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Adversarial Attacks on Multi-Agent Reinforcement Learning Preliminary Exam

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November 17, 2022



Thesis Statement

• There are vulnerabilities in multi-agent systems and attackers can influence the behavior of multi-agent reinforcement learners through data or environment poisoning.

Introduction

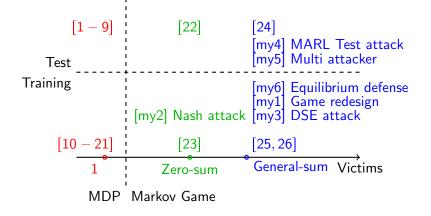
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List of Projects

- Completed:
- my1 Game Redesign: Training time, reward poisoning, online victims.
- my2 Nash Attack: Training time, reward poisoning, two offline victims.
- my3 DSE Attack: Training time, reward poisoning, multiple offline victims.
 - Future work:
- my4 MARL Test Attack: Test time, state or action manipulation, pre-trained victims.
- my5 Equilibrium Defense: Training or test time, attack-aware victims.
- my6 Multi Attacker: Training or test time, multiple attackers.



Adversarial RL Literature





[my1] Game Redesign

- my1 Joint work (\approx 15% contribution) with Yuzhe Ma (main author), and Jerry Zhu.
 - Victim setting:
 - The victims are no-regret online learners with O(T^α) regret, e.g. EXP3.P.
 - The victims participate in an *n*-player general-sum bandit game with original reward r^o (a) ∈ [-1,1]ⁿ for action profile a = (a₁, a₂, ..., a_n).

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Attacker Setting

- Attacker setting:
- The attacker wants the victims to take a target (deterministic) policy $\pi^{\dagger} = \left(\pi_{1}^{\dagger}, \pi_{2}^{\dagger}, ..., \pi_{n}^{\dagger}\right)$ as often as possible, i.e. maximize $\sum_{t=1}^{T} \mathbb{1}_{\left(a_{t}=\pi^{\dagger}\right)}$.
- The attacker can modify the rewards that the victims see from r^o (a) to r[†] (a).

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• The attacker wants sublinear design cost $\sum_{t=1}^{T} \left\| r^{o} \left(a_{t} \right) - r_{t}^{\dagger} \left(a_{t} \right) \right\|_{p}.$

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Interior Design Example

• Suppose $\pi^{\dagger} = (1, 1)$, the attacker can redesign the game r^{o} to r†, $r^{o} = \begin{bmatrix} (0,0) & (-1,\underline{1}) & (\underline{1},-1) \\ (\underline{1},-1) & (0,0) & (-1,\underline{1}) \\ (-1,\underline{1}) & (\underline{1},-1) & (0,0) \end{bmatrix},$ $r_1^{\dagger} = r_2^{\dagger} = \dots = \begin{bmatrix} (0, 0) & (0.1, -0.1) & (0.1, -0.1) \\ (-0.1, 0.1) & (0, 0) & (0, 0) \\ (-0.1, 0.1) & (0, 0) & (0, 0) \end{bmatrix}.$

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Interior Design Algorithm

- Given $r^{o}(a) \in [-1, 1]$, first consider the interior case when $r^{o}(\pi^{\dagger}) > -1$.
- Assumption: $r^{o}\left(\pi^{\dagger}\right) \geqslant -1 + \rho$, for some $\rho > 0$.

• Attack:
$$r_{i,t}^{\dagger}(a) = \begin{cases} r_i^o(\pi^{\dagger}) + \left(1 - \frac{d(a_t)}{n}\right)\rho & \text{if } a_{i,t} = \pi_i^{\dagger}, \\ r_i^o(\pi^{\dagger}) - \frac{d(a_t)}{n}\rho & \text{if } a_{i,t} \neq \pi_i^{\dagger}, \end{cases}$$

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where
$$d(a_t) = \sum_{i=1} \mathbb{1}_{\left\{a_{i,t}=\pi_i^{\dagger}\right\}}.$$



Interior Design Result

Theorem

Using the interior design, π^{\dagger} is used $T - O(nT^{\alpha})$ times while incurring design cost $O(n^{1+1/p}T^{\alpha})$.

• For example, EXP3.*P* with L_1 cost can achieve π^{\dagger} being used $T - O\left(n\sqrt{T}\right)$ times with cost $O\left(n^2\sqrt{T}\right)$.

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Interior Design Proof Sketch

• Under this attack, we have,

$$r_{i,t}^{\dagger}(a) = \begin{cases} r_i^o\left(\pi^{\dagger}\right) + \left(1 - \frac{d\left(a_t\right)}{n}\right)\rho & \text{ if } a_{i,t} = \pi_i^{\dagger}\\ r_i^o\left(\pi^{\dagger}\right) - \frac{d\left(a_t\right)}{n}\rho & \text{ if } a_{i,t} \neq \pi_i^{\dagger} \end{cases}$$

• π^{\dagger} is strictly dominant: $r_{i,t}^{\dagger}\left(\pi_{i,t}^{\dagger}, \mathbf{a}_{-i,t}\right) = r_{i,t}^{\dagger}\left(a_{i,t}, \mathbf{a}_{-i,t}\right) + \left(1 - \frac{1}{n}\right)\rho, \forall a_{i,t} \neq \pi_{i,t}^{\dagger}.$

2 π^{\dagger} rewards remain unchanged: $r_{i,t}^{\dagger}(\pi^{\dagger}) = r_{i}^{o}(\pi^{\dagger})$.

 No-regret learners will use the optimal action profile π[†] in all but O(T^α) rounds while incurring O(T^α) design cost.

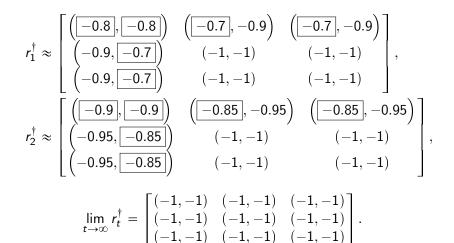
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Boundary Design Example

- When $r^{o}(\pi^{\dagger}) = -1$, it is impossible to decrease other entries below -1: another design is needed.
- Suppose again $\pi^{\dagger} = (1,1)$, then,

$$\begin{split} r^o &= \begin{bmatrix} (-1,-1) & \left(-1,\boxed{1}\right) & \left(\boxed{1},-1\right) \\ \left(\boxed{1},-1\right) & (-1,-1) & \left(-1,\boxed{1}\right) \\ \left(-1,\boxed{1}\right) & \left(\boxed{1},-1\right) & (-1,-1) \end{bmatrix}, \\ r_1^{\dagger} &\approx \begin{bmatrix} \left(\boxed{-0.8},\boxed{-0.8}\right) & \left(\boxed{-0.7},-0.9\right) & \left(\boxed{-0.7},-0.9\right) \\ \left(-0.9,\boxed{-0.7}\right) & (-1,-1) & (-1,-1) \\ \left(-0.9,\boxed{-0.7}\right) & (-1,-1) & (-1,-1) \end{bmatrix}, \end{split}$$

Boundary Design Example Limit



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Boundary Design Algorithm

- Assumption: $r^{o}(\pi^{\dagger}) = -1$.
- Attack: $r_{i,t}^{\dagger}(a) = w_t r_{i, \text{ interior }}^{\dagger}(a) + (1 w_t) r^o(\pi^{\dagger})$, where $w_t = t^{\alpha + \varepsilon 1}$, for some $\varepsilon \in (0, 1 \alpha]$.

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Boundary Design Result

Theorem

Using the boundary deisng with $\varepsilon = \frac{1-\alpha}{2}$, π^{\dagger} is used $T - O\left(nT^{(1+\alpha)/2}\right)$ times while incurring design cost $O\left(n^{1/p}\left(1+n\right)T^{(1+\alpha)/2}\right)$.

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Boundary Design Proof Sketch

Game Redesign

cost.

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[my2] Nash Attack

- In [my1], the attacker modifies the victims' rewards during online learning.
- In [my2], and [my3], the attacker modifies the rewards in an offline data set.

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Victim Setting

my2 Joint work (≈ 75% contribution) with Jeremy McMahan, Jerry Zhu, Qiaomin Xie. (Thanks: Yudong Chen)

- Victim setting:
- The victims are uncertainty-aware offline learners that use additive bonus terms β when estimating the Q function, i.e. Q = R̂ − β + ℝ_{ρ̂} [V'].
- **2** The victims learn a two-player zero-sum Markov game from a training set $\left\{ \left(\left(s_t^{(k)}, a_t^{(k)}, r_t^{(k)} \right)_{t=1}^T \right) \right\}_{k=1}^K$, with $r_t^{(k)} \in [0, 1]$.

Attacker Setting

- Attacker setting:
- The attacker wants the victims to learn a target (deterministic) policy π[†] as the unique Markov perfect (Nash) equilibrium.
- 2 The attacker can modify the rewards in the training set from r^{o} to r^{\dagger} .
- **3** The attacker minimizes the reward modification cost $\|r^{\dagger} r^{o}\|$, e.g. $\sum_{k=1}^{K} \sum_{t=1}^{T} \|r_{t}^{\dagger,(k)} r_{t}^{o,(k)}\|_{1}$.
- $\begin{array}{|c|c|c|} \bullet & \text{The attacker does not know } \hat{R} \text{ and } \hat{P}, \text{ but assumes} \\ & \left| \hat{R} R^{(\mathsf{MLE})} \right| < \rho^{(R)} \text{ and } \left\| \hat{P} P^{(\mathsf{MLE})} \right\|_1 < \rho^{(P)}. \end{array}$

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iNash Formulation

• The attack can be formulated as

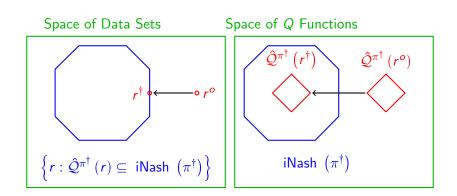
$$\begin{split} \min_{r^{\dagger}} \left\| r^{\dagger} - r^{o} \right\| \\ &\text{s.t.} \hat{\mathcal{Q}}^{\pi^{\dagger}} \left(r^{\dagger}; \rho^{(R)}, \rho^{(P)} \right) \subseteq \text{ iNash } (\pi^{\dagger}) , \end{split}$$

where,

- $\hat{\mathcal{Q}}^{\pi}\left(r\right)$ is the set of plausible Q functions computed based on r evaluated on π ,
- **2** iNash (π) is the inverse Nash polytope of Q functions such that π is the strict Markov perfect (Nash) equilibrium.



iNash Diagram



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Feasibility

Theorem

The attack is feasible if $\rho_t^{(R)}(s,a) + |\beta_t(s,a)| < \frac{1}{4T}, \forall t, s, and$ actions a such that $a_1 = \pi_{1,t}^{\dagger}(s)$ or $a_2 = \pi_{2,t}^{\dagger}(s)$.

• For example, if $\rho^{(R)} = 0$ and $\beta = \frac{c}{\sqrt{N_t(s,a)}}$, then the condition is a data coverage condition, $N_t(s,a) > 16cT^2$ for actions profiles in the same row or column as π^{\dagger} in the stage game matrices.

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Feasible Example

• Suppose $\pi^{\dagger} = (1, 1)$ in a stage game, then the following attack is feasible under the previous feasibility condition,

$a_1 ackslash a_2$	1	2	3	4
1	0.5	1	1	1
2	0	-	-	-
3	0	-	-	-
4	0	-	-	-

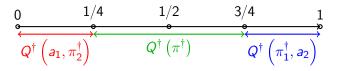
 Unspecified cells' corresponding rewards do not need to be poisoned.

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Feasibility Proof Sketch

- The condition $\rho_t^{(R)}(s, a) + |\beta_t(s, a)| < 1/(4T)$ implies that the cumulated confidence interval width for R and P in the future periods is bounded by 1/4.
- In period *t*, state *s*, for every $a_1 \neq \pi_1^{\dagger}$ and $a_2 \neq \pi_2^{\dagger}$, the *Q* values have the following relationship.



• Therefore, $\pi_t^{\dagger}(s)$ is the strict, thus unique, Nash equilibrium in every stage game (t, s).

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Linear Program Formulation

• The attacker's problem is given by,

$$\min_{r^{\dagger}} \sum_{k=1}^{K} \sum_{t=1}^{T} \left\| r_{t}^{\dagger,(k)} - r_{t}^{o,(k)} \right\|_{1}$$

s.t. for every t, s, and $Q_{t}^{\dagger} \in \hat{\mathcal{Q}}^{\pi^{\dagger}}\left(r^{\dagger}\right)$,

$$\begin{aligned} & Q_t^{\dagger}\left(s, \pi_t^{\dagger}\left(s\right)\right) > Q_t^{\dagger}\left(s, \left(a_1, \pi_{t,2}^{\dagger}\left(s\right)\right)\right), \forall \ a_1 \neq \pi_{t,1}^{\dagger}\left(s\right), \\ & Q_t^{\dagger}\left(s, \pi_t^{\dagger}\left(s\right)\right) < Q_t^{\dagger}\left(s, \left(\pi_{t,1}^{\dagger}\left(s\right), a_2\right)\right), \forall \ a_2 \neq \pi_{t,2}^{\dagger}\left(s\right). \end{aligned}$$

• Since $\hat{Q}^{\pi}(r)$ are polytopes, this problem can be formulated as a linear program and solved efficiently.



[my3] DSE Attack

- my3 Joint work (\approx 50% contribution) with Jeremy McMahan, Jerry Zhu, Qiaomin Xie. (Thanks: Yudong Chen)
 - The settings are similar to the Nash Attack [my2], except there are *n* victims learning general-sum Markov games.
 - iDSE (Markov perfect dominant strategy equilibrium) is used in place of iNash: the feasibility conditions are similar, and the attack can also be converted into a linear program.

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• [my4], [my5], [my6] are incomplete future work, and focus mostly on modes of attack other than training time reward poisoning.



[my4] MARL Test Attack

- The setting:
- The attacker knows the victims' pre-trained policy $\pi = (\pi_1, \pi_2, ..., \pi_n).$
- The attacker wants to minimize some function of the victims' rewards g ((r₁, r₂, ..., r_n)).

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• The attacker may poison the environment at test time, for example, modify the perceived states from s_t to s_t^{\dagger} .



• Single-agent test time attacks have been studied, but they can be extended to the multi-agent reinforcement learning setting.



Attacker Goal

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• The attacker wants to minimize some social welfare $g((r_1, r_2, ..., r_n))$ of the victims, for example,

3 Utilitarian:
$$g(r) = \sum_{i=1}^{n} r_i(\pi)$$
.

2 Rawlsian:
$$g(r) = \min_{i} r_i(\pi)$$
.

• Other functions of rewards such as $g(r) = \max_{i} r_i(\pi) - \min_{i} r_i(\pi).$

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Attacker Action

- The attacker may modify one of the following during test time,
- Perceived state, common to all victims, i.e. change $s_t \rightarrow s_t^{\dagger}$, but $P(s_{t+1}|a_t, s_t)$ stays the same.
- Perceived state, different to different victims, i.e. change $s_t \rightarrow \left(s_{t,1}^{\dagger}, s_{t,2}^{\dagger}, ..., s_{t,n}^{\dagger}\right), \text{ but } P\left(s_{t+1}|a_t, s_t\right) \text{ stays the same.}$

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- **3** True state, i.e. change $s_t \rightarrow s_t^{\dagger}$, and $P(s_{t+1}|a_t, s_t) \rightarrow P(s_{t+1}|a_t, s_t^{\dagger})$.
- Victim action, i.e. change $a_t \rightarrow a_t^{\dagger}$, and $P(s_{t+1}|a_t, s_t) \rightarrow P(s_{t+1}|a_t^{\dagger}, s_t)$.



Roadmap to Solve [my4]

- The original Markov game, given by its states, actions, transitions, and rewards, say G = (S, A, P, R).
- In the perceived common state attack, the attacker's problem can be formulated as a meta Markov decision process $\mathcal{M} = (\mathcal{S}', \mathcal{A}', P', R')$, where,
- The meta states S' = S.
- 2 The meta actions $\mathcal{A}' \subseteq \mathcal{S}$.
- **3** The meta transitions $P'_t \left((s_{t+1}, a_{t+1}) | (s_t, a_t), s_t^{\dagger} \right) = P_t \left(s_{t+1} | s_t^{\dagger}, a_t \right) \pi_{t+1} (a_{t+1} | s_{t+1}).$
- The meta rewards $R'_t\left((s_t, a_t), s^{\dagger}_t\right) = -g\left(r_t\left(s^{\dagger}_t, a_t\right)\right).$



Roadmap, Continued

- The meta MDP \mathcal{M} can be solved using any reinforcement learning or planning algorithms.
- There might be special algorithms to solve \mathcal{M} more efficiently since the meta action space might be large.
- The setting where the attacker does not know ${\mathcal G}$ or π could be studied.

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• Experiments could be implemented.



[my5] Equilibrium Defense

- The setting:
- The attacker wants to minimize some social welfare g ((r₁, r₂, ..., r_n)).
- The victims want to maximize expected discounted individual rewards r_i.
- The attacker and victims simultaneously select and commit to a perceived state attack ν : S → S and a policy π = (π₁, π₂, ..., π_n), with π_i : S → A_i.

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Motivation

 In MARL Test Attack [my4] and most of the attack-defense literature, either the victim has a fixed policy, or the attacker has a fixed attack algorithm, and the other agent best responds to the fixed action. Both models are not realistic and equilibrium attack-defense should be studied instead.

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Roadmap to Solve [my5]

- The problem can be formulated as a static game $\mathcal{G} = (\mathcal{A}', \mathcal{R}')$, where,
- The meta actions $\mathcal{A}' = (\mathcal{S} \to \mathcal{S}, \mathcal{S} \to \mathcal{A}_1, \mathcal{S} \to \mathcal{A}_2, ..., \mathcal{S} \to \mathcal{A}_n).$
- The meta rewards
 R' (ν, π₁, π₂, ..., π_n) = (−g (V (π (ν))), V (π (ν))), where
 V (π (ν)) = ∑_{t=1}[∞] γ^t 𝔼_P [R (s_t, π₁ (ν (s_t)), ..., π_n (ν (s_t)))].

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Roadmap, Continued

- The meta game \mathcal{G} can be solved using a Nash solver, e.g. a linear program when the game is zero-sum.
- There might be special classes of equilibria that are easier to solve since the meta action space might be large.
- The equilibrium policy π might correspond to some robust policy in the RL literature.
- Training time equilibrium attack defense could also be studied.

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[my6] Multi Attacker

- The setting:
- Multiple attackers j ∈ [m], each attacks a subset of the victims.
- Each attacker wants to minimize a different social welfare g_j ((r₁, r₂, ..., r_n)).
- Seach attacker may modify the perceived state of the set of victims it attacks.

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• In MARL Test Attack [my4] and most of the training or test time attack literature, there is only one attacker. The problem with multiple attackers with different objectives is an interesting problem with many real-world applications.

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Roadmap to Solve [my6]

- The original Markov game, given by its states, actions, transitions, and rewards, say G = (S, A, P, R).
- In the perceived state attack where each attacker attacks a single victim, the attackers' problem can be formulated as a meta Markov game $\mathcal{M} = (\mathcal{S}', \mathcal{A}', \mathcal{P}', \mathcal{R}')$, where,

• The meta states
$$S' = S$$
.

- 2 The meta actions $\mathcal{A}' \subseteq \mathcal{S}^{n}$.
- **3** The meta transitions $P'_t \left((s_{t+1}, a_{t+1}) | (s_t, a_t), s_t^{\dagger} \right) = P_t \left(s_{t+1} | s_t^{\dagger}, a_t \right) \pi_{t+1} (a_{t+1} | s_{t+1}).$

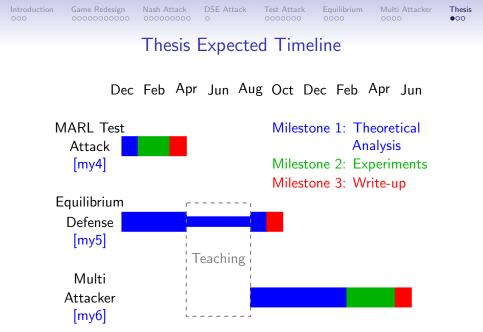
• The meta rewards $R'_t\left((s_t, a_t), s^{\dagger}_t\right) = -g\left(r_t\left(s^{\dagger}_t, a_t\right)\right).$



Roadmap, Continued

- The meta Markov Game \mathcal{G} can be solved using any multi-agent reinforcement learning or planning algorithms.
- There might be special algorithms to solve \mathcal{G} more efficiently since the meta action space might be large.
- The setting where the attackers do not know ${\mathcal G}$ or π could be studied.
- Experiments could be implemented.
- Training time attacks with multiple attackers could also be studied.

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Thank you!

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