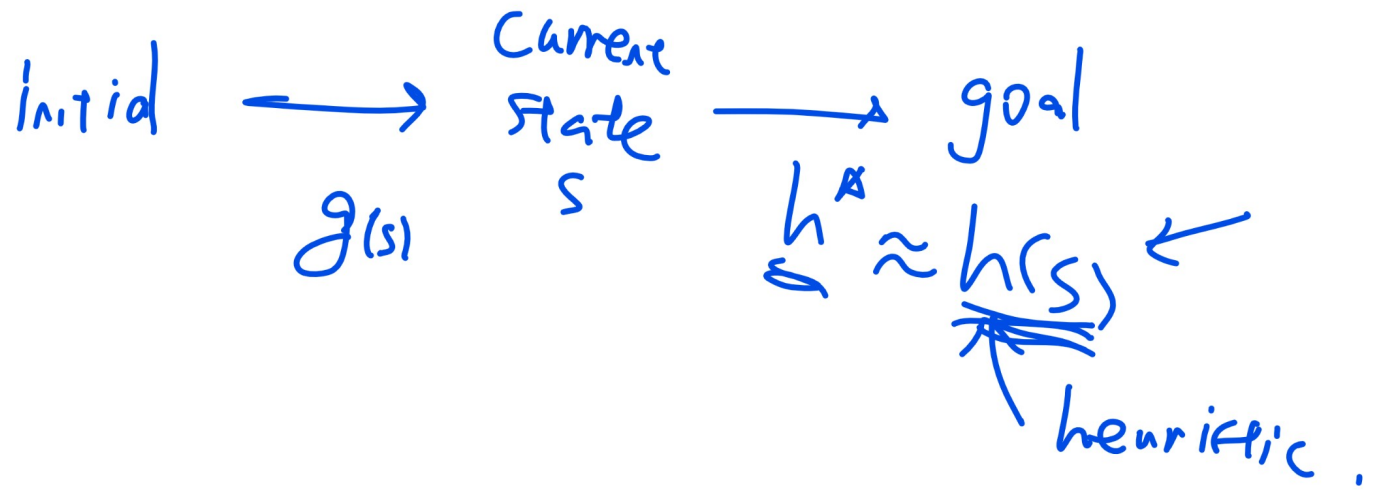


Heuristic Diagram

Motivation



UCS Example 1

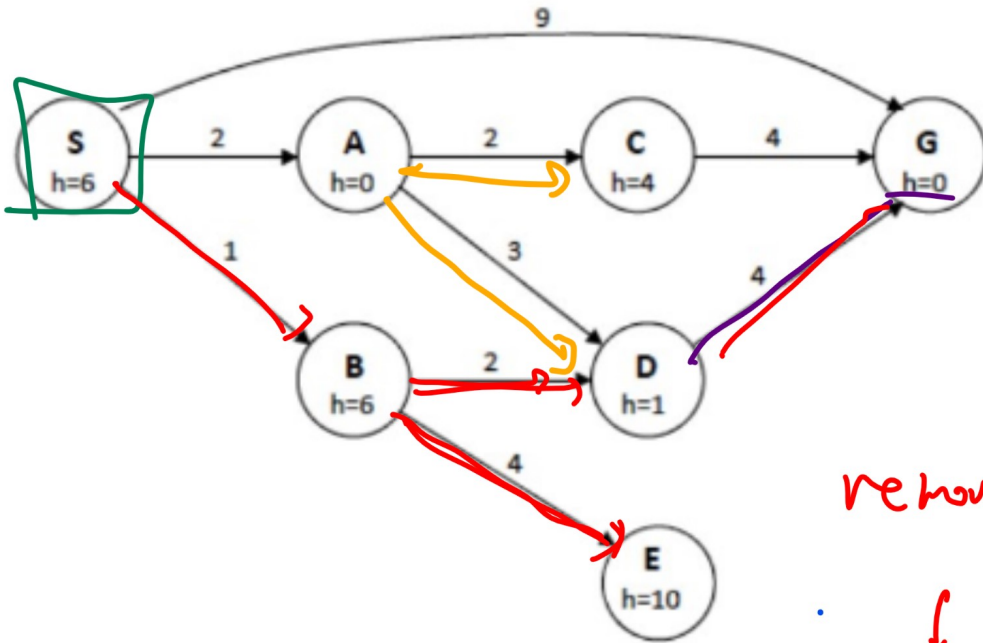
Quiz

- Given the following adjacency matrix. Find UCS expansion path.

—	<i>S</i>	<i>A</i>	<i>B</i>	<i>C</i>	<i>D</i>	<i>E</i>	<i>G</i>
<i>S</i>	$h = 6$	2	1	—	—	—	9
<i>A</i>	—	$h = 0$	—	2	3	—	—
<i>B</i>	—	—	$h = 6$	—	2	4	—
<i>C</i>	—	—	—	$h = 4$	—	—	4
<i>D</i>	—	—	—	—	$h = 1$	—	4
<i>E</i>	—	—	—	—	—	$h = 10$	—
<i>G</i>	—	—	—	—	—	—	$h = 0$

UCS Example 1 Diagram

Quiz



remove highest priority (lowest cost)

list (queue)

~~S~~
g=0

~~A~~
g=2

~~B~~
g=1

~~G~~
g=0

~~D~~
g=3

~~E~~
g=5

~~C~~
g=4

~~G~~
g=0

7

←

$g_{S \rightarrow D} + g_{D \rightarrow G}$

$g_{S \rightarrow A} + g_{A \rightarrow C}$
 $g_{S \rightarrow C} + g_{C \rightarrow A}$

UCS Example 2

Quiz

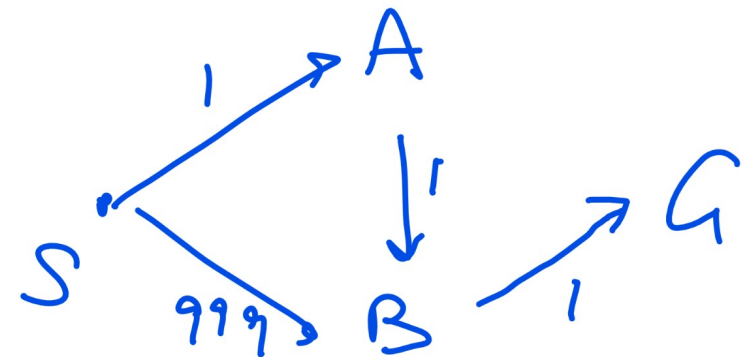
- Find UCS expansion path → list of removed states

—	S	A	B	G
S	$h = 3$	1	999	—
A	—	$h = 1000$	1	—
B	—	—	$h = 1$	1
G	—	—	—	$h = 0$

Q2

- A: S, A, B, G
- B: S, B, G
- C: S, B, A, G
- D: S, B, A, B, G

E: I don't understand.



Greedy Example 1

Quiz

- Given the following adjacency matrix. Find Greedy Search expansion path.

—	<i>S</i>	<i>A</i>	<i>B</i>	<i>C</i>	<i>D</i>	<i>E</i>	<i>G</i>
<i>S</i>	$h = 6$	2	1	—	—	—	9
<i>A</i>	—	$h = 0$	—	2	3	—	—
<i>B</i>	—	—	$h = 6$	—	2	4	—
<i>C</i>	—	—	—	$h = 4$	—	—	4
<i>D</i>	—	—	—	—	$h = 1$	—	4
<i>E</i>	—	—	—	—	—	$h = 10$	—
<i>G</i>	—	—	—	—	—	—	$h = 0$

Greedy Example 2

Quiz

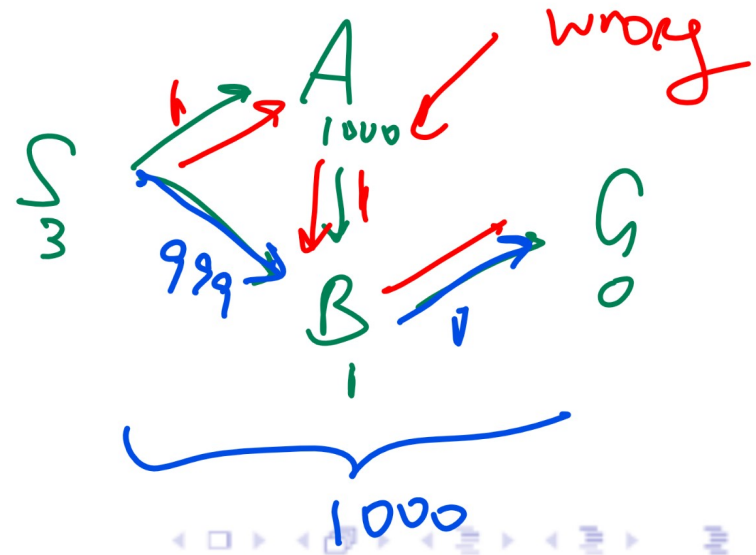
- Find Greedy expansion path

—	S	A	B	G
S	$h = 3$	1	999	—
A	—	$h = 1000$	1	—
B	—	—	$h = 1$	1
G	—	—	—	$h = 0$

Q}

- A: S, A, B, G
- B: S, B, G
- C: S, B, A, G
- D: S, B, A, B, G

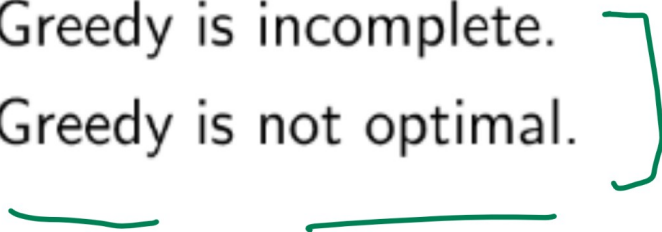
E:



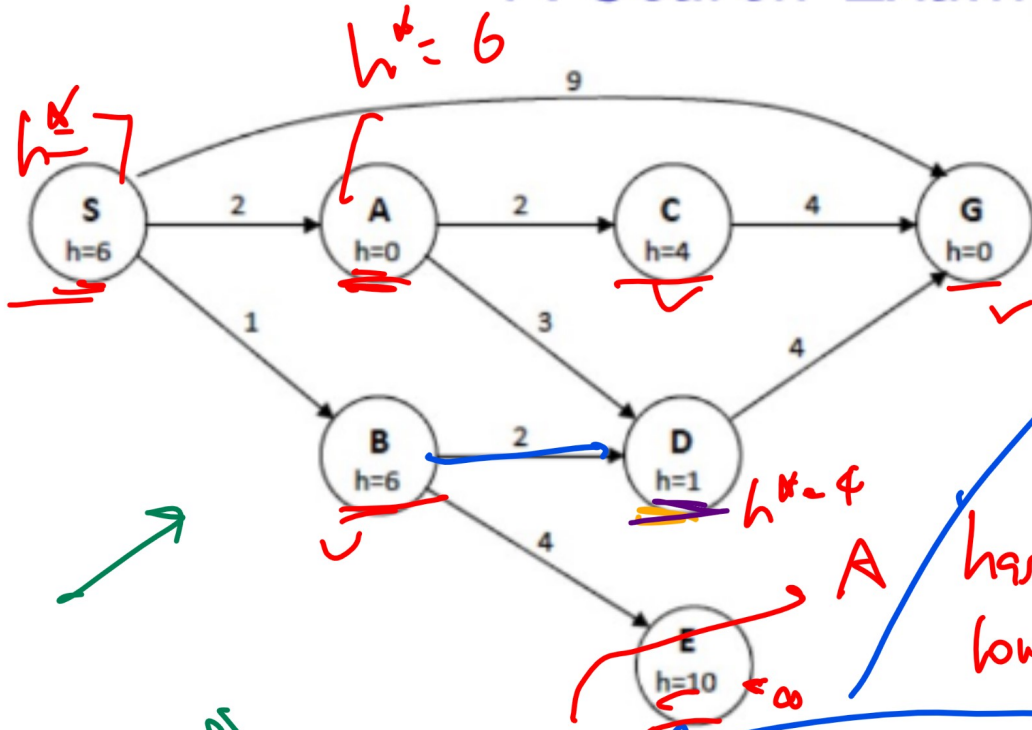
Best First Greedy Search Performance

Discussion

- Greedy is incomplete.
- Greedy is not optimal.



A Search Example 1 Diagram



replace not part of expansion.
Stop when G is removed.

A has higher priority
lowest $g+h$ (total estimated cost)

$S \rightarrow A \rightarrow G$

S	A	B
g	0	2
h	6	6
$g+h$	6	7

C	D	E	G
g	2	2	3
h	4	1	0
$g+h$	8	6	7

Opt soln $S \rightarrow B \rightarrow D \rightarrow G$

S, A, D, B, D, G

from list, $S \rightarrow D \rightarrow G$

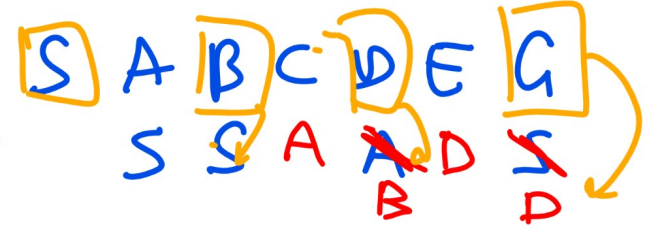
$$g_{S \rightarrow B} + g_{B \rightarrow D} = 3$$

A Search Example 2

1 + 2 Quiz

from diagram table

predecessor



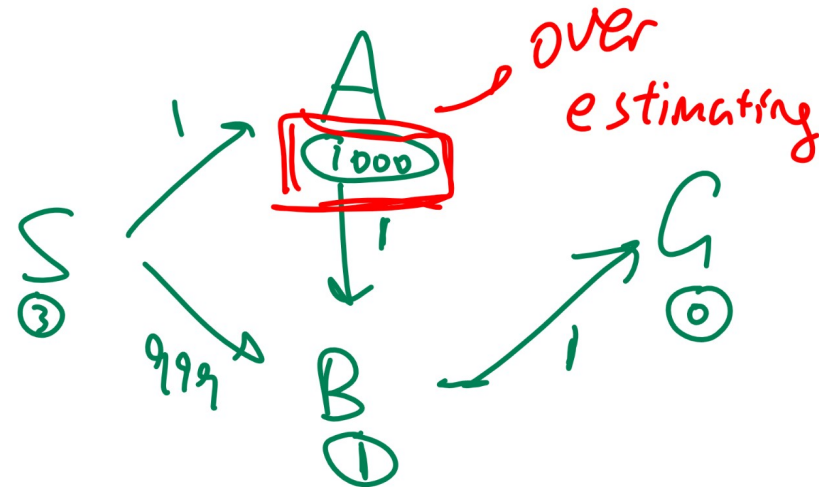
- Find A search expansion path

-	S	A	B	G
S	$h = 3$	1	999	-
A	-	$h = 1000$	1	-
B	-	-	$h = 1$	1
G	-	-	-	$h = 0$

7:40

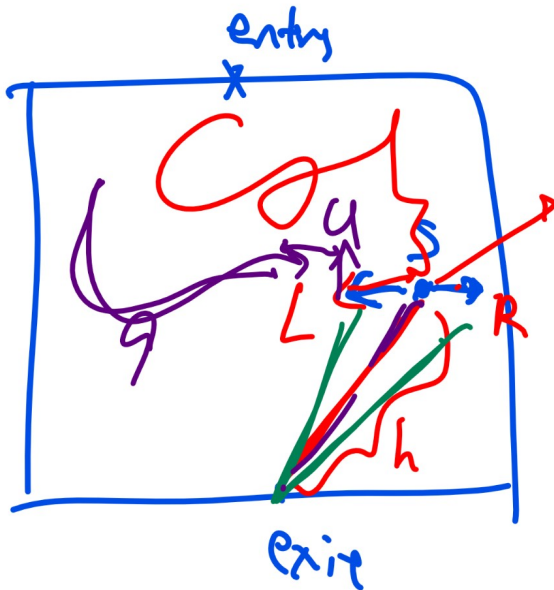
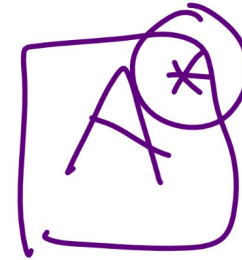
Q4

- A: S, A, B, G
- B: S, B, G**
- C: S, B, A, G
- D: S, B, A, B, G



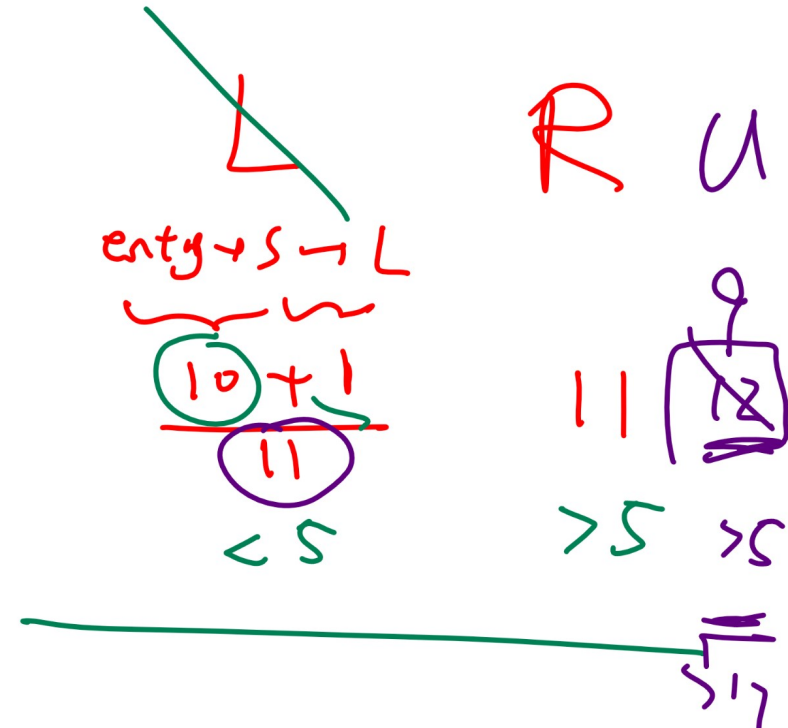
A Search Example 2 Diagram

		A	B	C	Quiz
g	3	1	999	1000	}
h	3	1000	1	0	
Sum	3	1001	1000	1000	



$g_{entry \rightarrow S} = 10$

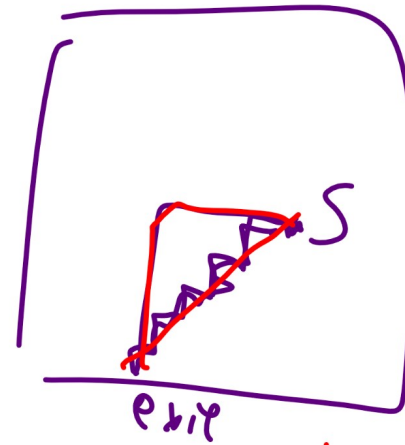
$h_{S \rightarrow exit} \approx 5$



A Search Performance

Discussion

- A is complete.
- A is not optimal.



always
underestimating
always positive

A Star Search

Description

- A^* search is A search with an admissible heuristic.

Admissible Heuristic

Definition

- A heuristic is admissible if it never over estimates the true cost.

$$0 \leq h(s) \leq h^*(s)$$

Admissible Heuristic 8 Puzzle Example

Definition

Admissible Heuristic 8 Puzzle Example

Quiz

• Which ones (select multiple) of the following are admissible heuristic function for the 8 Puzzle?

• A: $h(s)$ = number of tiles in the wrong position.

16

$\leq h^*(s)$

• B: $h(s) = 0$.

$\rightarrow ? UCS \leq h^*(s)$

$\rightarrow 0$

• C: ~~$h(s) = 1$~~

$s \neq goal$

• D: $h(s)$ = sum of Manhattan distance between each tile and its goal location.

$\leq h^*(s)$

• E: $h(s)$ = sum of Euclidean distance between each tile and its goal location.

$\leq h^*(s)$



$h(s) \leq h^*(s)$

$s \neq goal$

$h(goal) = 1$

$h^*(goal) = 0$

A Star Search Example 1 Diagram

Quiz

Admissible Heuristic General Example 1

Quiz

$$0 \leq h \leq h^*$$

Q5

- Which ones (select multiple) of the following are admissible heuristic function?

~~A~~: $h(s) = h^*(s) \cdot 2 \Rightarrow h^*$

~~B~~: $h(s) = \sqrt{h^*(s)} \in h^*$

~~C~~: $h(s) = h^*(s) + 1 \Rightarrow h^*$

D: $h(s) = \min\{1, h^*(s)\}$

E: $h(s) = h^*(s) \cdot \frac{1}{2}$

~~F~~: $h(s) = h^*(s)^2$

~~G~~: $h(s) = \max\{1, h^*(s)\}$

~~H~~: $h(s) = h^*(s) - 1$

$\sqrt{0.25} = 0.5 \leq 0.25$

$h^* \leq 1 \Rightarrow h \leq h^*$
 $h^* > 1 \Rightarrow h = h^*$

$h^* = 0 \Rightarrow h = 0$

$0 \leq h^* \cdot \frac{1}{2} \in h^*$

$h^*(s) = 2 \quad h(s) = 4$

$h^*(s) = 0 \quad h(s) = 1$

$h^*(s) = 0 \Rightarrow h(s) = -1$

$0 \leq h(s) \leq h^*(s)$

admissible

2
0.5

A Search Performance

Discussion

- A^* is complete.
- A^* is optimal.

Iterative Deepening A Star Search

Discussion



IDS

- A^* can use a lot of memory.
- Do path checking without expanding any vertex with $g(s) + h(s) > 1$.
- Do path checking without expanding any vertex with $g(s) + h(s) > 2$.
- ...
- Do path checking without expanding any vertex with $g(s) + h(s) > d$.

Iterative Deepening A Star Search Performance

Discussion

- IDA* is complete.
- IDA* is optimal.
- ~~IDA*~~ is more costly than A*.



Beam Search

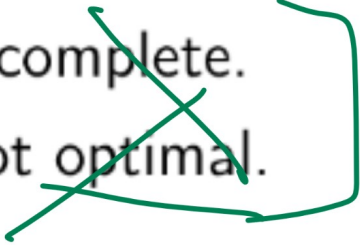
Discussion

- Version 1: Keep a priority queue with fixed size k . Only keep the top k vertices and discard the rest.
- Version 2: Only ^{store} keep the vertices that are at most ϵ worse than the best vertex in the queue. ϵ is called the beam width.



Beam Search Performance

Discussion

- Beam is incomplete.
 - Beam is not optimal.
- 
- A hand-drawn green bracket groups the two bullet points. A green slash is drawn over the entire list.

