



FoundationDB

**A Distributed Unbundled Transactional
Key Value Store**

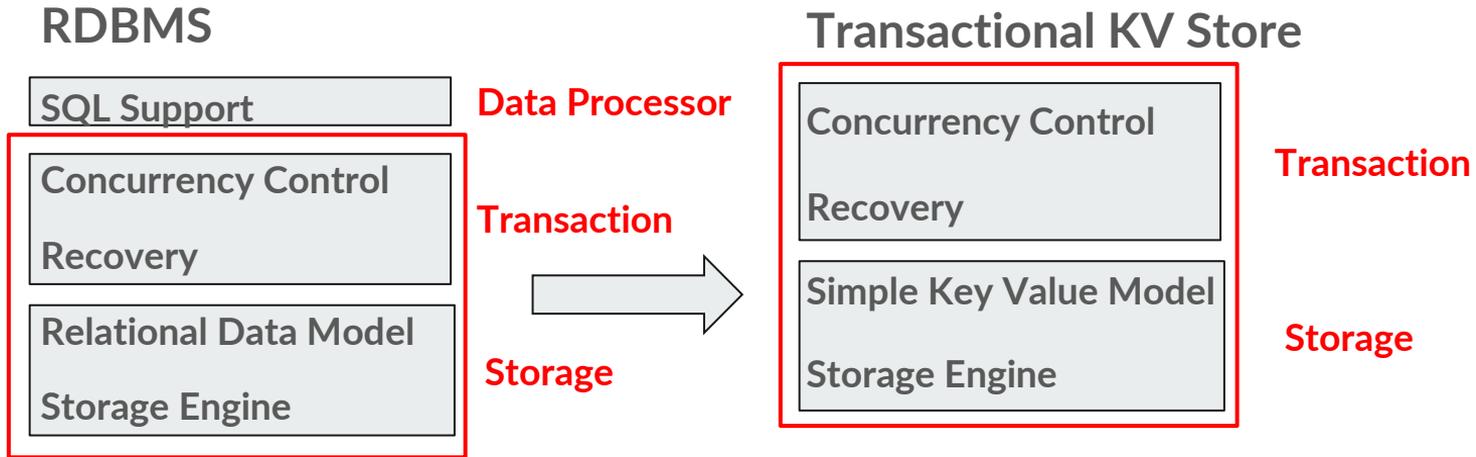
Yuhan Wang



Agenda

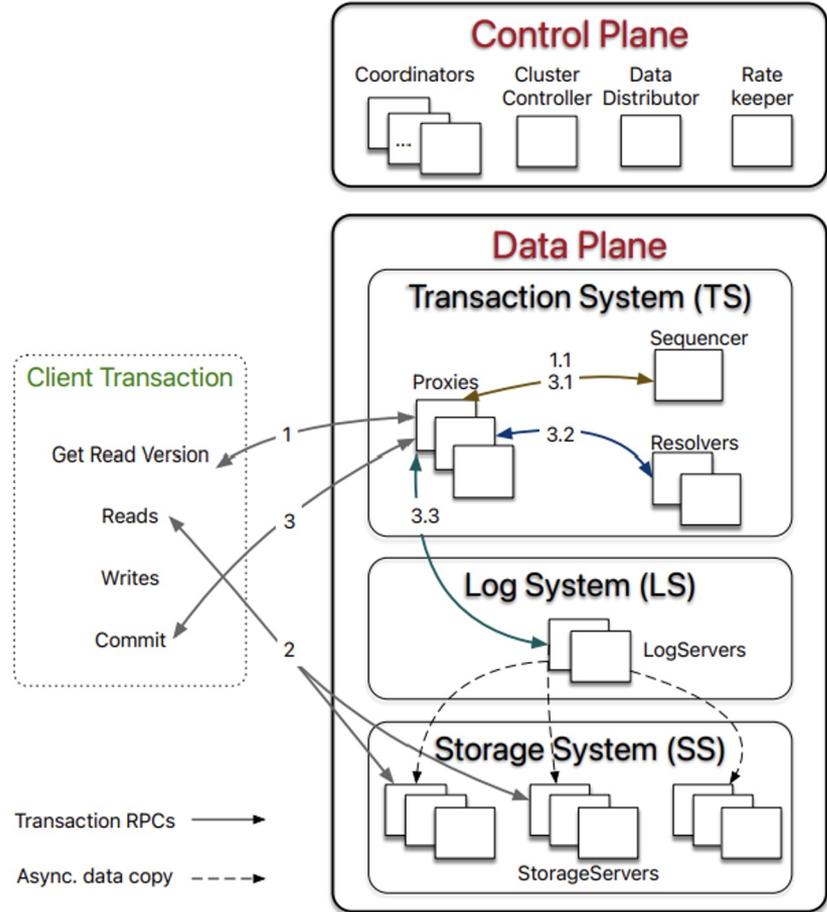
- Background
- Design
- Contributions
- Limitations

Background: why a transactional key value store?



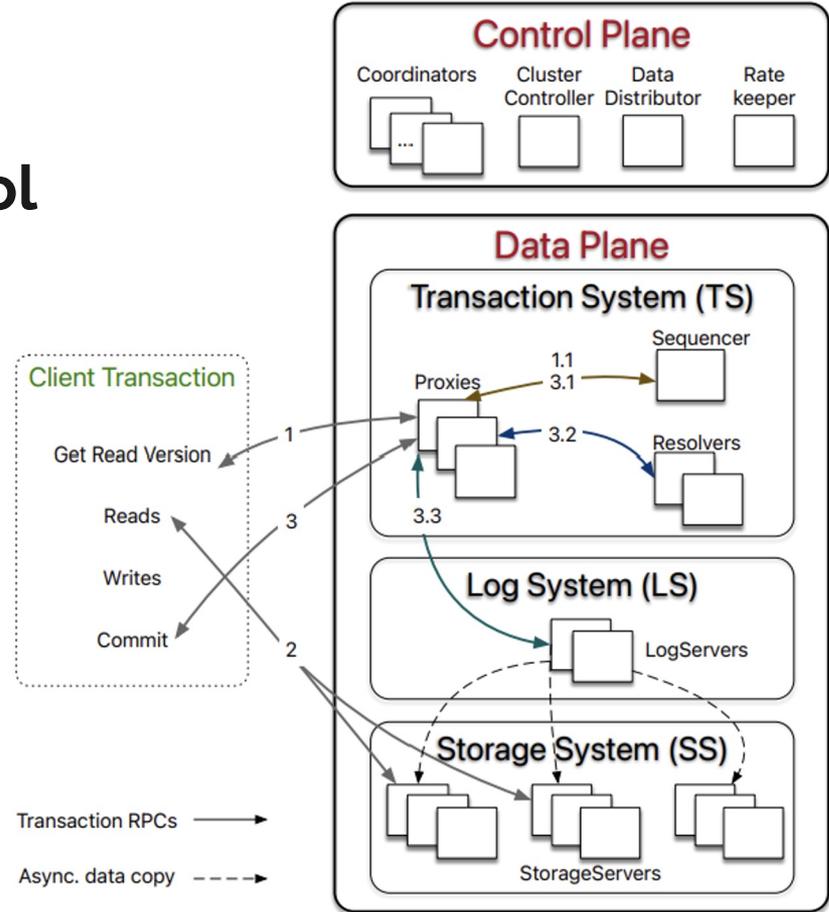
Design: Architecture

- Coordinators: Paxos
- Transaction System: Stateless
- Log System: Stateful
- Storage System: Stateful



Design: Concurrency Control

- Strict Serializability
- Backward Validation OCC
- MVCC
- Commit Version :
 - prev version
 - current version



Design: Concurrency Control

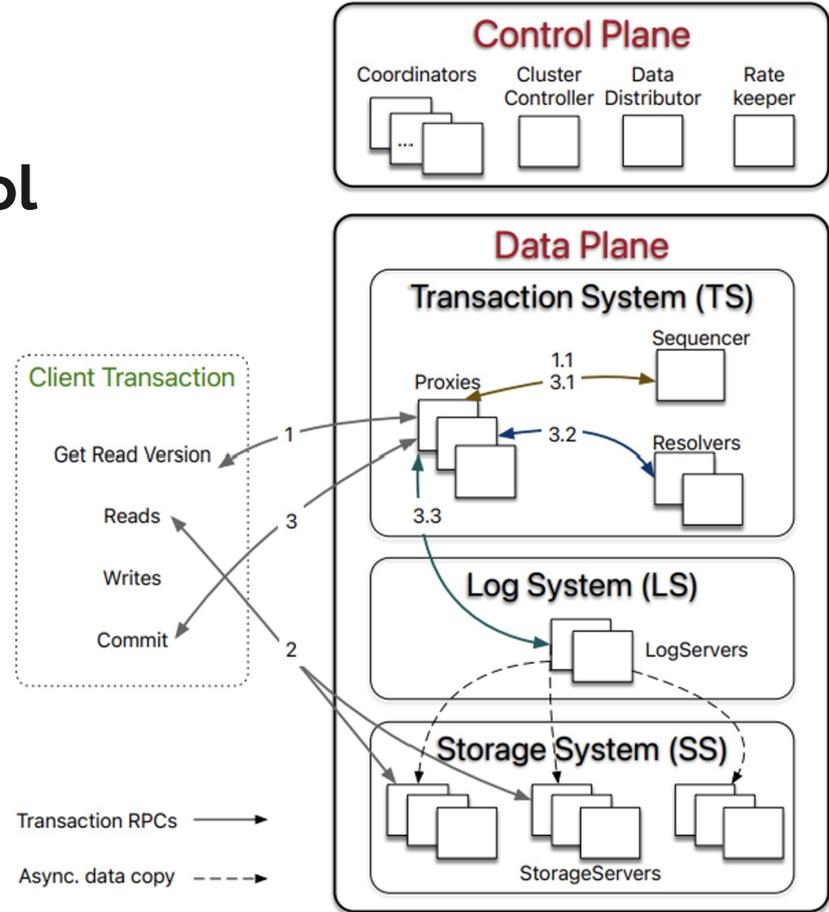
Algorithm 1: Check conflicts for transaction T_x .

Require: $lastCommit$: a map of key range \rightarrow last commit version

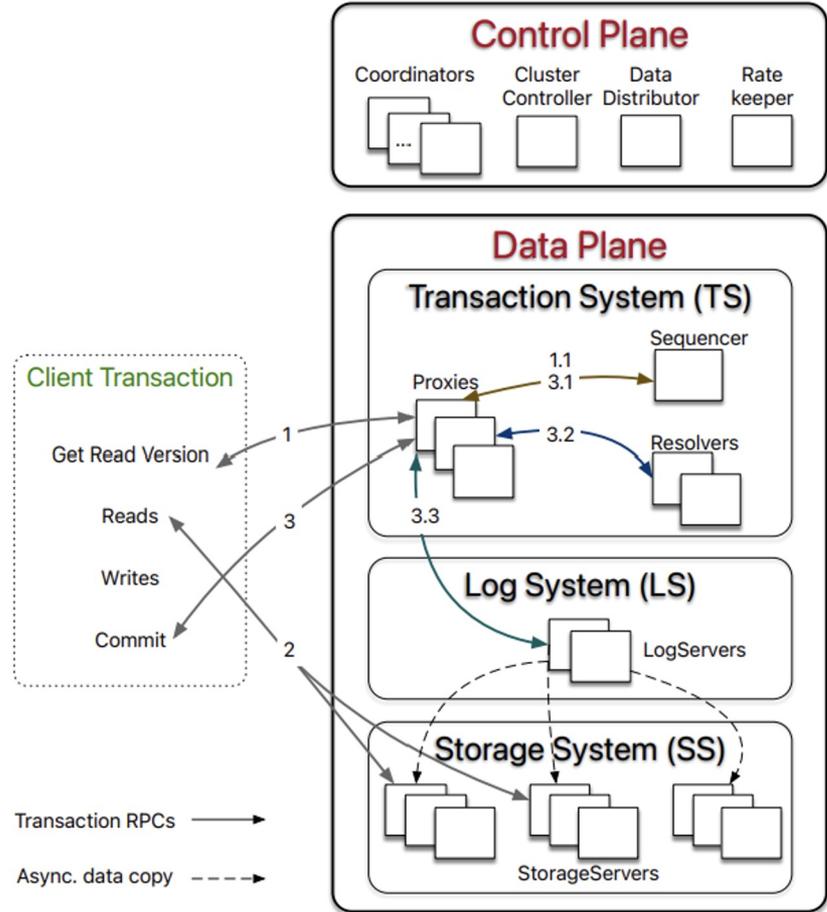
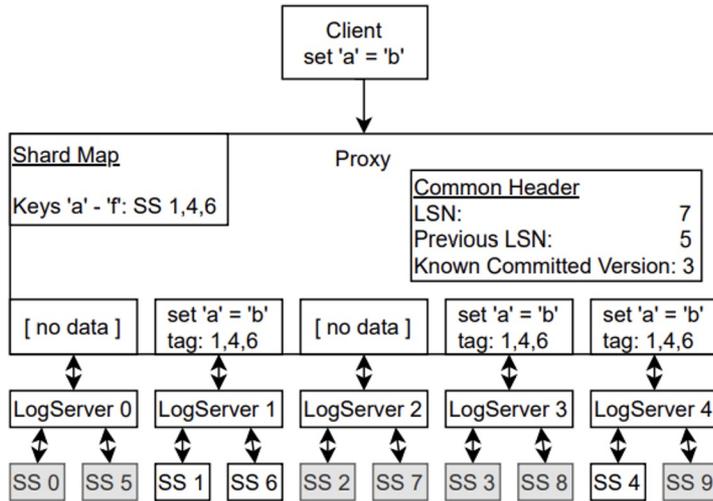
```
1 for each range  $\in R_r$  do
2   ranges = lastCommit.intersect(range)
3   for each r  $\in$  ranges do
4     if lastCommit[r] >  $T_x.readVersion$  then
5       return abort;
```

// commit path

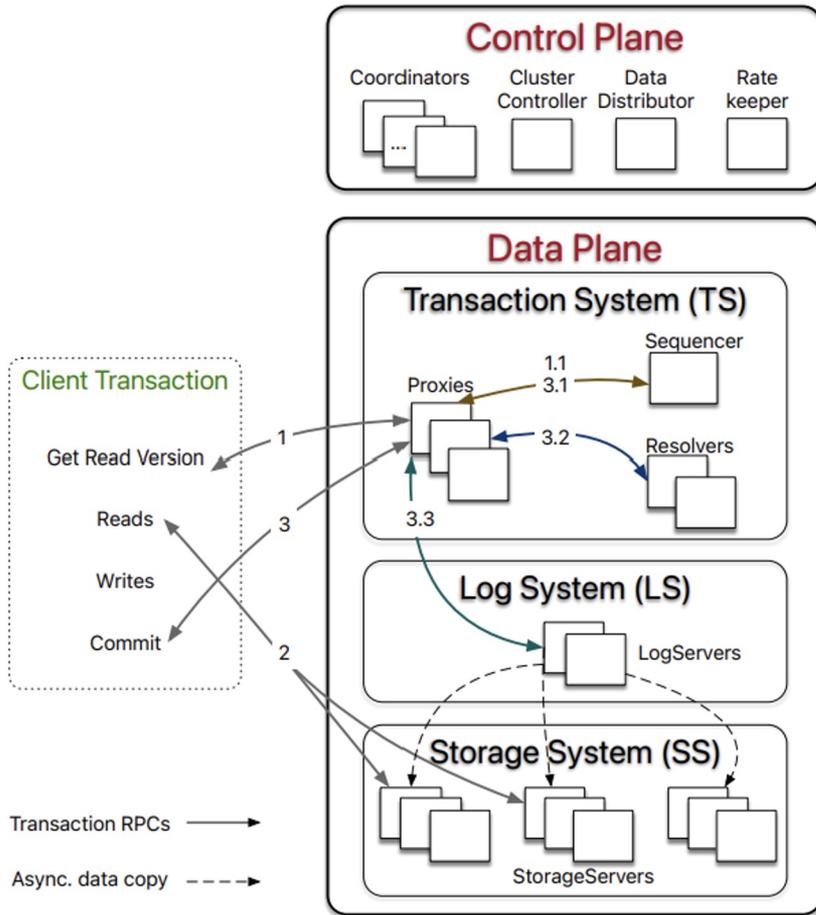
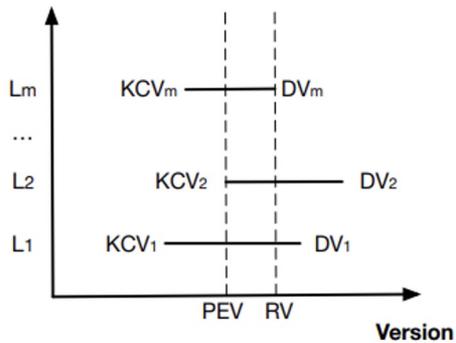
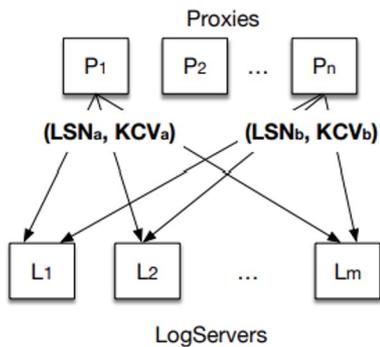
```
6 for each range  $\in R_w$  do
7   lastCommit[range] =  $T_x.commitVersion$ ;
8 return commit;
```



Design: Log System



Design: Recovery



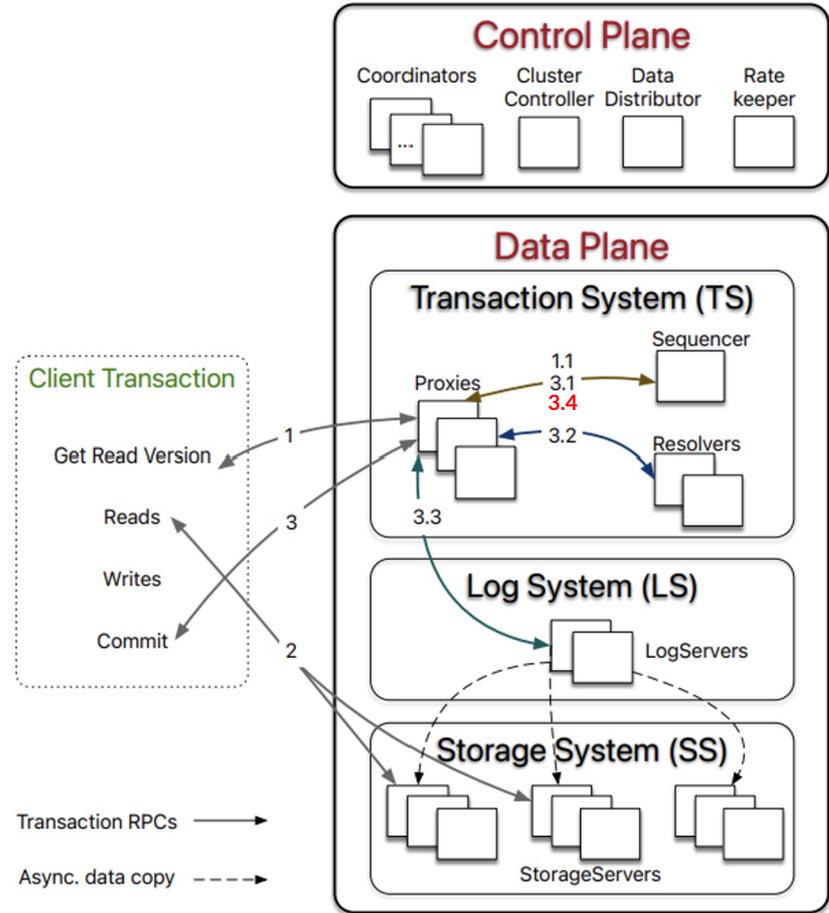
Design: Tricky Problem

Assign versions in Sequencer?

- Strict Serializability
- SS Aggressive Fetching

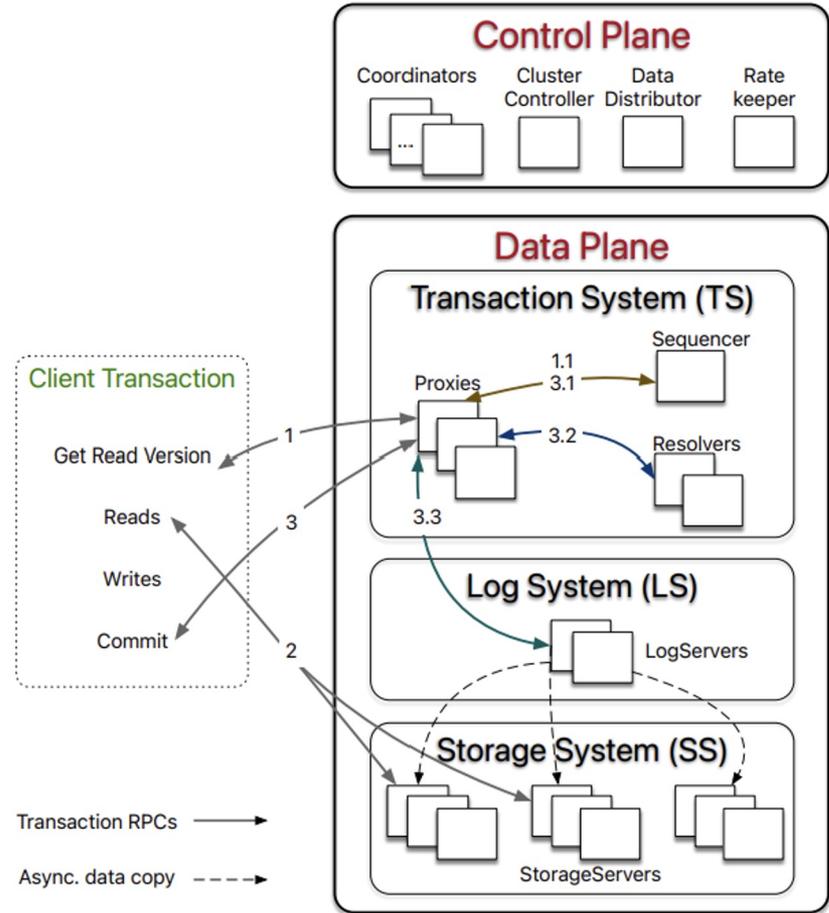
Solution:

- Increasing Commit Versions
- Lasted Committed Version (as Read Version)



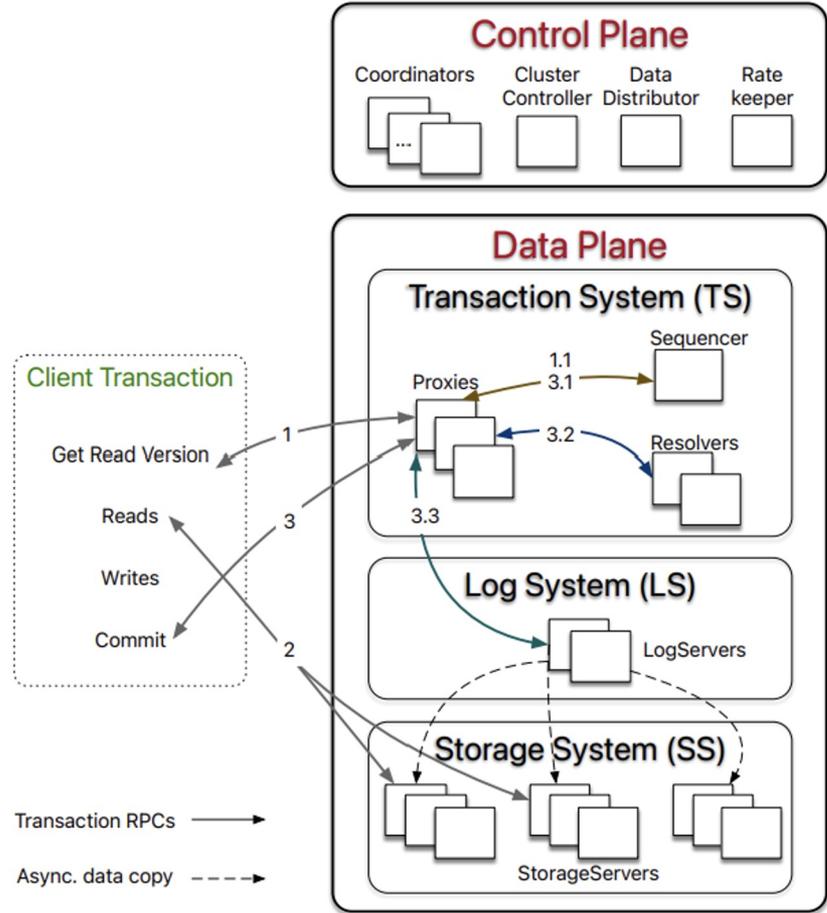
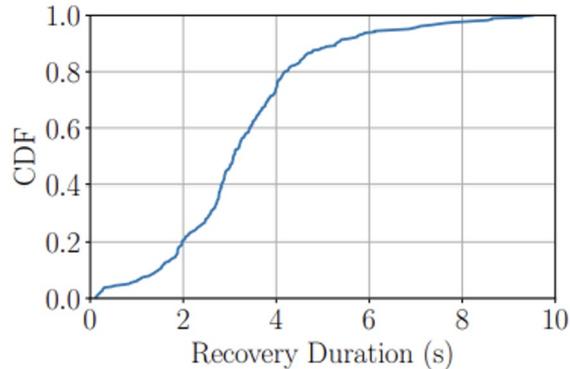
Contributions

- OCC vs Two Phase Locking
- Fast Recovery vs Always On
- Optimize For Happy Case



Limitations

- Transaction Size (OCC)
- 5s MVCC Window
- No RW Transactions in Recovery



Thanks!
And Questions?

