WISCONSIN
UNIVERSITY OF WISCONSIN–MADISON

CS 764: Topics in Database Management Systems

Lecture 2: Join

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Today's Paper: Join

Join Processing in Database Systems with Large Main Memories

LEONARD D. SHAPIRO
North Dakota State University

We study algorithms for computing the equijoin of two relations in a system with a standard architecture but with large amounts of main memory. Our algorithms are especially efficient when the main memory available is a significant fraction of the size of one of the relations to be joined; but they can be applied whenever there is memory equal to approximately the square root of the size of one relation. We present a new algorithm which is a hybrid of two hash-based algorithms and which dominates the other algorithms we present, including sort-merge. Even in a virtual memory environment, the hybrid algorithm dominates all the others we study.

Finally, we describe how three popular tools to increase the efficiency of joins, namely filters, Babb arrays, and semijoins, can be grafted onto any of our algorithms.

Categories and Subject Descriptors: H.2.0 [**Database Management**]: General; H.2.4 [**Database Management**]: Systems—*query processing*; H.2.6 [**Database Management**]: Database Machines

General Terms: Algorithms, Performance

Additional Key Words and Phrases: Hash join, join processing, large main memory, sort-merge join

ACM Transactions on Database Systems, 1986

Agenda

System architecture and notations

Join algorithms

- Sort merge join
- Simple hash join
- GRACE hash join
- Hybrid hash join

Partition overflow and additional techniques

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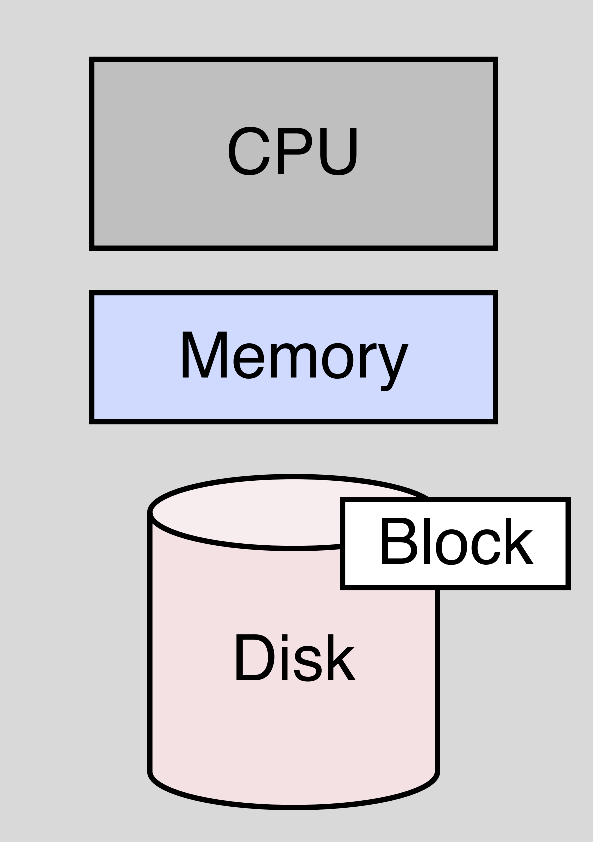
System Architecture and Assumptions



CPU: uniprocessor

- No multi-core synchronization complexity
- Could be built on systems of the day

Memory

- Tens of Megabytes
- Good for **both sequential and random** accesses
- Capacity is smaller than disk

Disk

- Good for **only sequential** accesses

Notation

Relations: R, S (| R | < | S |)

Join: S ⋈ R

Memory: M

| R |: number of blocks in relation R (similar for S and M)

F: hash table for R occupies | R | * F blocks

Focus only on equi-join

Notation

Relations: R, S ($|R| < |S|$)

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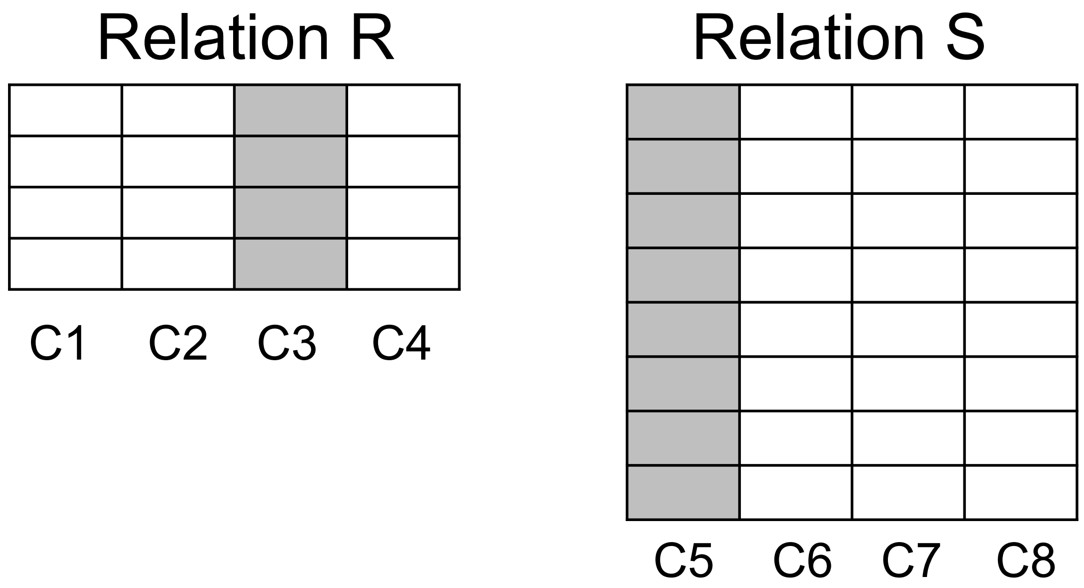


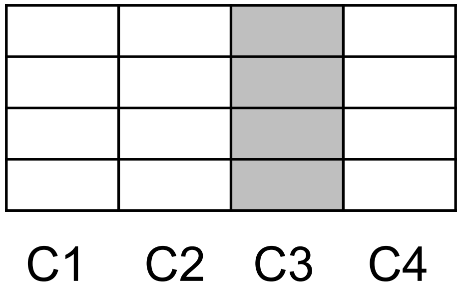

```
SELECT *
FROM R, S
WHERE R.C3 = S.C5
```

Notation

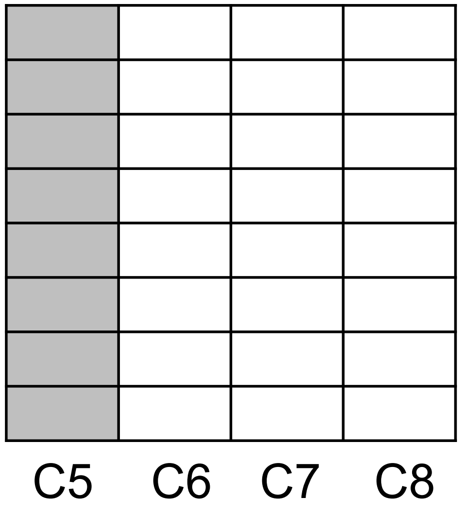
```
answer = {}
for t1 in R do
  for t2 in S do
    if R.C3 = S.C5
      then answer = answer ∪ {(C1,…,C8)}
return answer
```

Vanilla query executor

Relation R



Relation S



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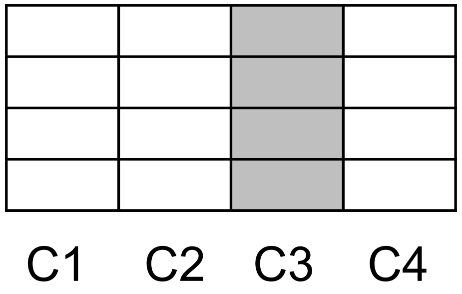
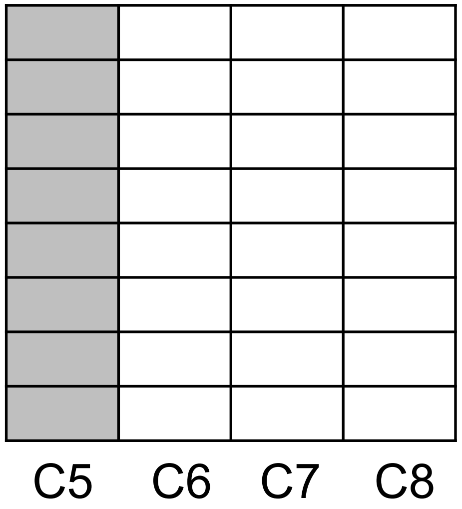
Vanilla query executor

Key question: How to execute a join fast?

Relation R



Relation S

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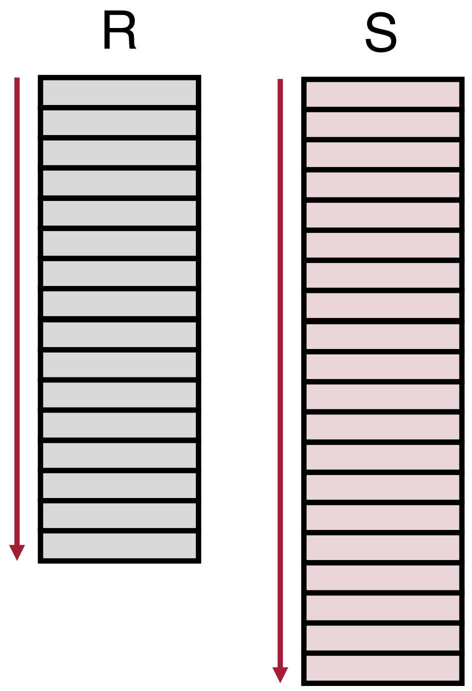
Partition overflow and additional techniques

Sort Merge Join

Key idea: sort both relations based on join attributes, then traverse both relations in the sorting order



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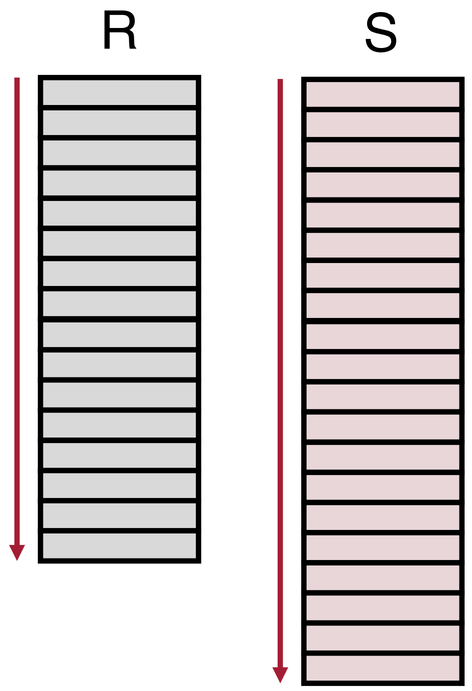
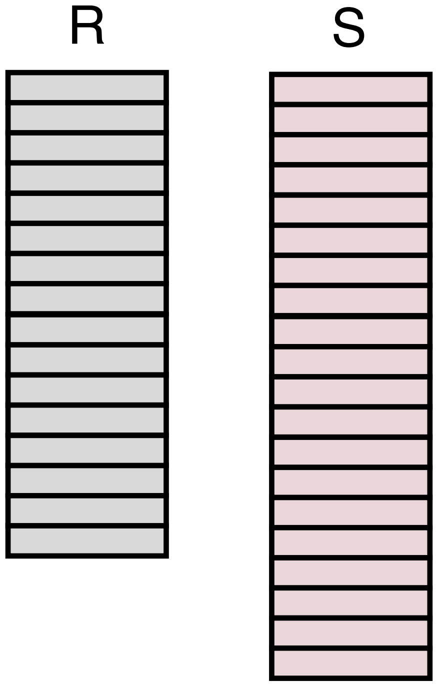


Challenge: If a relation does not fit in memory, need to sort data on disk

Sort Merge Join

Phase 1: Produce sorted runs of S and R

Phase 2: Merge runs of S and R, output join result



Unsorted R and S

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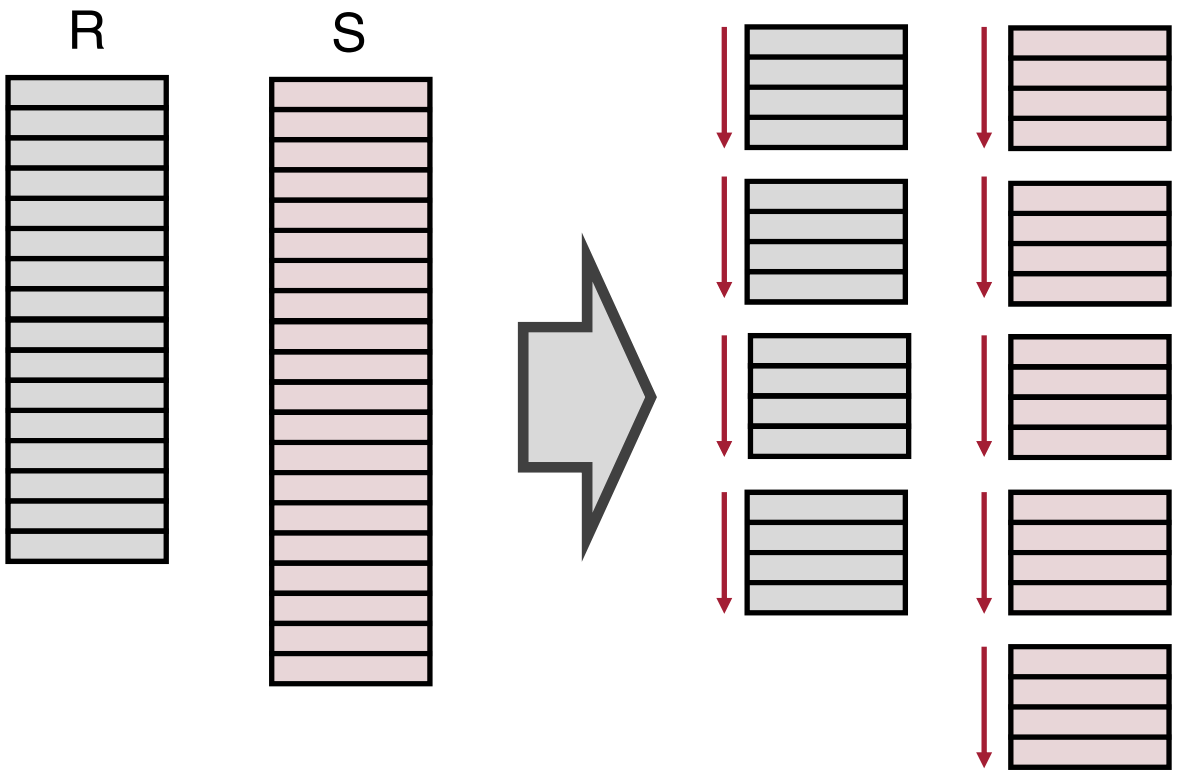

Each sorted run can fit in memory

Unsorted R and S

Sorted runs of R and S

Sort Merge Join

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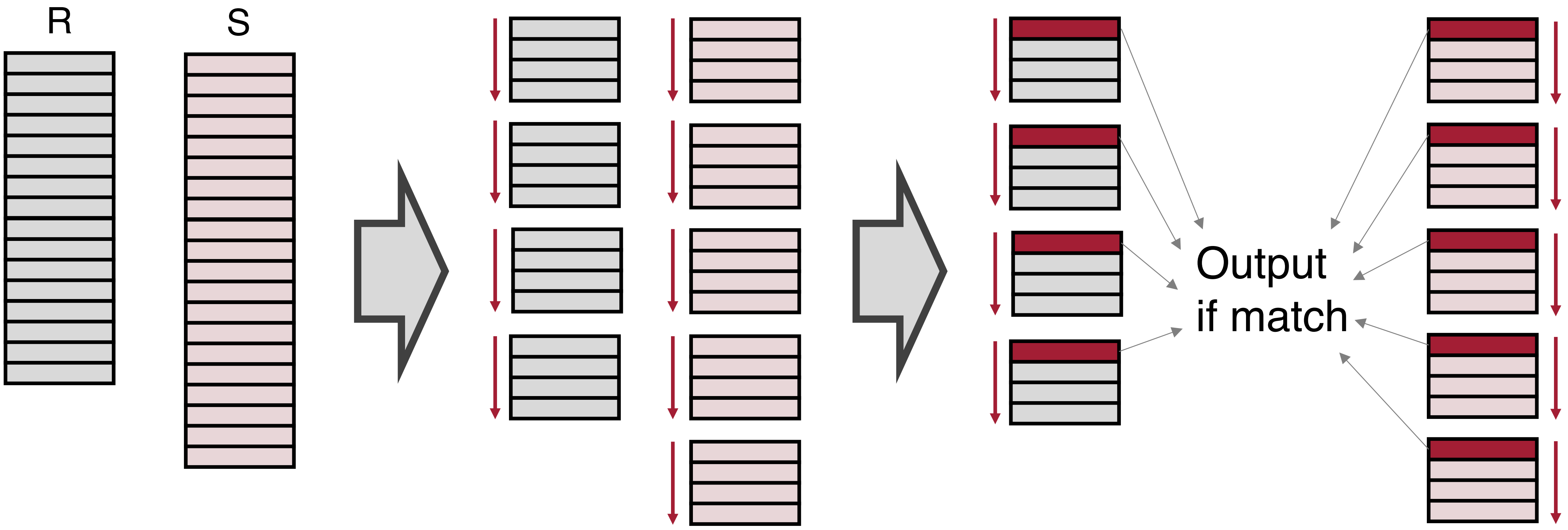


Unsorted R and S

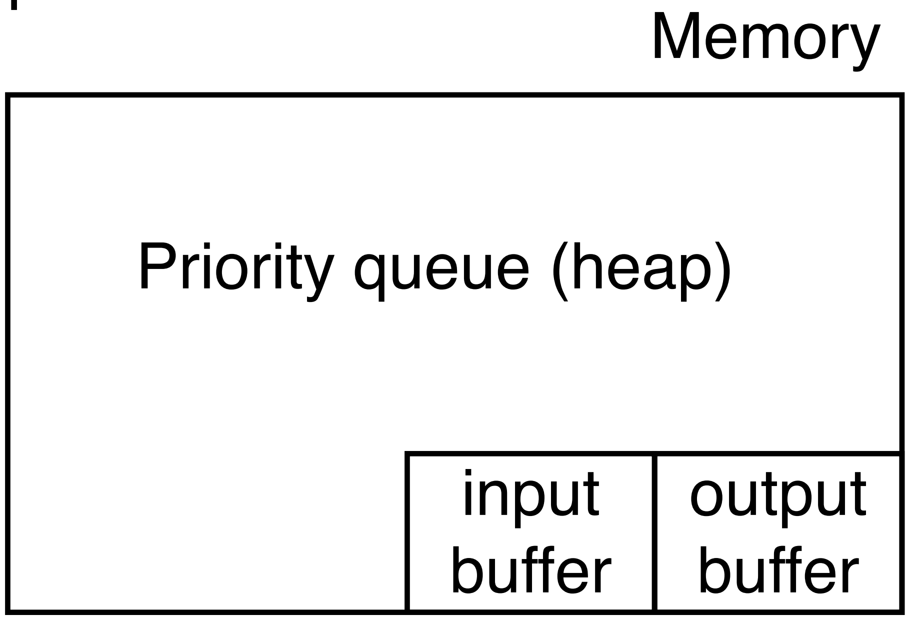
Sorted runs of R and S

Find matches in sorted runs

Sort Merge Join – Phase 1

Phase 1: Produce sorted runs of S and R

- Each run of S will be **2 × | M |** average length



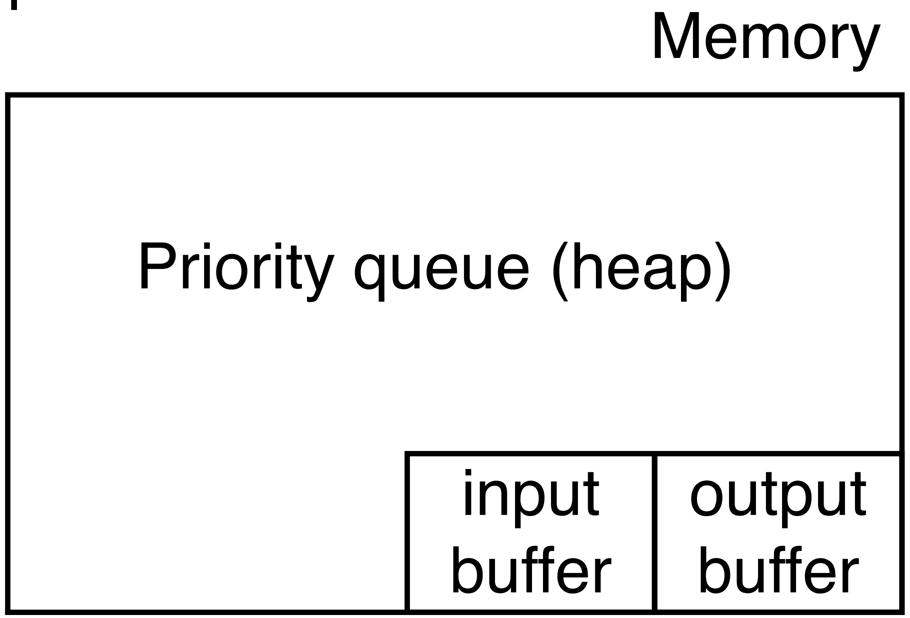
Memory layout in Phase 1

Sort Merge Join – Phase 1

Phase 1: Produce sorted runs of S and R

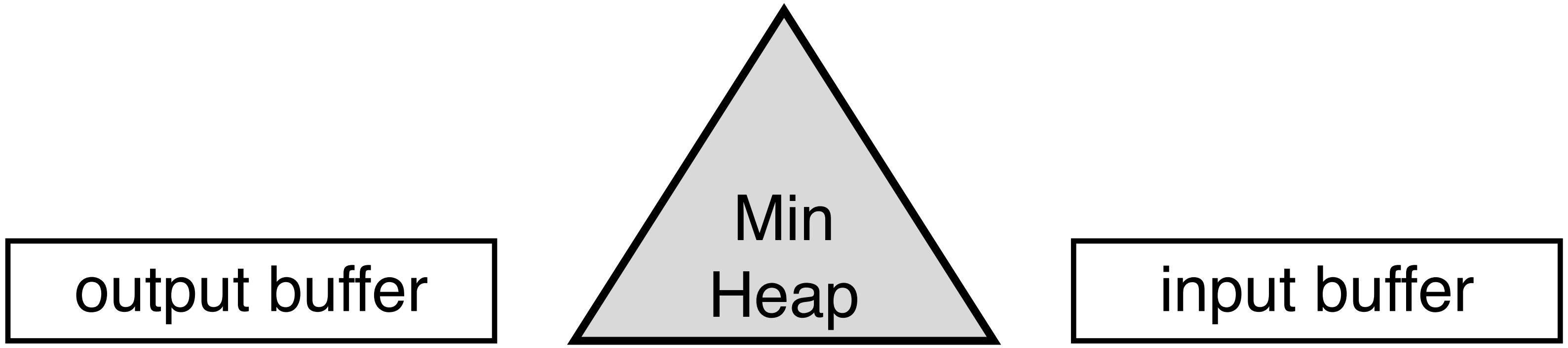
- Each run of S will be **2 × | M |** average length

Q: Where does 2 come from?
A: Replacement selection



Memory layout in Phase 1

Sort Merge Join – Replacement Selection

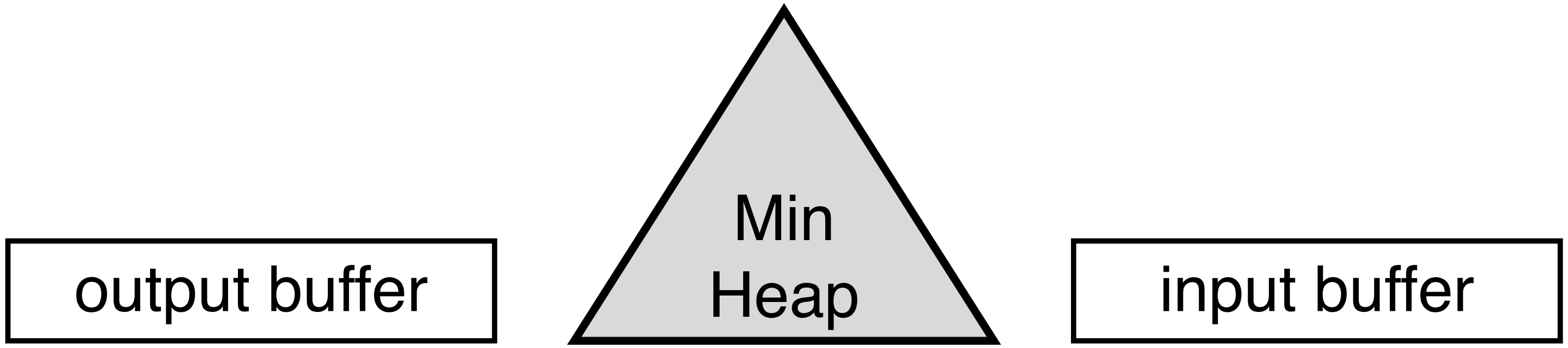


Naïve solution:

- Load $|M|$ blocks
- Sort
- Output $|M|$ blocks

Each run contains $|M|$ blocks

Sort Merge Join – Replacement Selection

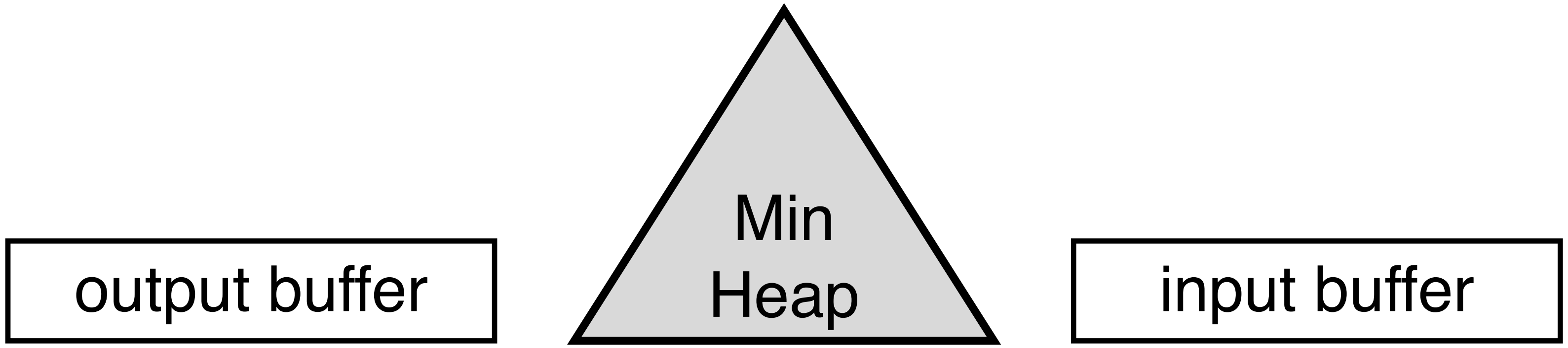


Replacement selection:

- load | M | blocks and sort

```
While heap is not empty
    If new tuple ≥ all tuples in output
        add new tuple to heap
    else
        save new tuple for next run
```

Sort Merge Join – Replacement Selection



Replacement selection:

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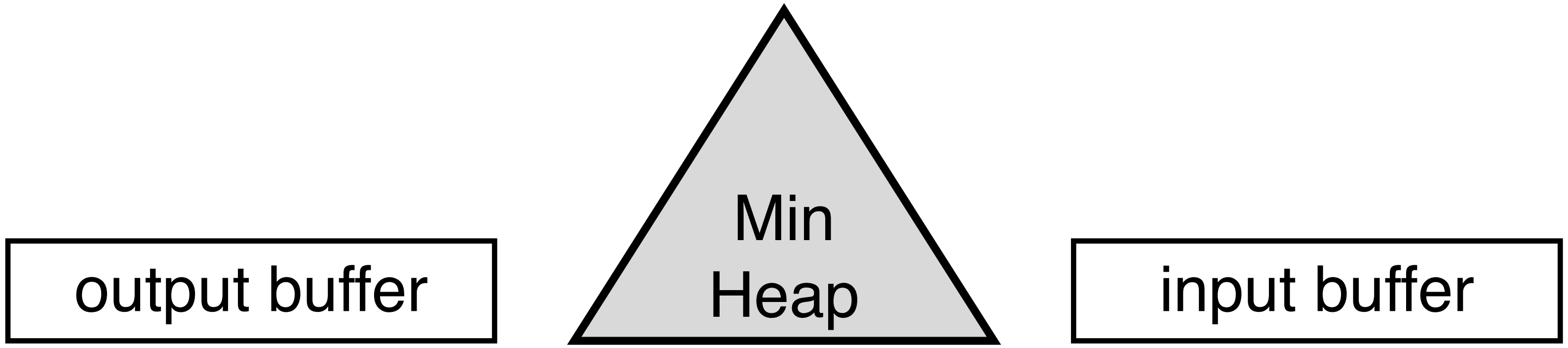
Sort Merge Join – Replacement Selection



Replacement selection:

- load | M | blocks and sort

A run contains 2 × | M | blocks on average

While heap is not empty
    If new tuple ≥ all tuples in output
        add new tuple to heap
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Total number of runs

$$= \frac{|S|}{2 \times |M|} + \frac{|R|}{2 \times |M|} \leq \frac{|S|}{|M|}$$

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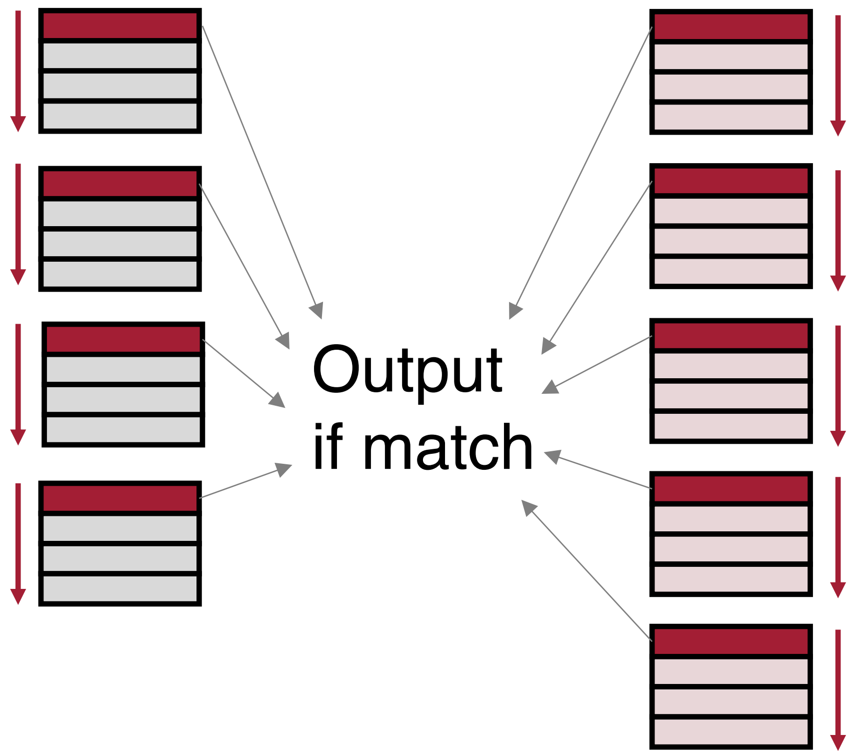
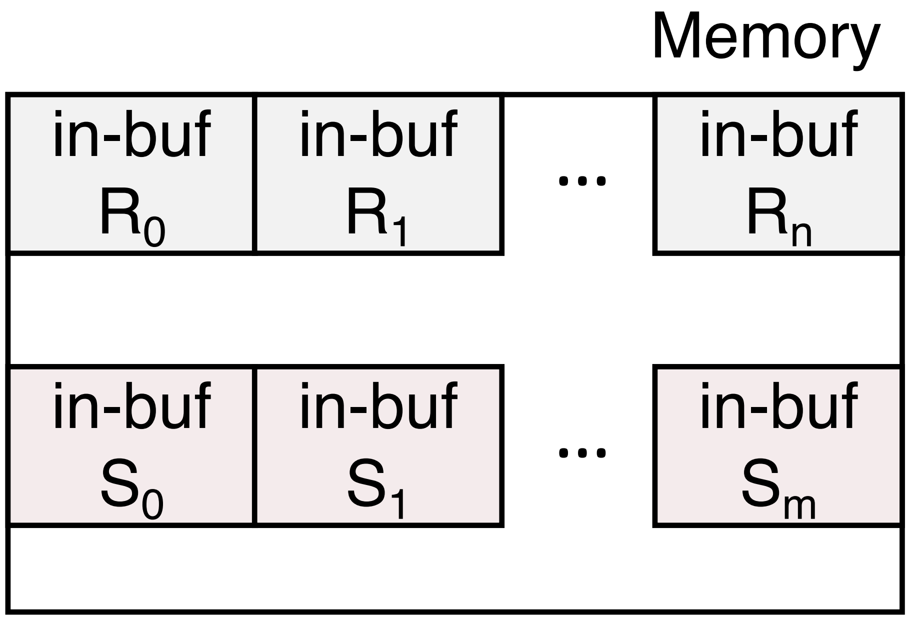
Sort Merge Join – Phase 2

Phase 2: Merge runs of S and R, output join result

- One input buffer required for each run



Find matches in sorted runs



Memory layout in Phase 2

Sort Merge Join – Phase 2

Phase 2: Merge runs of S and R, output join result

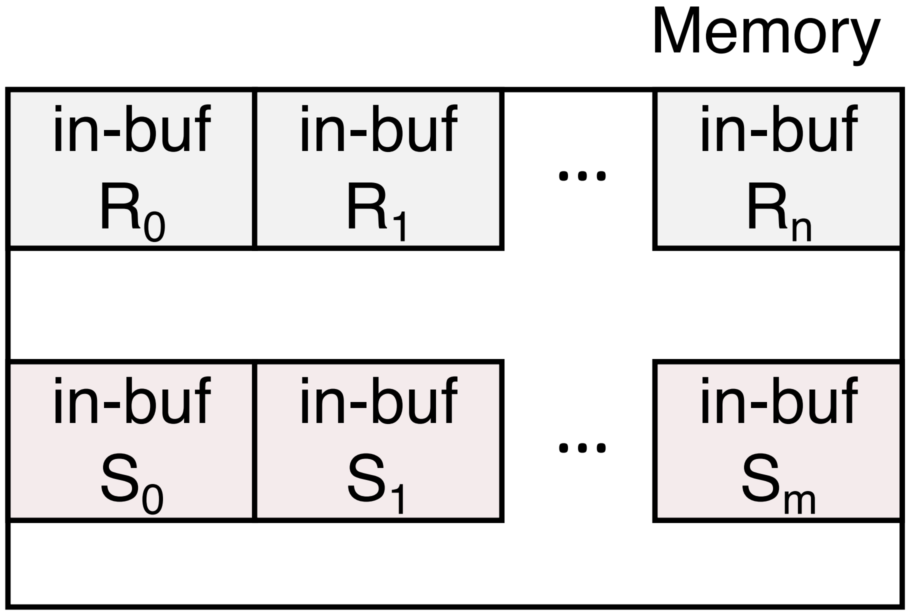
- One input buffer required for each run

Requirement

$$| M | \geq \text{total number runs}$$

Satisfied if $| M | \geq \frac{| S |}{| M |}$

namely $| M | \geq \sqrt{| S |}$



Memory layout in Phase 2

Hash Join

Build a hash table on the smaller relation (**R**) and probe with larger (**S**)

Hash tables have overhead, call it **F**

When **R** doesn't fit fully in memory, partition hash space into ranges

Hash table on **R**
(size = | **R** | × **F**)

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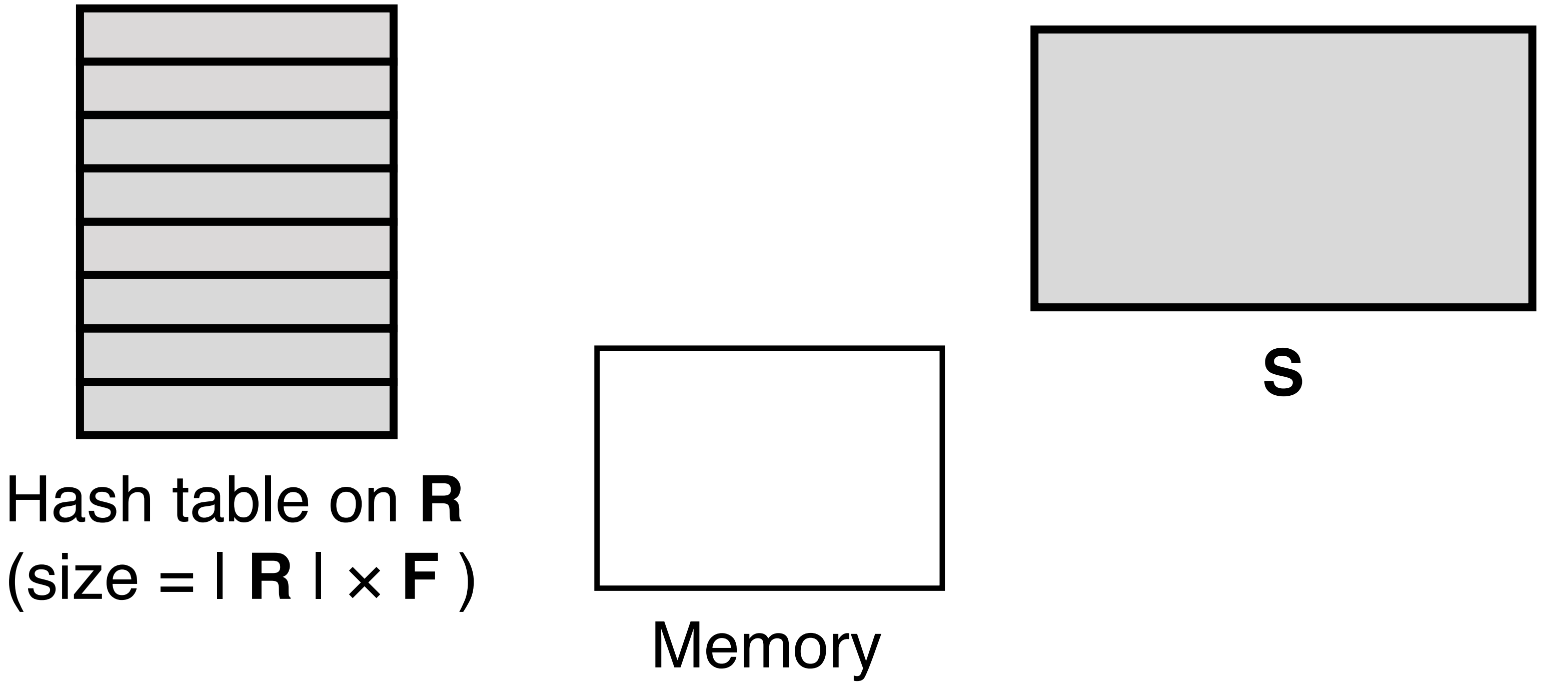
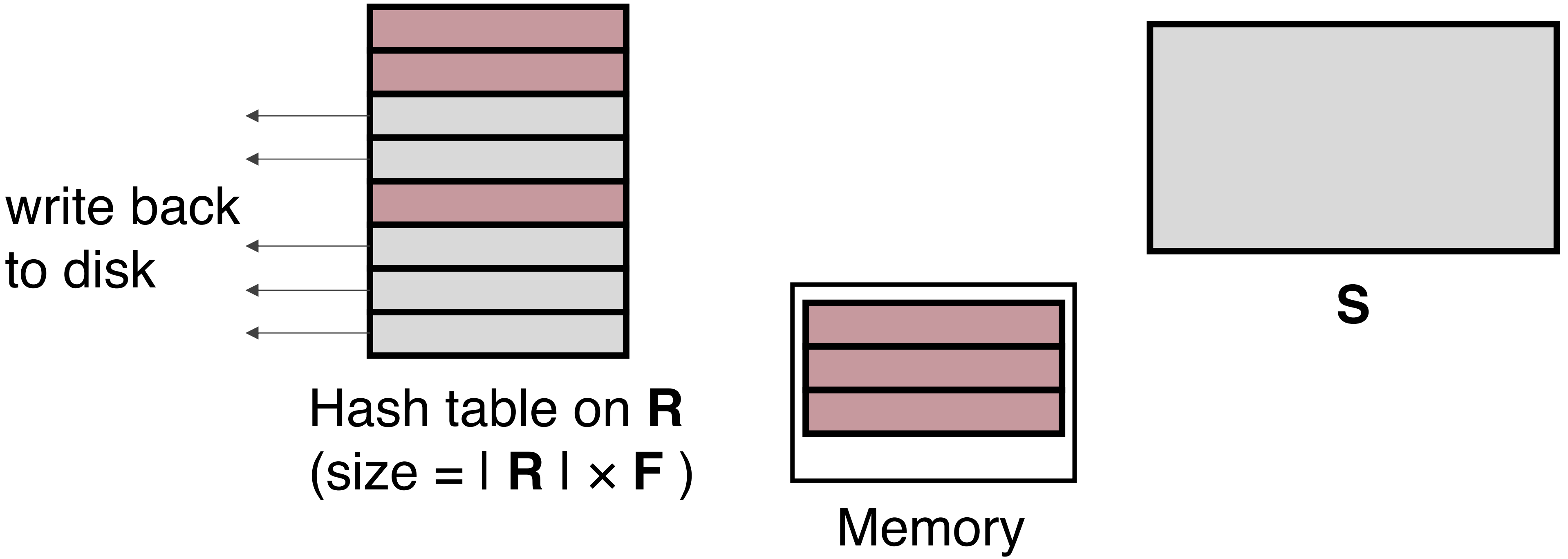
Simple Hash Join

- Build a hash table on **R**



Simple Hash Join – 1st pass

- Build a hash table on **R**
- If **R** does not fit in memory, find a subset of buckets that fit in memory



Simple Hash Join – 1st pass

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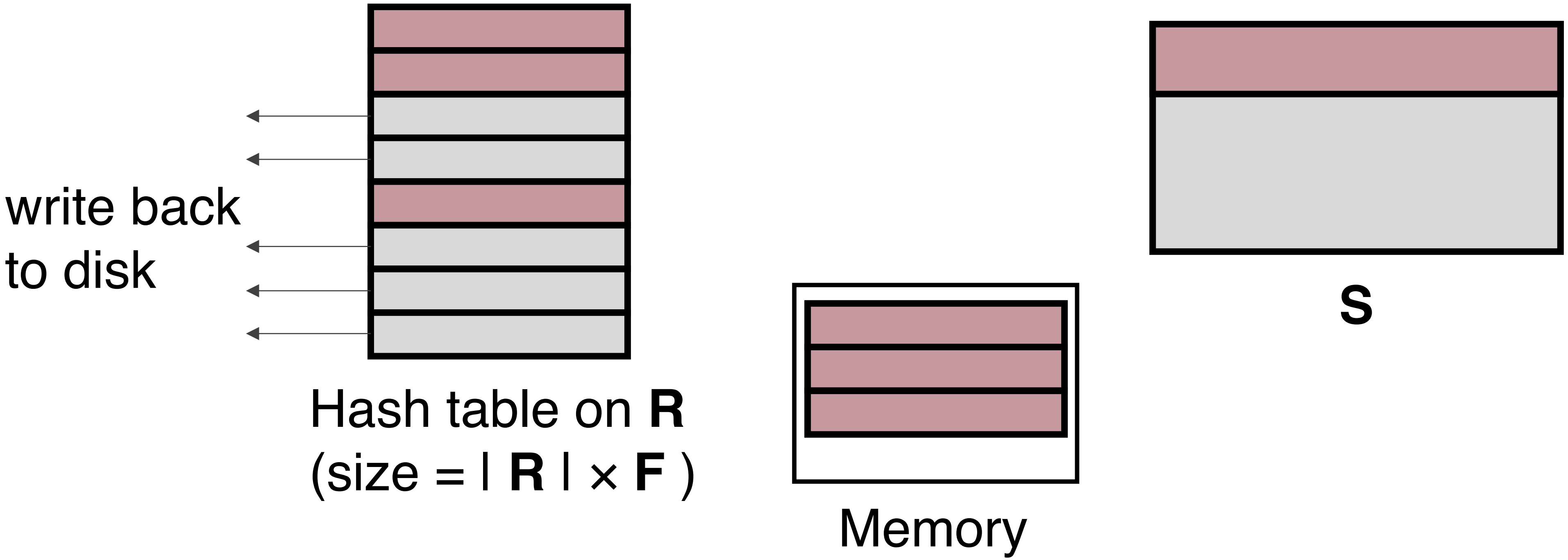
- Read in **S** to join with the subset of **R**



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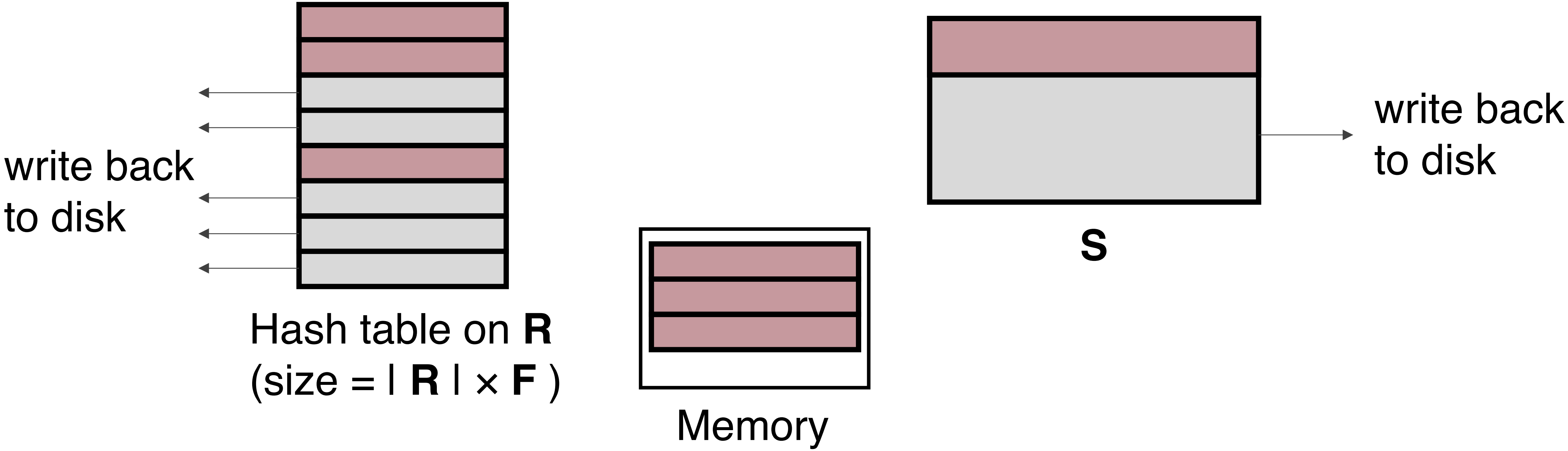
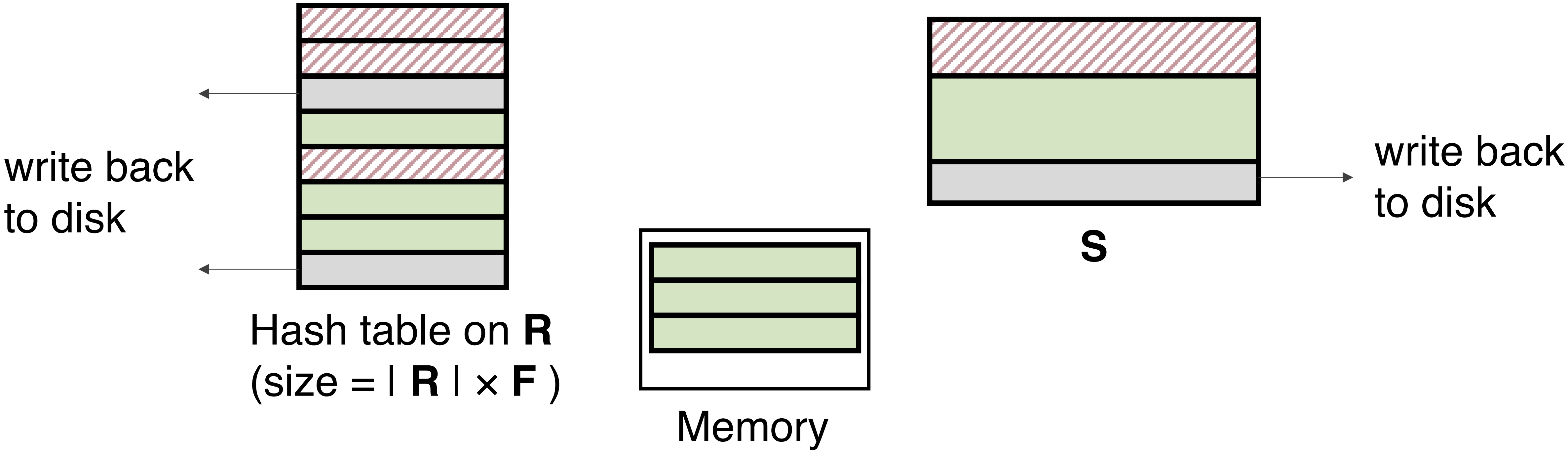


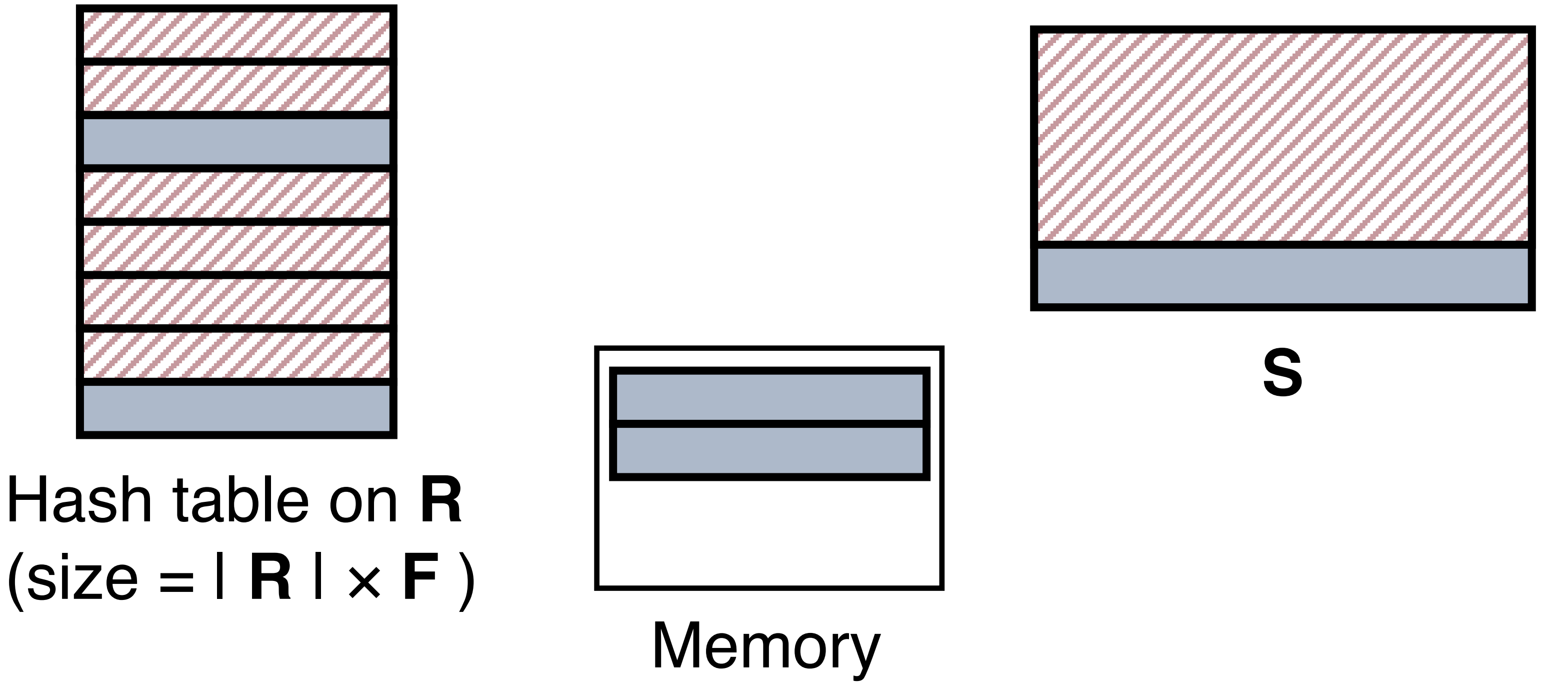
Simple Hash Join – 2nd pass

- Build a hash table on **R**
- If **R** does not fit in memory, find a subset of buckets that fit in memory
- Read in **S** to join with the subset of **R**
- The remaining tuples of **S** and **R** are written back to disk



Simple Hash Join – 3rd pass

- Build a hash table on **R**
- If **R** does not fit in memory, find a subset of buckets that fit in memory
- Read in **S** to join with the subset of **R**
- The remaining tuples of **S** and **R** are written back to disk



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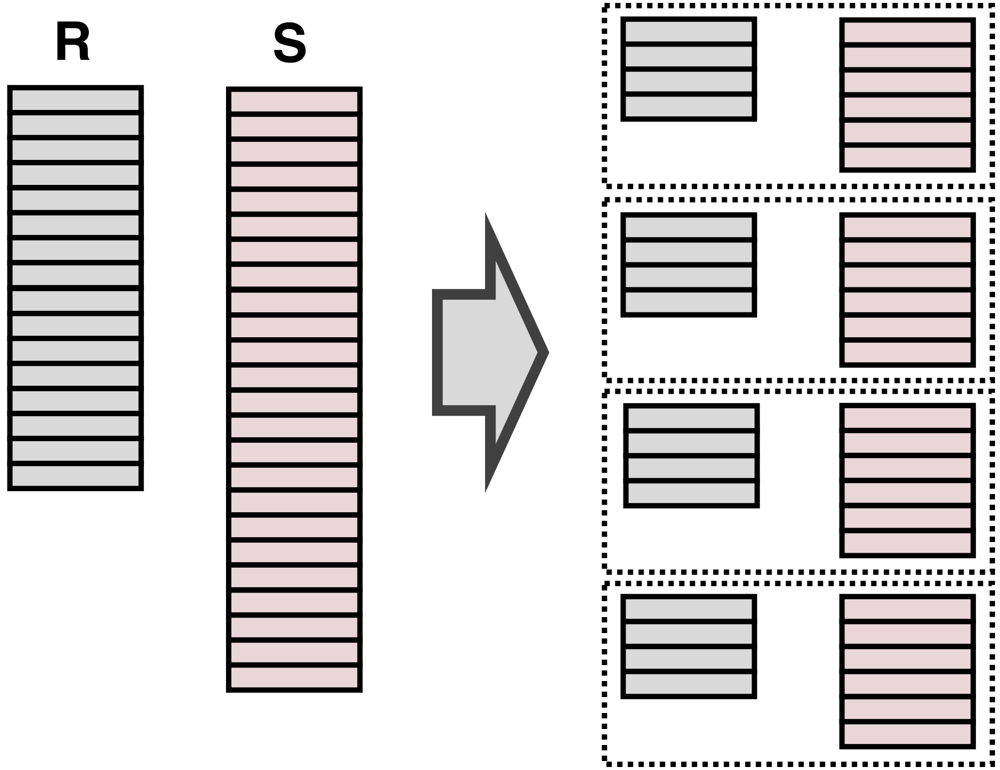
GRACE Hash Join

Phase 1: Partition both R and S into pairs of k shards

Phase 2: Separately join each pairs of partitions



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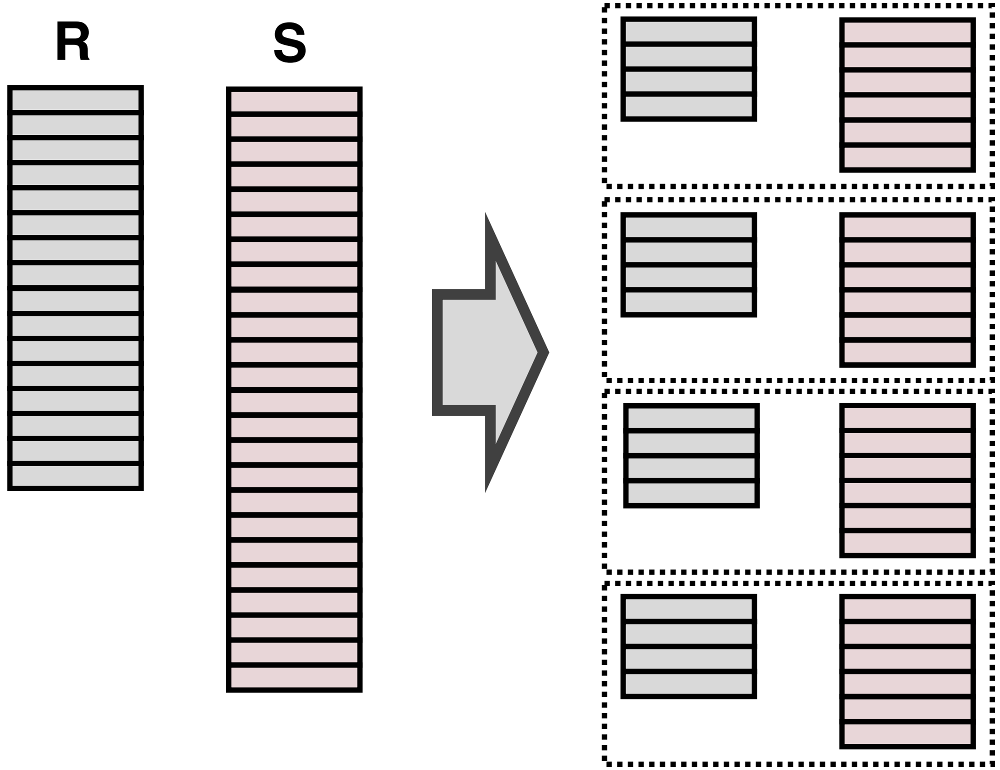
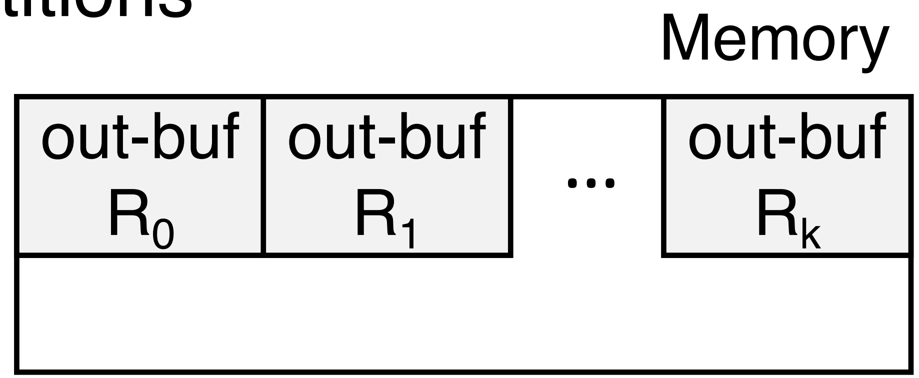
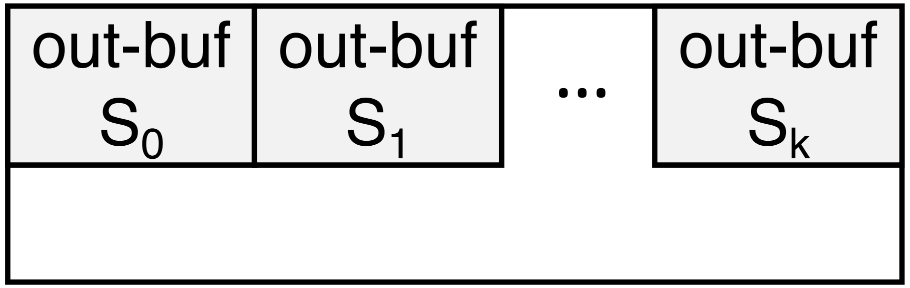




Memory layout when Partitioning R



Memory layout when Partitioning S

GRACE Hash Join

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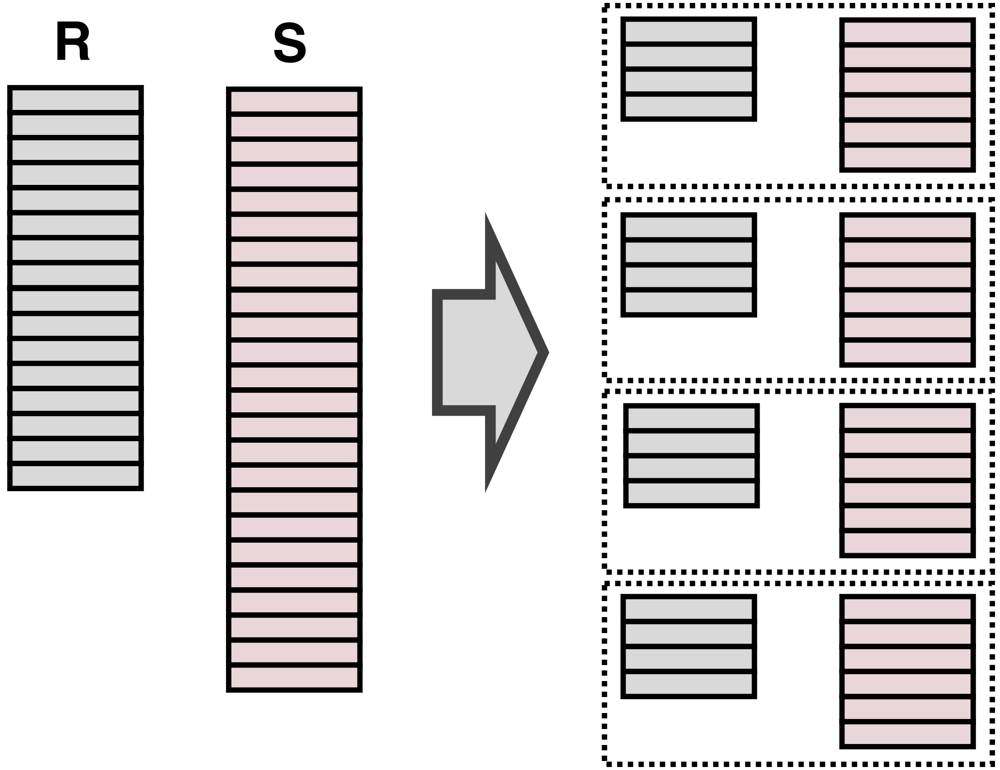
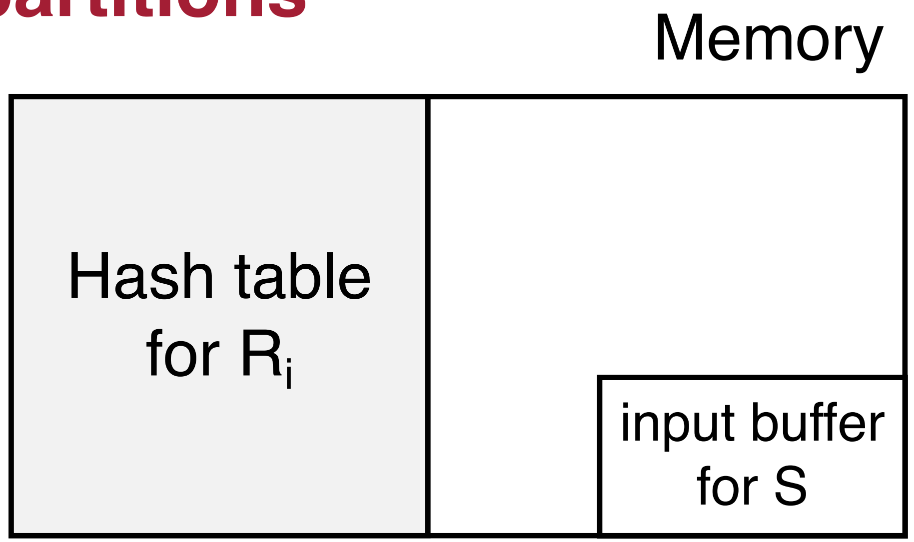
Phase 2: Separately join each pairs of partitions





Memory layout in Phase 2

GRACE Hash Join

Assume **k** partitions for **R** and **S**

In phase 1, needs one output buffer (i.e., block) for each partition

$$k \leq |M|$$

GRACE Hash Join

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In phase 2, the hash table of each shard of **R** must fit in memory

$$\frac{|R|}{k} \times F \leq |M|$$

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In phase 2, the hash table of each shard of **R** must fit in memory

$$\frac{|R|}{k} \times F \leq |M|$$

The maximum size of **R** to perform Grace hash join:

$$|R| \leq \frac{|M|}{F} k \leq \frac{|M|^2}{F} \qquad |M| \geq \sqrt{|R| \times F}$$

GRACE vs. Simple Hash Join

When $| R | \times F < | M |$

- Simple hash join incurs no IO traffic (better)
- GRACE hash join writes and reads each table once
- Trivial optimization to GRACE: use simple hash join when $| R | \times F < | M |$

When $| M |^2 \geq | R | \times F >> | M |$

- Simple hash join incurs significant IO traffic
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Discussion Question:
What if $| R | \times F > | M |^2$?

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Partition overflow and additional techniques

Hybrid Hash Join

When two algorithms are good in different settings, create a hybrid!

Hybrid Hash Join

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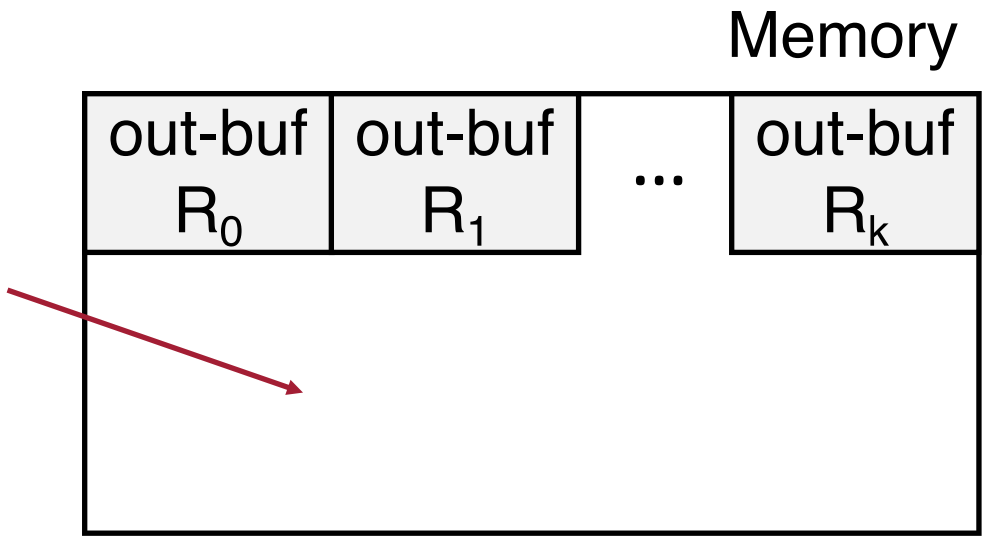
Key observation: when | R | is relatively small (e.g., | R | = 2 | M |), significant memory capacity is unused in Phase 1 of GRACE join



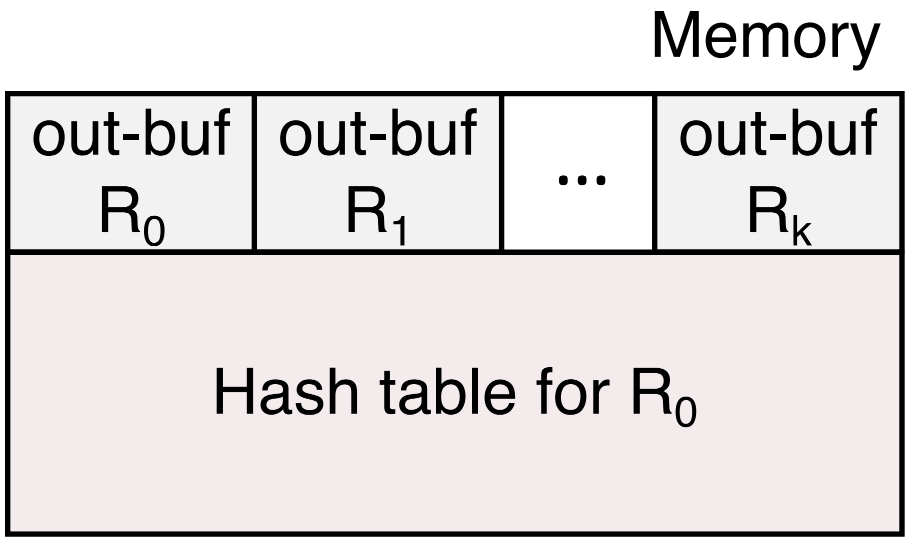
Memory layout in Phase 1 of GRACE hash join

Hybrid Hash Join

When two algorithms are good in different settings, create a hybrid!

Key observation: when | R | is relatively small (e.g., | R | = 2 | M |), significant memory capacity is unused in Phase 1 of GRACE join

Key idea: Use the otherwise-unused memory to build hash table for R_0

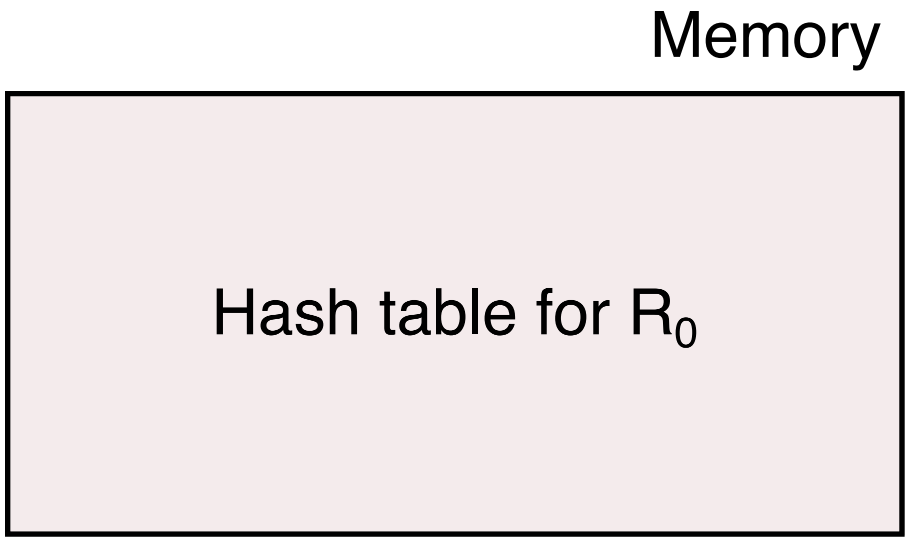


Memory layout in Phase 1 of GRACE hash join

Hybrid Hash Join

Case 1: $|R| \times F < |M|$

- No need to partition R
- Identical to simple hash join



Memory layout in Phase 1 of hybrid hash join

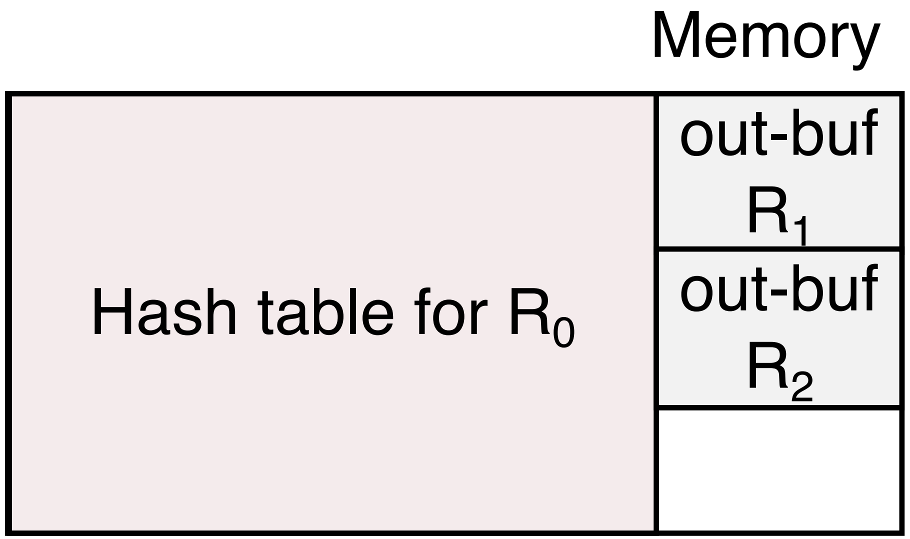
Hybrid Hash Join

Case 1: | R | × F < | M |

- No need to partition R
- Identical to simple hash join

Case 2: | R | × F = α | M | (α is small)

- R_0 is a significant fraction of R
- R_0 is not written to disk
- Performance is like simple hash join



Memory layout in Phase 1 of hybrid hash join

Hybrid Hash Join

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- R_0 is a significant fraction of R
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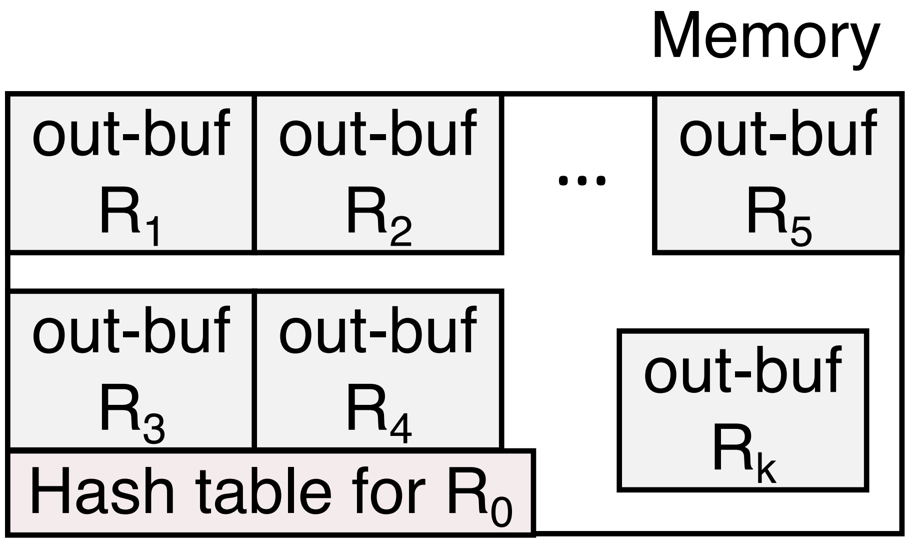
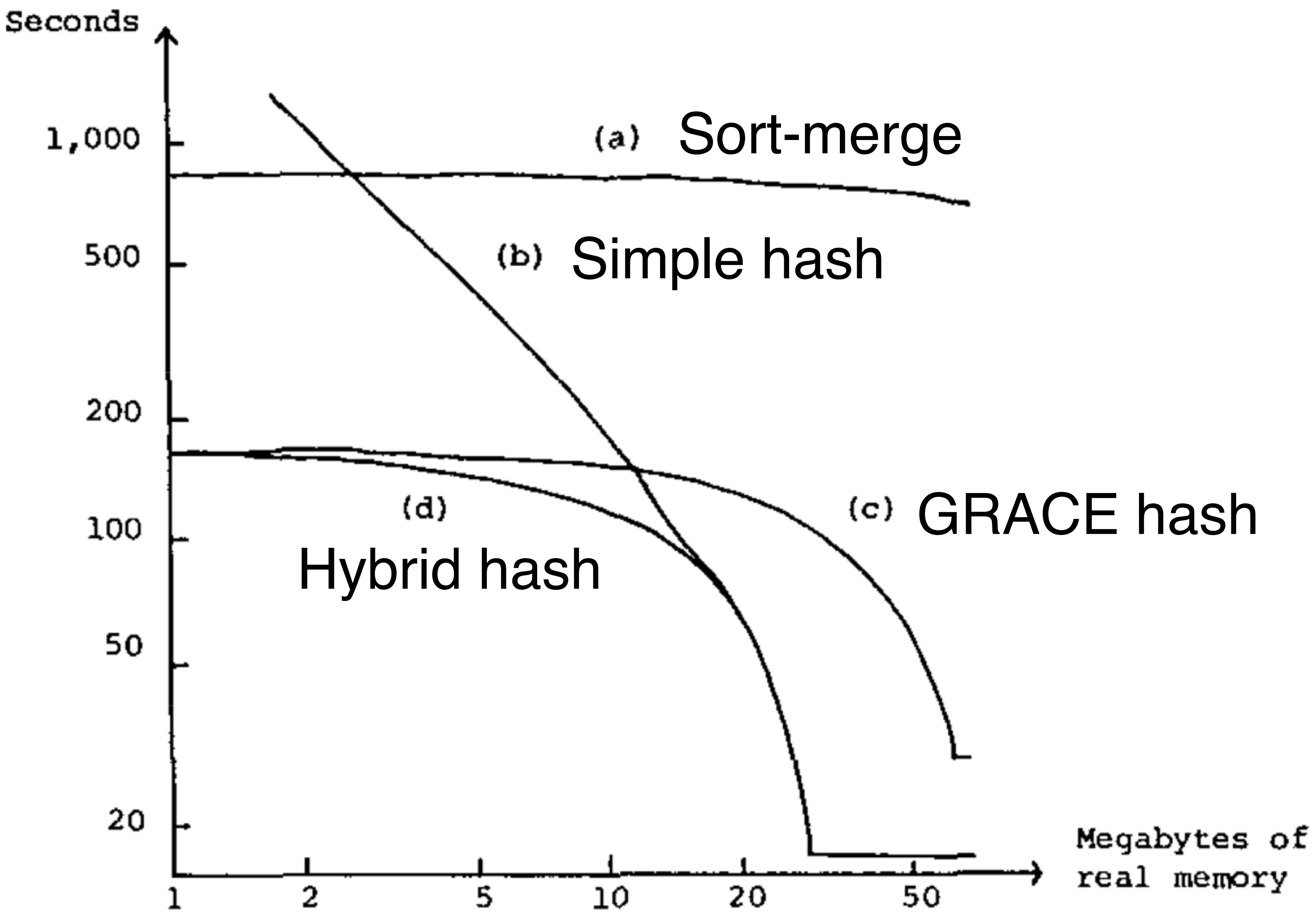
Case 3: | R | × F >> | M |

- R_0 is an insignificant fraction of R
- Performance is like GRACE hash join



Memory layout in Phase 1 of hybrid hash join

Evaluation



Conclusion 1: Hash join is generally better than sort-merge join

Conclusion 2: Hybrid hash join is strictly better than simple and GRACE hash joins

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Partition overflow and additional techniques

Partition Overflow

So far we assume uniform random distribution for **R** and **S**

What if we guess wrong on size required for R hash table and a partition does not fit in memory?

Solution: further divide into smaller partitions range

Additional Techniques

Babb array (or bitmap filter)

- One bit per hash bucket in R
- Set the bit if a tuple in R maps to the bucket
- When scanning S, if a tuple hashes to a bucket where the bit is unset, can discard the tuple immediately

Additional Techniques

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- One bit per hash bucket in R
- Set the bit if a tuple in R maps to the bucket
- When scanning S, if a tuple hashes to a bucket where the bit is unset, can discard the tuple immediately

Semi-join

- Project join attributes from R, join to S, then join that result back to R
- Useful if full R tuples won't fit into memory, but join will be selective and filter many S tuples
- Can be added to any join algorithm above

Join – Comments and Q/A

- How will the join algorithms change in parallel system?
- Is simple hash better since modern systems have large memories?
- Is the assumption l M l > sqrt(l S l) realistic?
- How to select a good hash function?
- Babb arrays used in practice?
- How do new storage devices (e.g., PM, SSD, tiered memory) change the story?
- Difficult to understand math.
- Lack of experiments.

Group Discussion

In some modern in-memory DBMSs, the entire database can fit in memory. In such a system, can similar optimizations be applied to on-chip SRAM caches vs. DRAM? What are the key challenges compared to a DRAM vs. Disk setting?

Before Next Lecture

Submit review for

Peter Boncz, et al., Database Architecture Optimized for the new Bottleneck: Memory Access. VLDB, 1999